

Building Composable Domain-specific Language Extensions for Java

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Extensible Languages: motivation and goals

- Domain specific languages (DSLs) provide
 - high-level notations for abstractions
 - analyses (e.g. error checking at domain level),
 - optimizationsspecific to a problem domain.
 - But, problems often cover more than one domain.
- Libraries are “composable” --- programmers import the ones they want to use, but these have none of the advantages that DSLs provide.
- Extensible languages in which
 - language extensions add domain-specific language features
 - with composable extensions, programmers can import the set of features needed to address all aspects of a particular task.

Example: SQL embedded in Java

- In JDBC, SQL commands are sent as **strings** to the database server via a connection.
- One writes code like the following:

```
Int value = 100 ;  
ResultSet rs ;  
Statement stmt = conn.createStatement() ;  
rs = stmt.execute("SELECT cust_name FROM " +  
    "customers WHERE quantity > " + value );
```
- SQL error checking and optimization is done by the server at **run-time**, not by the compiler at compile-time.
- A better solution: extend Java with SQL language constructs.

Example: SQL embedded in Java

- Using an extensible language . . .

```
Int value = 100 ;  
ResultSet rc ;  
rc = using conn query  
    ( SELECT cust_name FROM  
      customers WHERE quantity > value ) ;
```

- Here `on ... query` is a new expression which takes a data base connection and an SQL query expression.
- An extended language processor can statically type check the SQL queries.
- Requires **semantic analysis** and **translation** to Java.
- This is not LINQ or SQLJ.

Example: Computational Geometry

- Geometric algorithms make extensive use of **primitive expressions** that return a qualitative result over geometric entities
 - *e.g.* is point **p** inside circle (c,r)
 - *e.g.* is point **p** to the left or right of vertical line **l**
 - implemented by taking the sign of an expression “ $\text{sign}(z - x * y)$ ”
- Writing **efficient robust CG** programs is **difficult**.
 - Round off errors due to limited precision numbers.
 - Import fast unbounded-precision integers from LN
 - Degeneracies in the input data – point **p** is on line **l**
 - Perform transformations to handle degeneracies
- Can **statically calculate** size needed for intermediate values.
- Again, require **semantic analysis** and **translation**.

Challenges to building extensible languages

- Composable specifications of **feature semantics**
 - Attribute grammars + forwarding + . . .
 - Get both **explicit** and **implicit** (via translation) specification of semantics.
 - Tool support: **Silver** – extensible AG system.
- Composable specifications of **feature syntax**
 - Context-aware scanning with LR-parsing
 - Deterministic, yet handles large class of languages
 - Monolithic and modular determinism analyses
 - Tool support: **Copper** – parser and scanner generator.

AbleJ

- An extensible implementation of Java
- Built from declarative specifications used by Silver and Copper.
- Several composable language extensions
 - SQL, Computational Geometry, Condition Tables, Pizza constructs, . . .
- Supports automatic composition of language extensions to create domain-adapted versions of Java.
- Specifies all of Java 1.4 syntax and several semantic analyses such as type checking.

More information

- ... on Tuesday and Thursday
- ... Thanks for your attention.