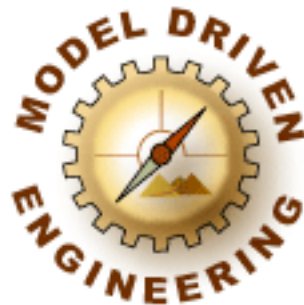
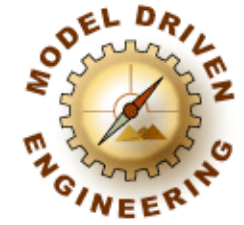

On the use of graph transformations for **model refactoring**

Tom Mens



Software Engineering Lab
University of Mons-Hainaut
<http://w3.umh.ac.be/genlog>





Tutorial outline

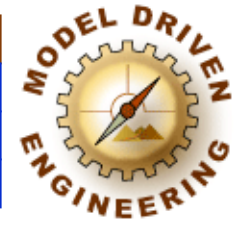
- Introduction
 - What is model-driven engineering, model transformation, model refactoring?
 - Where does graph transformation fit in?
- Graph transformation theory
- Graph transformation experiments
 - In *Fujaba*: model refactoring plug-in
 - In *AGG*: critical pair analysis
 - With pencil and paper: behaviour preservation
- Conclusion

Introduction



Model-driven engineering

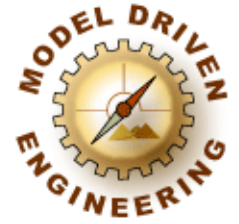
Introduction
GT theory
Experiments
Conclusion



- Goal: Raise the level of software development from *source code* to *models*
 - models = software artifacts at higher level of abstraction
 - e.g. UML diagrams = design models
- Principle: "Everything is a model"
 - Uniform approach to all kinds of software artifacts
 - source code is a kind of model
 - the syntax of a model is described by a *metamodel*

Model Transformation

Introduction
GT theory
Experiments
Conclusion

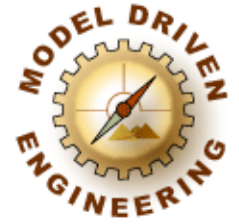


- Goal: Apply transformation techniques to modify, refine and evolve models
- Classification of model transformations

	<i>endogenous</i>	<i>exogenous</i>
<i>horizontal</i>	refactoring	language migration, bridging techn. spaces
<i>vertical</i>	formal refinement	code generation

Model Transformation

Introduction
GT theory
Experiments
Conclusion



- Endogenous versus exogenous
 - Endogenous transformations
 - transformations between models expressed within the same metamodel
 - Exogenous transformations
 - transformations between models expressed in different metamodels
- Horizontal versus vertical
 - Horizontal transformation
 - transformation between models residing at the same level of abstraction
 - Vertical transformation
 - transformation between models at different abstraction levels

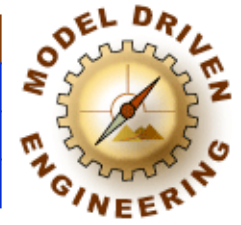
Model Evolution

Introduction

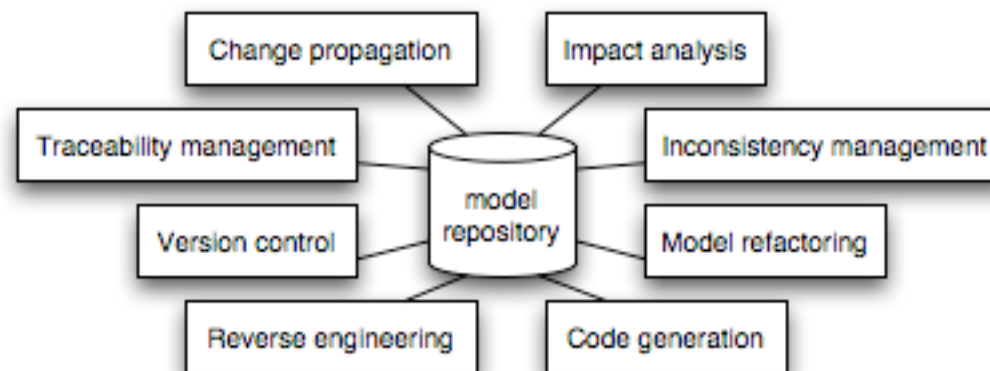
GT theory

Experiments

Conclusion



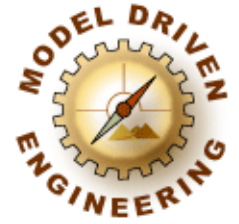
- Goal: Provide support for software evolution at the level of models



- Better tool support needed for all these activities
- Formalisms can be helpful for some of these tools

Model Refactoring

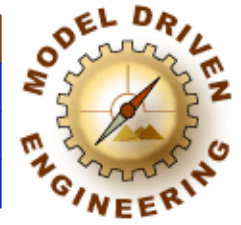
Introduction
GT theory
Experiments
Conclusion



- Goal: special kind of model evolution that improves the structure of the model, while **preserving** (certain aspects of) its **behaviour**
- Model refactoring is an example of an endogenous, horizontal model transformation
- Model refactoring is based on the idea of program refactoring
 - *"the process of changing a program in such a way that it does not alter the external behavior of the code, yet improves its internal structure" [Martin Fowler, 1999]*

Model Refactoring

Introduction
GT theory
Experiments
Conclusion

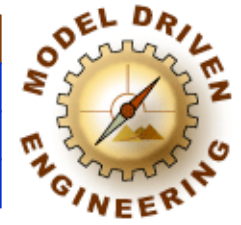


- Model refactorings can be applied to different views of a UML model
 - class diagrams
 - sequence diagrams
 - statecharts
 - activity diagrams
- Some model refactorings have been proposed by Boger et al.

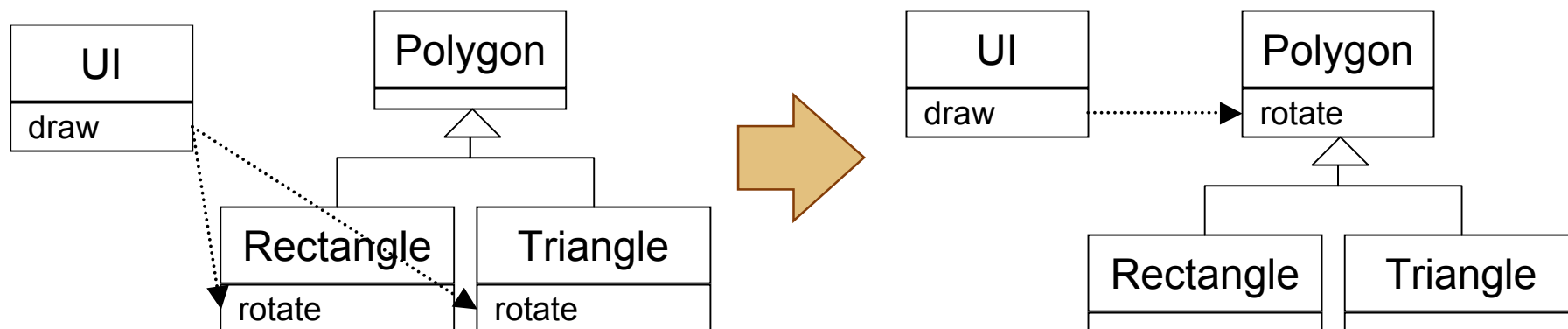
M. Boger, T. Sturm, P. Fragemann. *Refactoring Browser for UML*.
Proc. 3rd Int'l Conf. on eXtreme Programming, pp. 77-81, 2002

Model Refactoring

Introduction
GT theory
Experiments
Conclusion



- Examples of class diagram refactorings
 - Pull Up Method



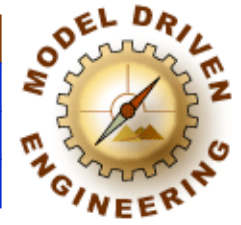
Model Refactoring

Introduction

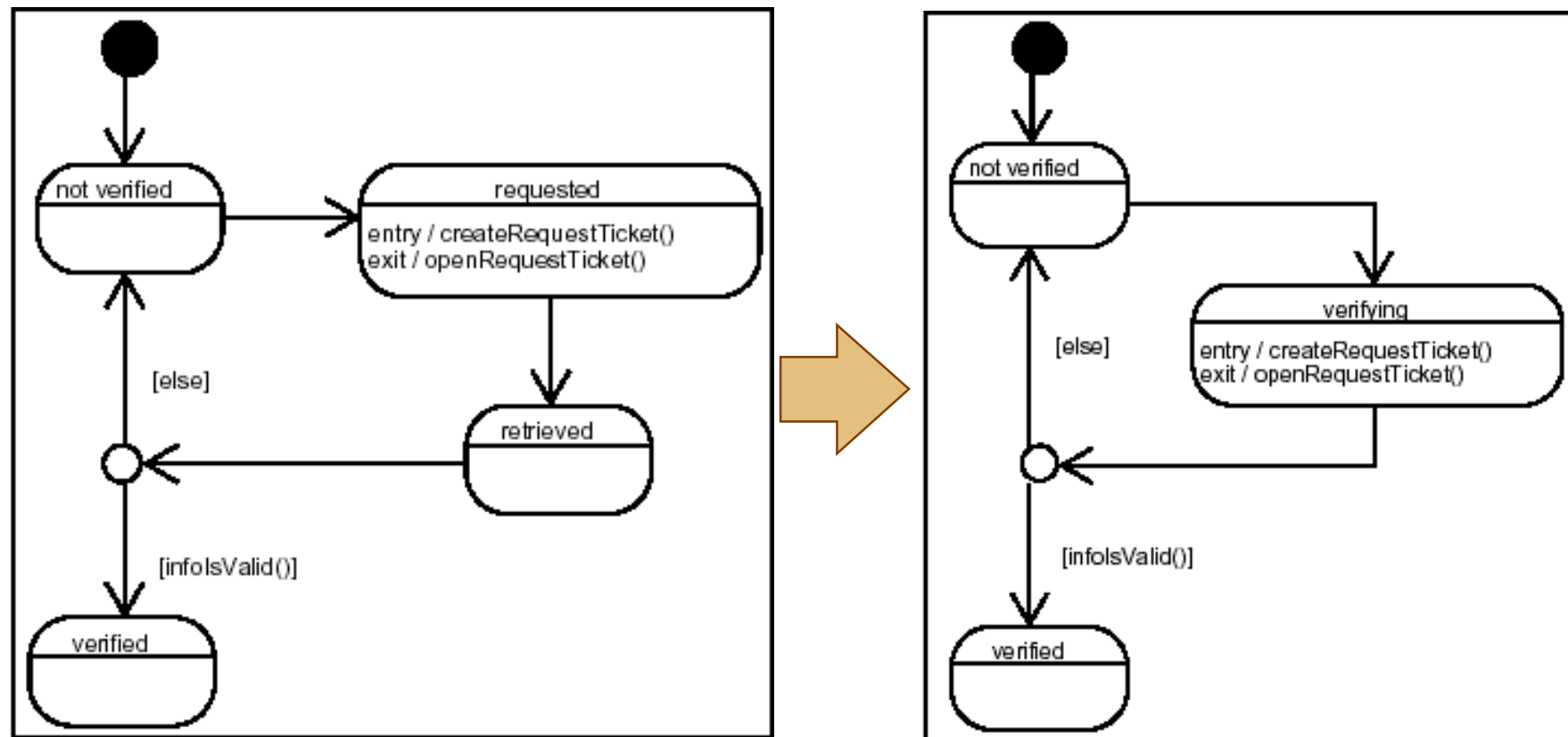
GT theory

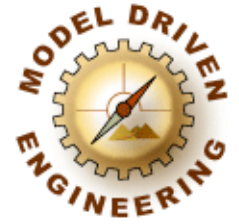
Experiments

Conclusion



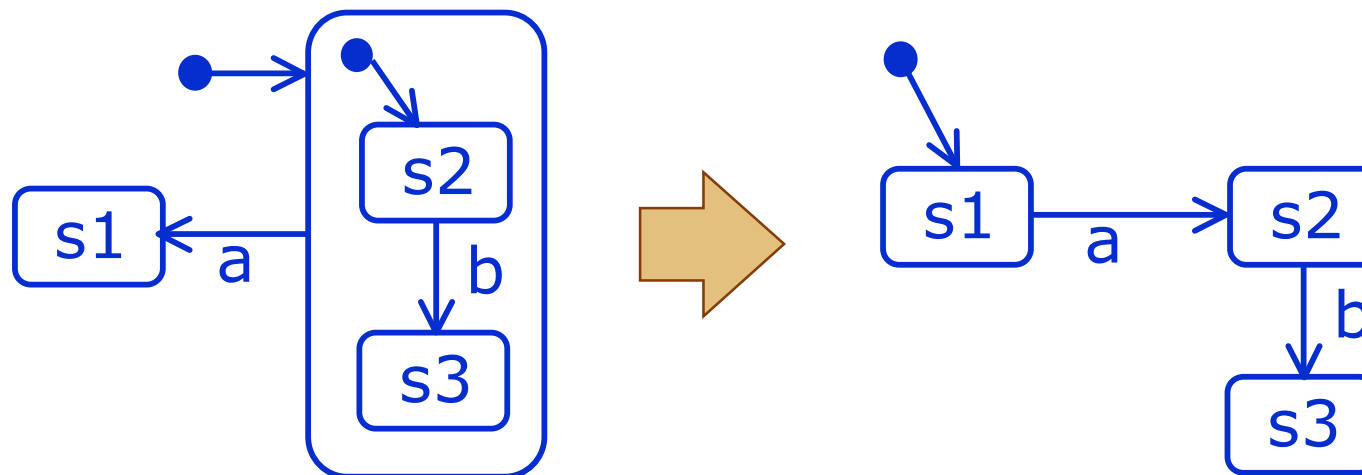
- Examples of statechart refactorings
 - Merge states



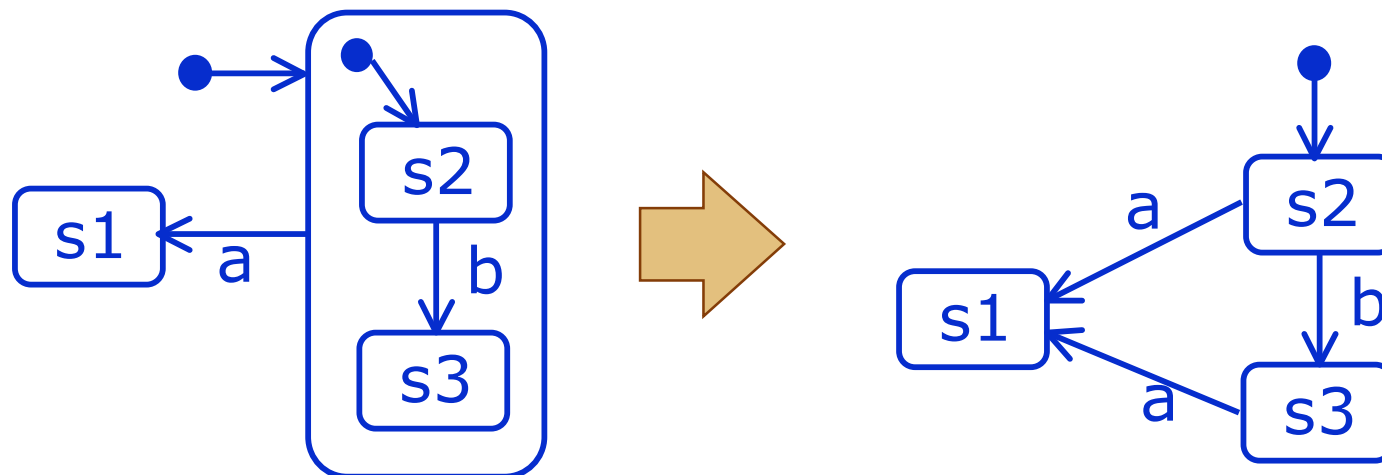


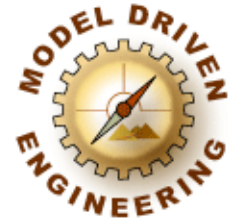
- Examples of statechart refactorings
 - Merge states
 - combine a set of states into a single composite state
 - Decompose sequential composite state
 - remove a composite state but keep its internal states
 - Create composite state
 - create a new composite state and move selected states to the interior
 - Sequentialise concurrent state
 - replace a concurrent state by a product automaton
 - Flatten states
 - see next slide

- Examples of statechart refactorings
 - Flatten states: Incoming transitions
 - Transition from state $s1$ to the boundary of a complex state represents a transition from $s1$ to the initial state of the complex state



- Examples of statechart refactorings
 - Flatten states: Outgoing transitions
 - Transition from boundary of complex state to state *s1* represents corresponding transitions from all substates to *s1*





- Examples of activity diagram refactorings
 - Make actions concurrent
 - Create a fork and a join pseudostate, and move several sequential groups of actions between them, thus enabling their concurrent execution
 - Sequentialize concurrent actions
 - Removes a pair of fork and join pseudostates, and links the enclosed group of actions to another

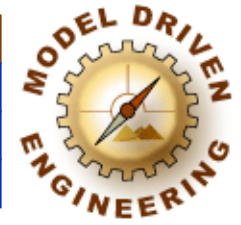
Model Refactoring

Introduction

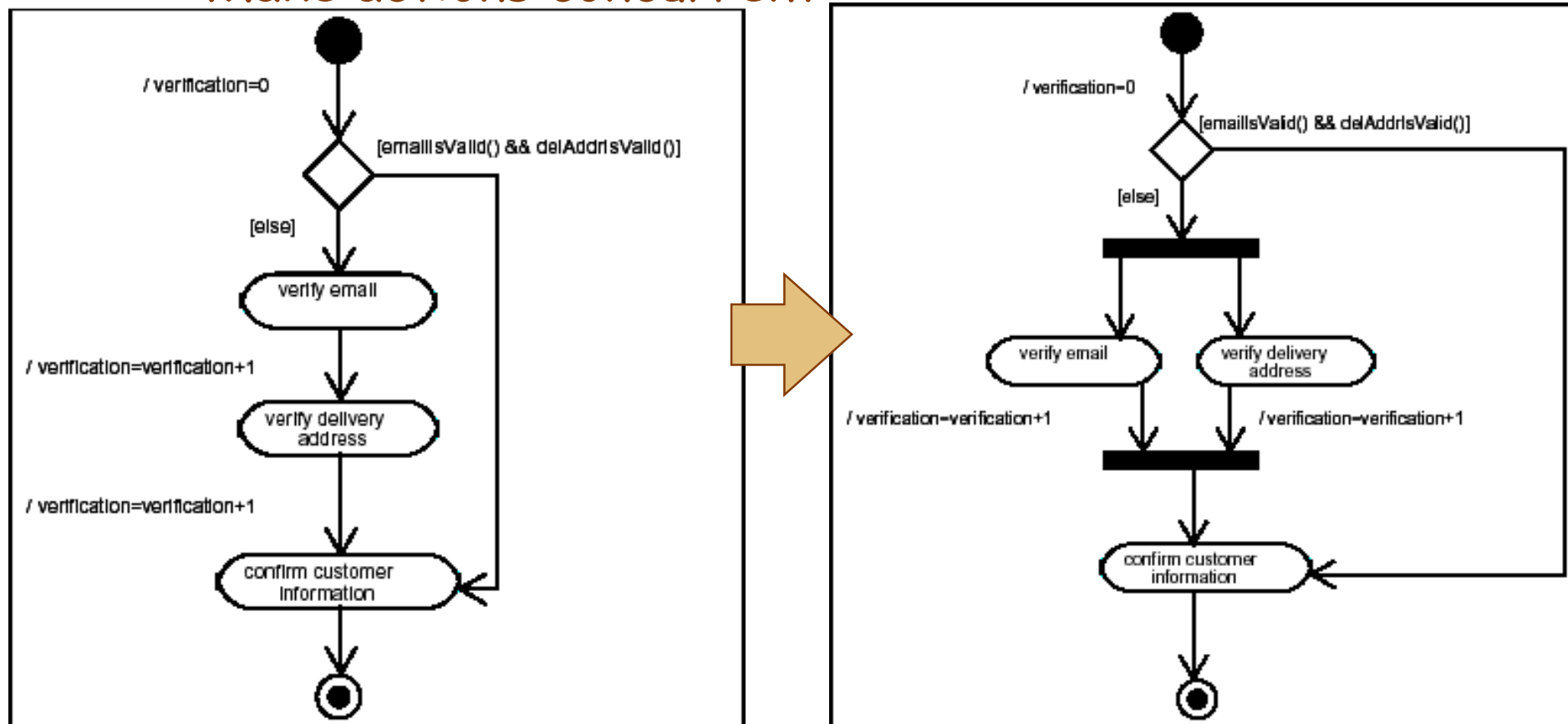
GT theory

Experiments

Conclusion

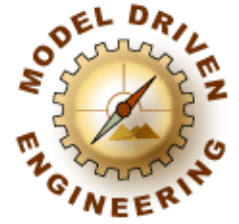


- Examples of activity diagram refactorings
 - Make actions concurrent



Model Refactoring

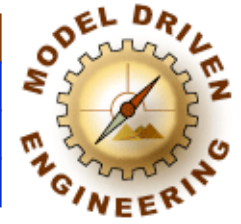
Introduction
GT theory
Experiments
Conclusion



- Other related work
 - Sunye *et al.* [UML 2001]
 - statechart refactorings expressed using OCL pre- and postconditions
 - Van Gorp *et al.* [UML 2003]
 - UML extension to support source consistent refactoring
 - integrated as plug-in in Fujaba tool
 - Correa and Werner [UML 2004]
 - UML refactorings in OCL-script, and extension of OCL
 - ...

Research Context

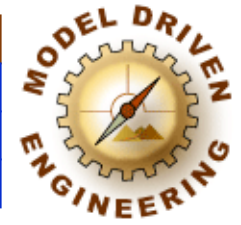
Introduction
GT theory
Experiments
Conclusion



- Several ongoing research projects
 - "*A Formal Foundation for Software Refactoring*"
 - Financed by FWO - Flanders, Belgium
 - Duration: January 2003 - December 2006
 - In collaboration with: Serge Demeyer and Dirk Janssens, University of Antwerp
 - "*Research Center on Structural Software Improvement*"
 - Financed by FNRS-FRFC, Belgium
 - Duration: January 2005 - December 2008
 - In collaboration with: Kim Mens, UCL - Roel Wuyts, ULB
 - For more information, see <http://www.info.ucl.ac.be/ingidocs/people/km/FRFC>

Research Context

Introduction
GT theory
Experiments
Conclusion

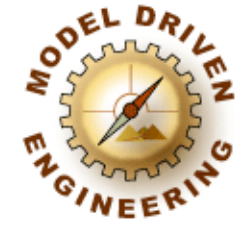


- Several international research networks
 - **ESF** Scientific Network RELEASE
 - « Research Links to Explore and Advance Software Evolution »
 - <http://www.esf.org/release/>



- **ERCIM** Working Group on Software Evolution
 - <http://w3.umh.ac.be/evol>





Tutorial outline

✓Introduction

- Graph transformation theory
- Graph transformation experiments
- Conclusion

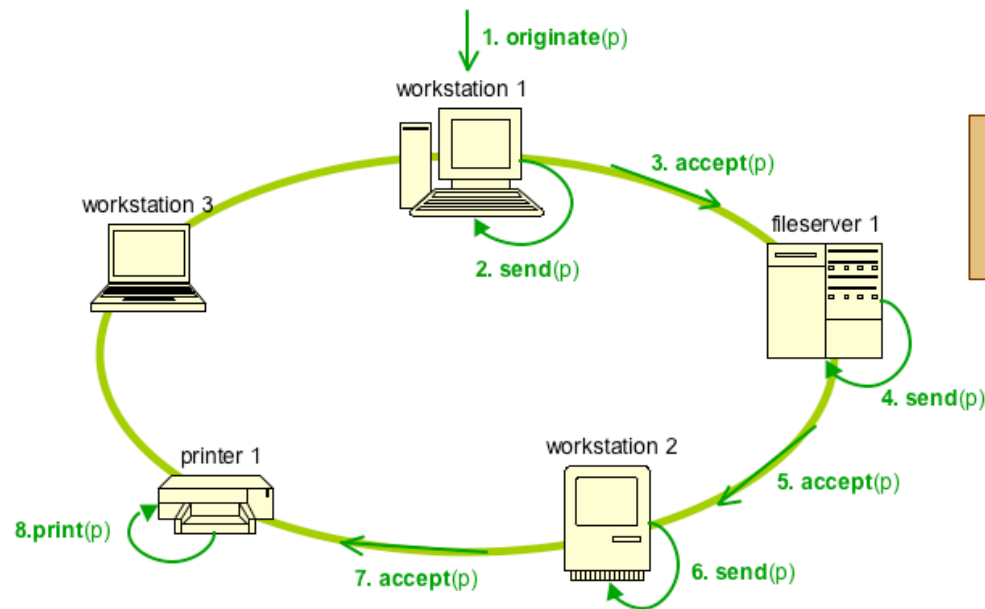
Graph Transformation Theory



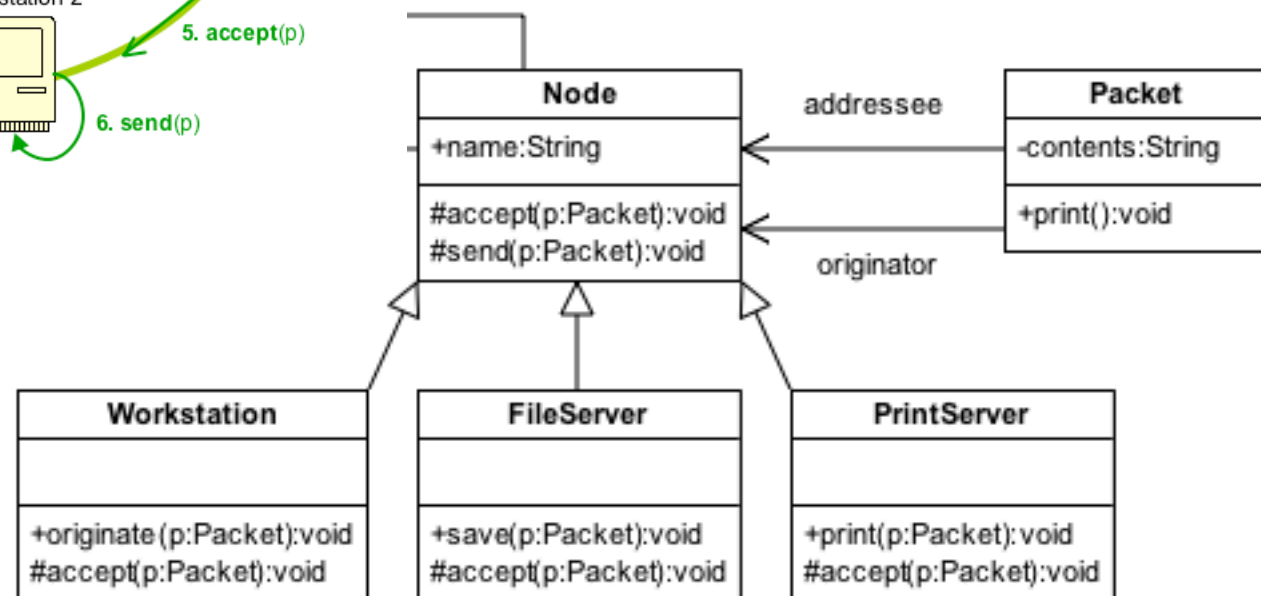
Models are Graphs

- Models can be represented naturally as graphs
 - many diagrams are intrinsically graph-based
 - class diagrams, statecharts, collaboration diagrams, Petri-nets, database schemas

Models are Graphs



Simple example: class diagram
 of a LAN simulation



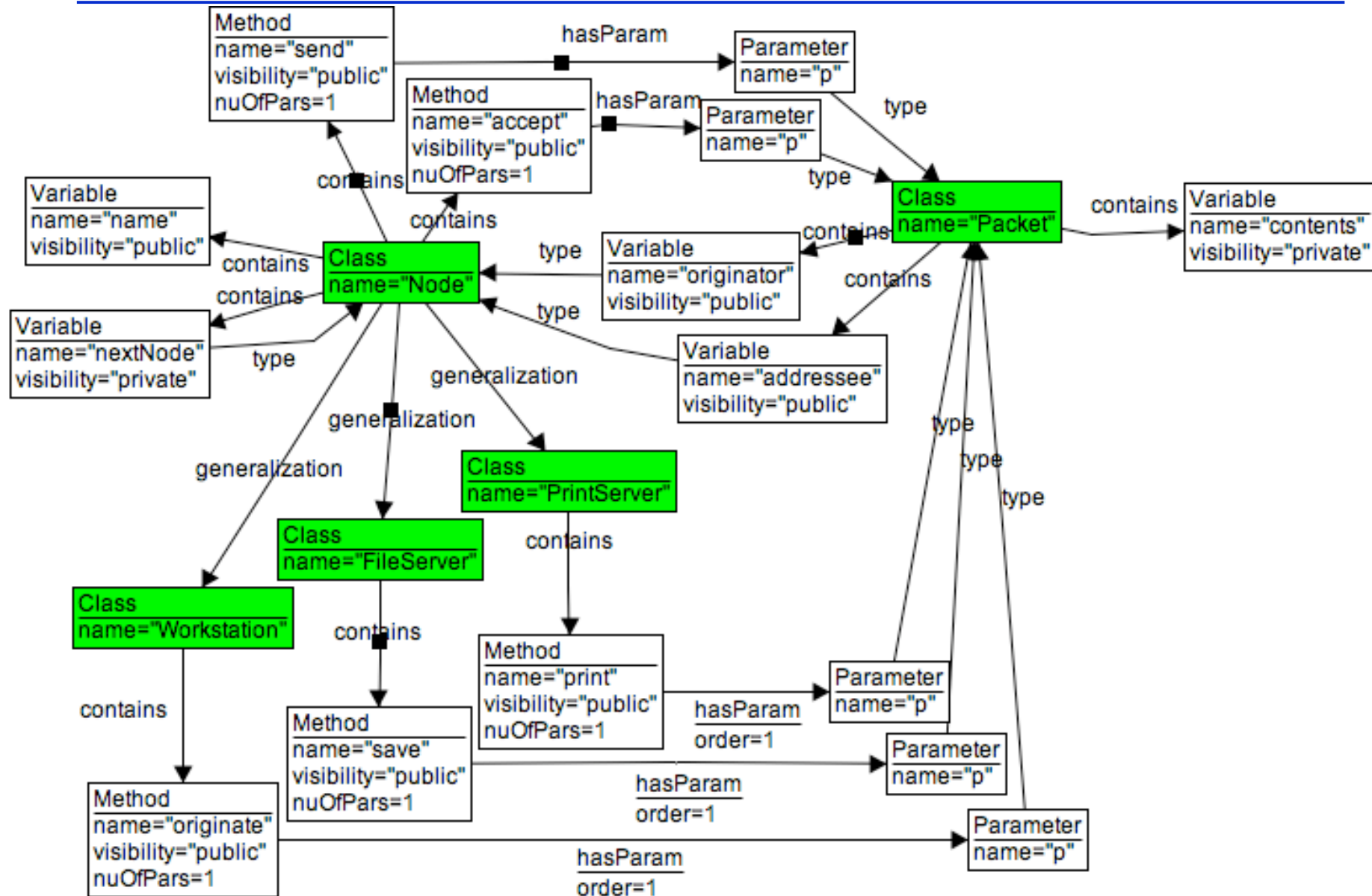
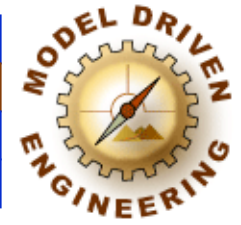
Models are Graphs

Introduction

GT theory

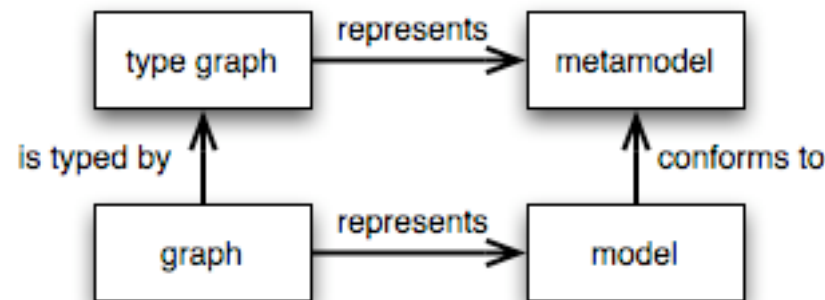
Experiments

Conclusion



Metamodels are type graphs

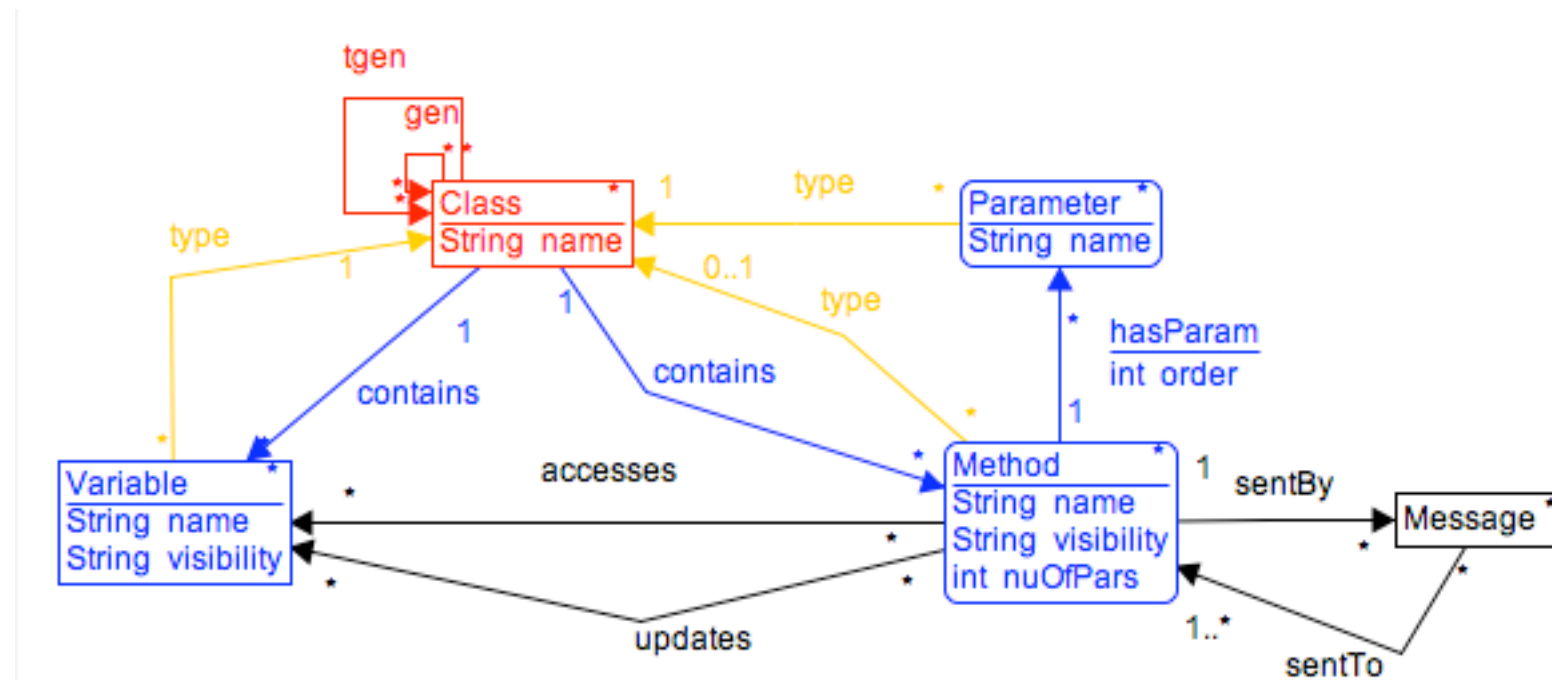
- All models conform to a metamodel that specifies their syntax
- All graphs conform to a type graph that specifies their well-formedness constraints



- Hence, type graphs are the graph-theoretic equivalent of metamodels

Metamodels are type graphs

- Example of a type graph
 - represents a simplified metamodel for UML class diagrams



Definitions

- A (directed) graph is an algebraic structure $G = (V, E, s: E \rightarrow V, t: E \rightarrow V)$
- A graph homomorphism is a mapping $h: G_1 \rightarrow G_2$ where $h = (h_V: V_1 \rightarrow V_2, h_E: E_1 \rightarrow E_2)$ and h_E preserves source and target nodes
- A graph G is typed by a type graph TG if there is a homomorphism $g: G \rightarrow TG$
- Direct extension of definitions to labeled graphs
 - where each node and edge may be labeled

Definitions

- A *labeled graph* is a graph where each node and edge is labeled over an alphabet L
 - labeling function $l: VUE \rightarrow L$
- An *attributed graph* is a graph where nodes and edge are labeled over an abstract data type

Method
int nuOfPars

Class
String name

type level

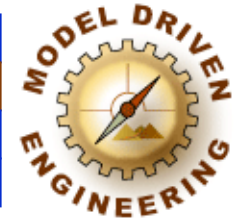
<u>m</u> : Method
int nuOfPars

<u>c</u> : Class
name = "Packet"

instance level

Model transformations are Graph transformations

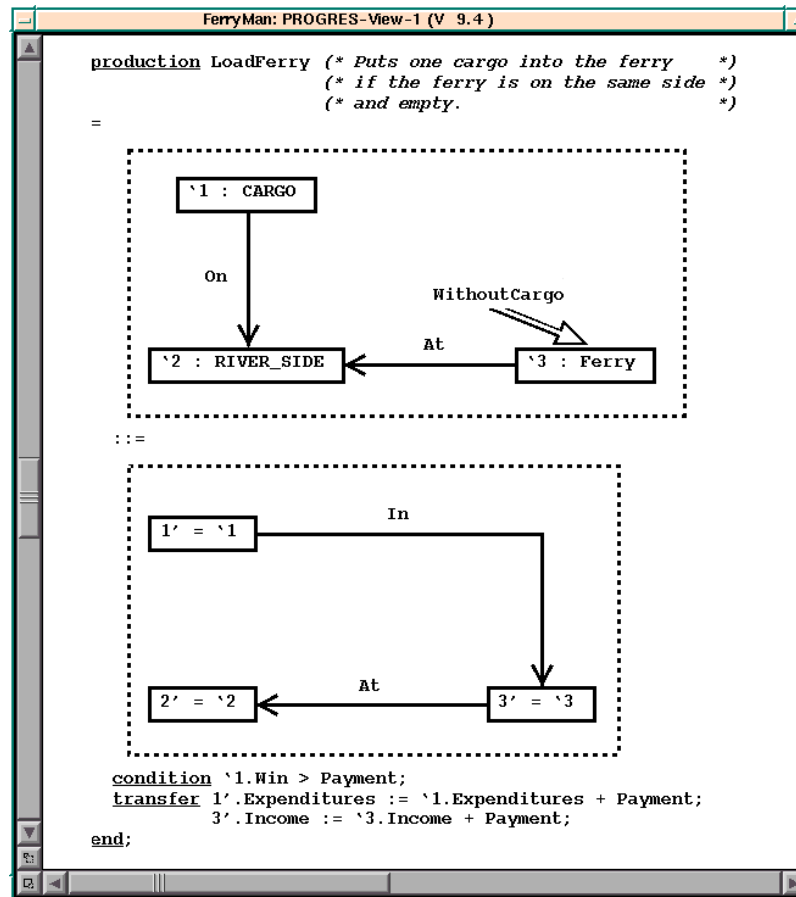
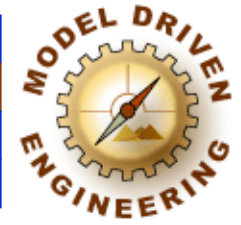
Introduction
GT theory
Experiments
Conclusion



- Model transformations can be naturally represented as graph transformations
- **GT theory** offers many theoretical results that can help during analysis
 - type graph, negative application conditions, parallel and sequential (in)dependence, confluence, critical pair analysis
- **GT tools** allow us to perform concrete experiments
 - Fujaba, AGG, Progres, ...

GT tools - PROGRES

Introduction
GT theory
Experiments
Conclusion



- Graphical/textual language to specify graph transformations
- Graph rewrite rules with complex and negative conditions
- Cross compilation in Modula 2, C and Java

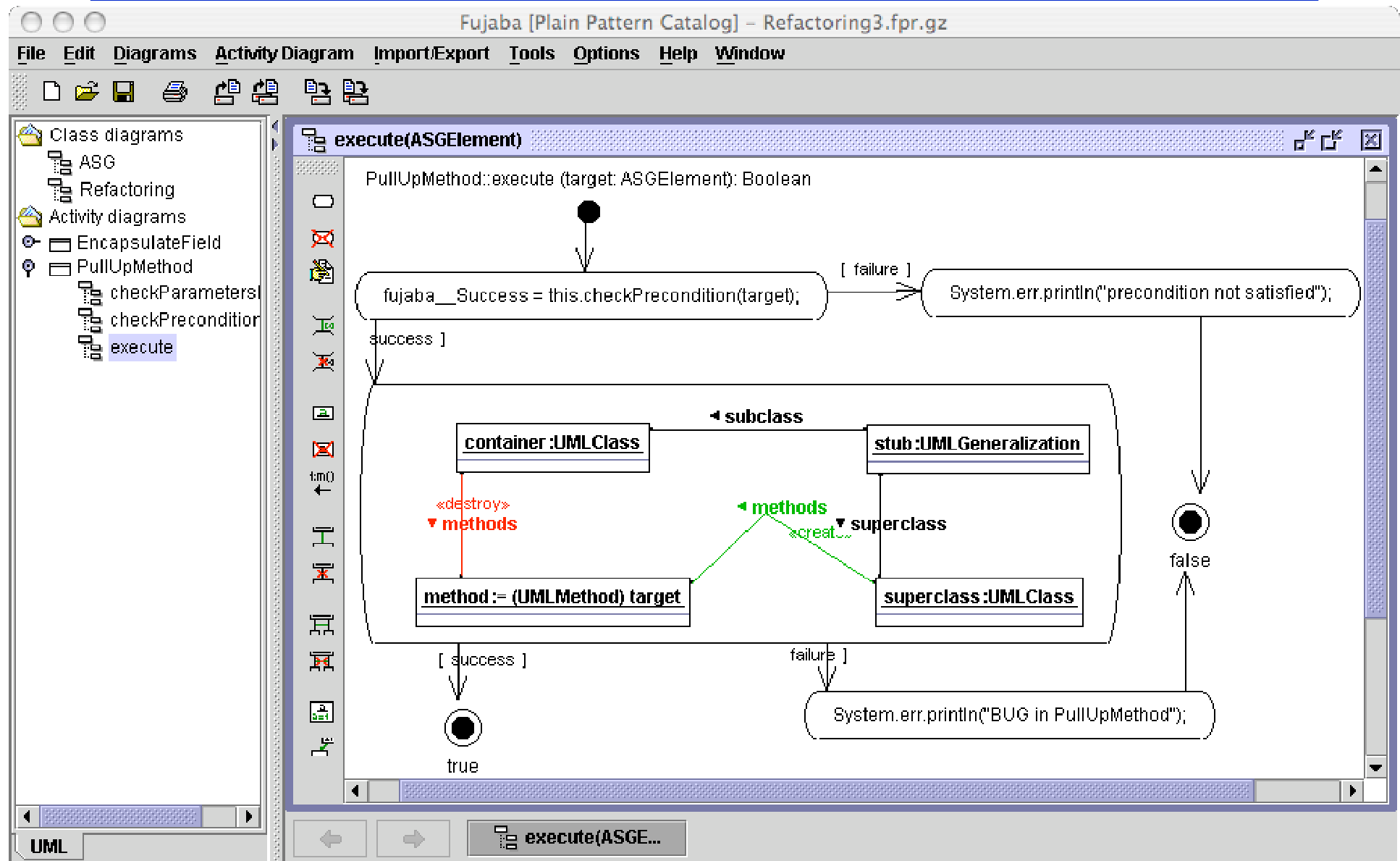
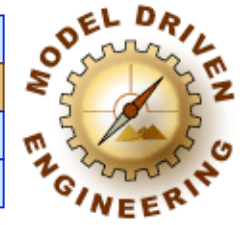
GT tools - Fujaba

Introduction

GT theory

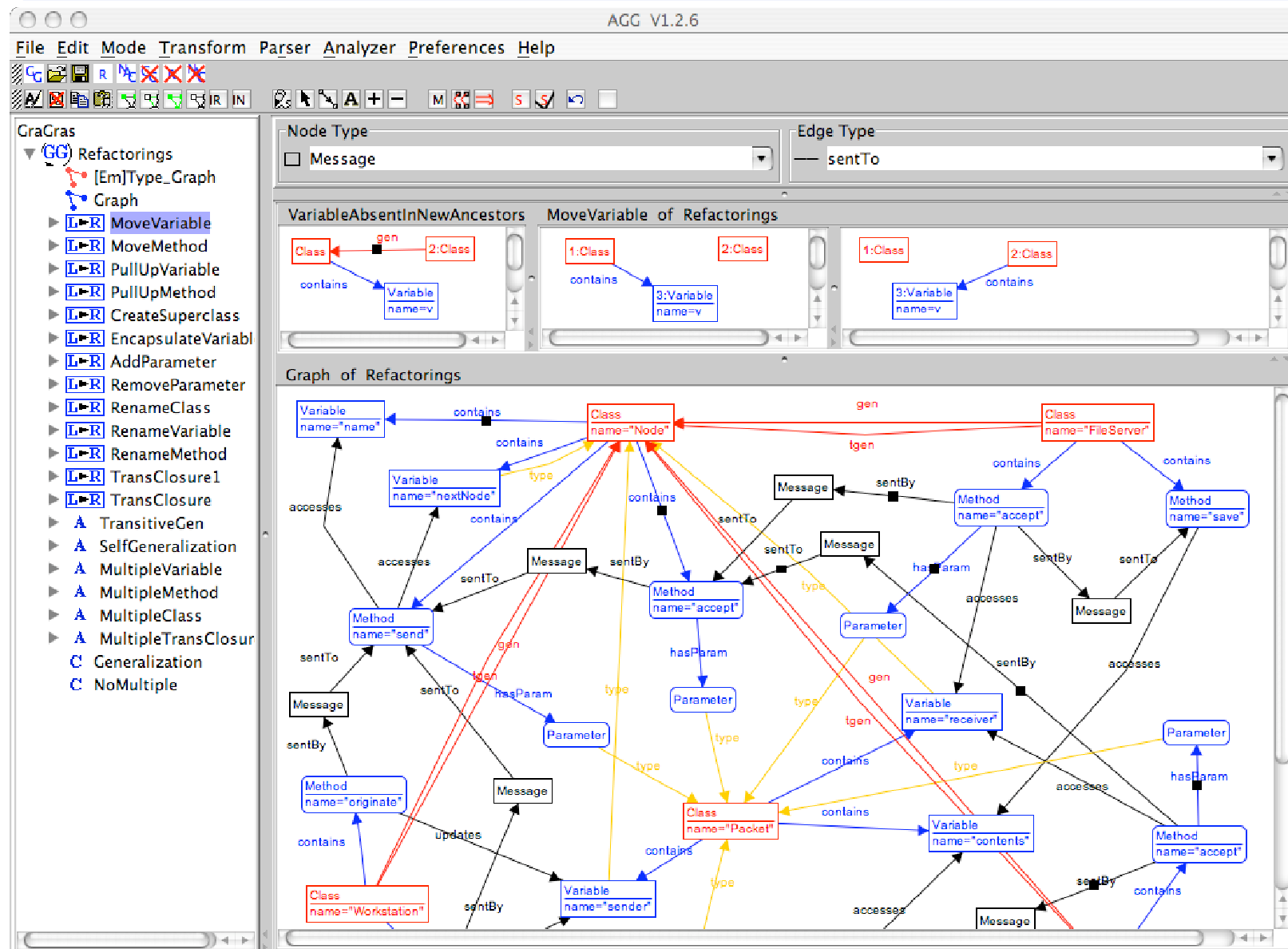
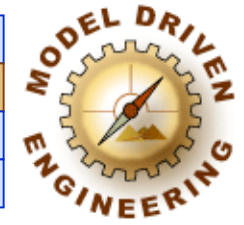
Experiments

Conclusion



GT tools - AGG

Introduction
GT theory
Experiments
Conclusion



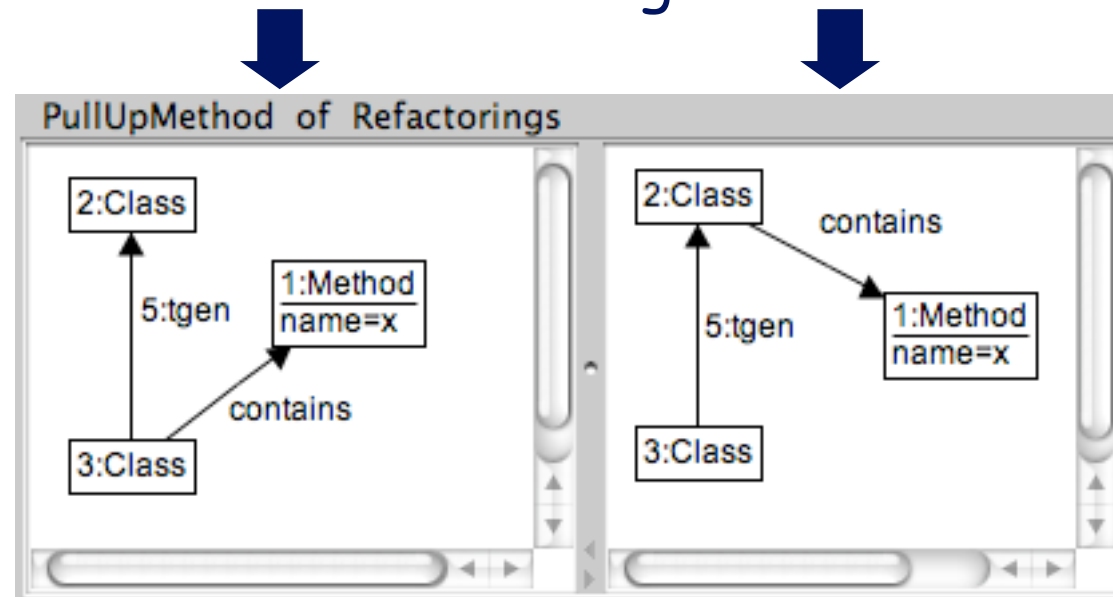
Graph Production

- A graph production $p: L \rightarrow R$ is a structure-preserving partial mapping between (directed, labeled, typed) graphs
 - Preserves sources and targets of edges
 - Preserves node and edge types
 - Preserves node and edge labels
 - Partial means that nodes or edges may be deleted

Graph Production

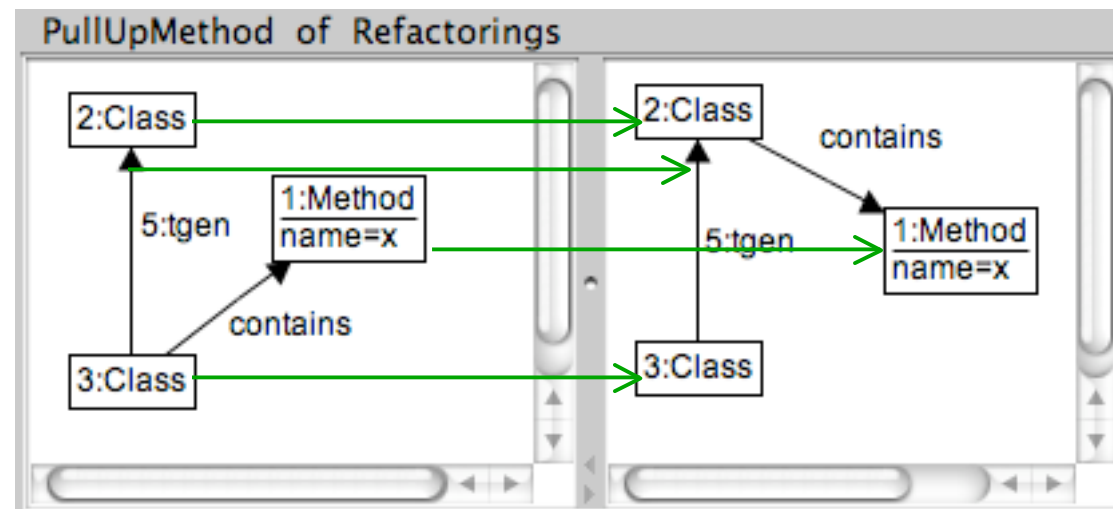
- Example: Pull Up Method refactoring

left-hand side L right-hand side R



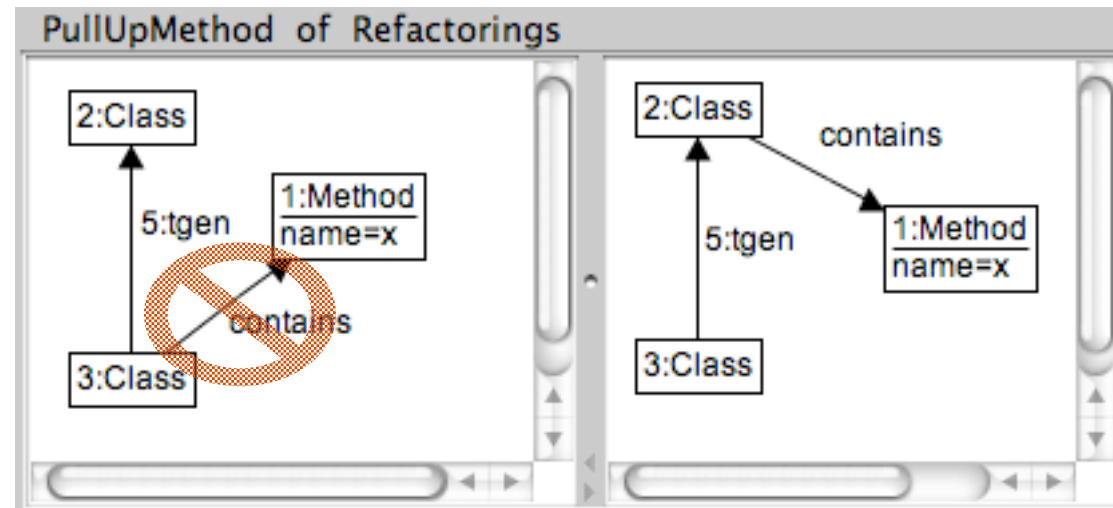
Graph Production

- Example: Pull Up Method refactoring
 - Some nodes and edges are **preserved**



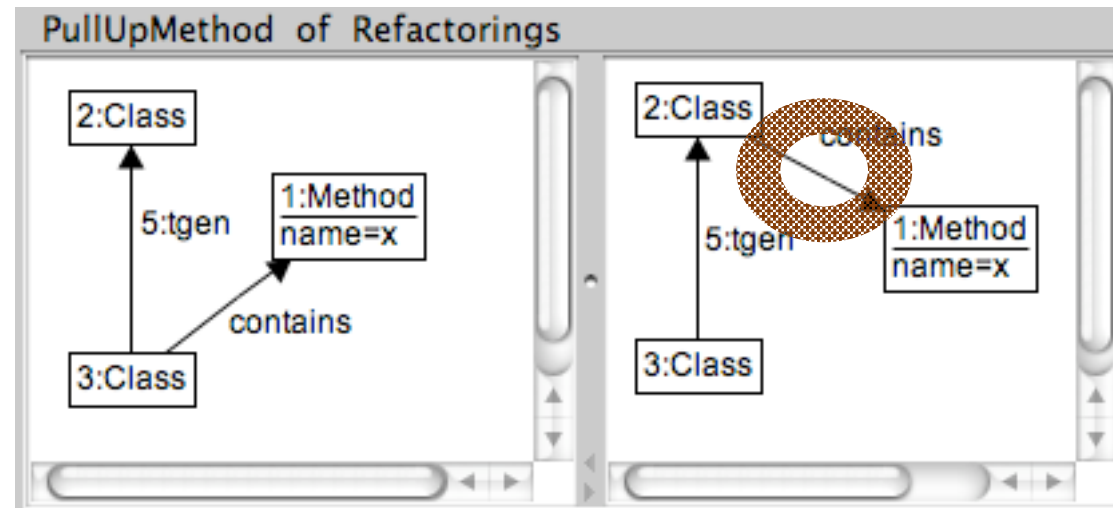
Graph Production

- Example: Pull Up Method refactoring
 - Some nodes and edges are **deleted**



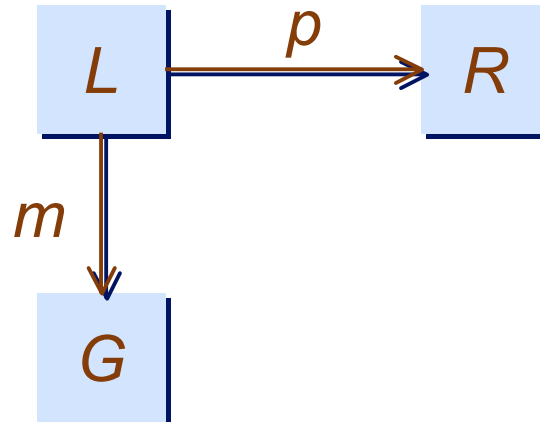
Graph Production

- Example: Pull Up Method refactoring
 - Some nodes and edges are **added**



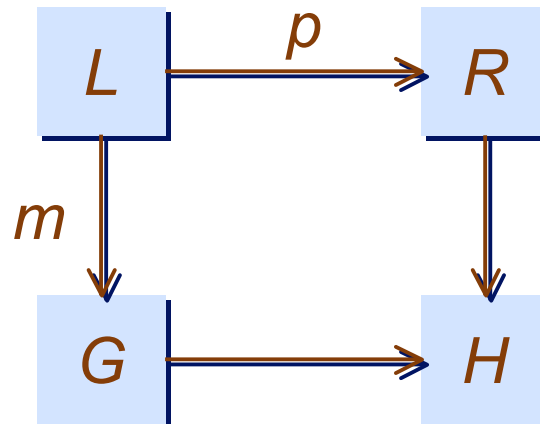
Graph Transformation

- A graph transformation $t: G \Rightarrow H$ is the application of a graph production $p: L \rightarrow R$ that is matched in the context of a given graph G
 $t = (p, m)$ where $m: L \rightarrow G$ is an injective graph morphism (match)



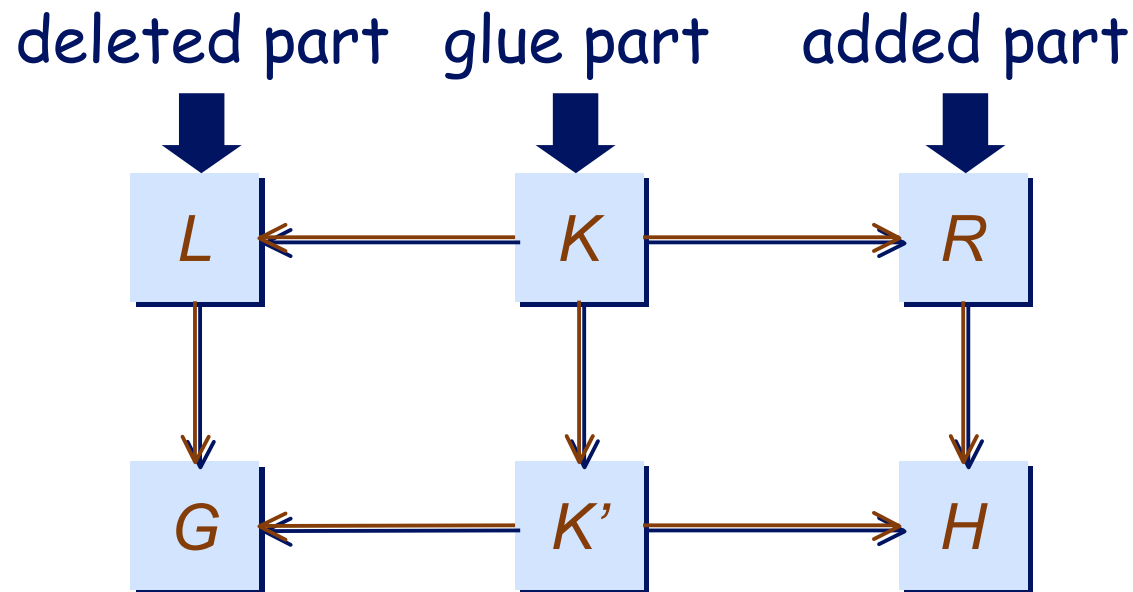
Graph Transformation

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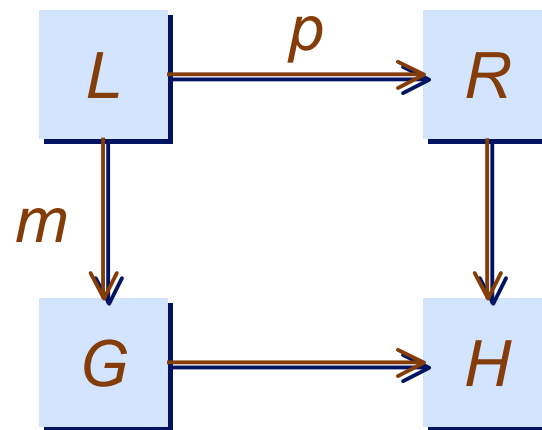
Graph Transformation

- Alternative definition: algebraic double-pushout (DPO) approach
 - Explicit presentation of intersection $K = L \cap R$



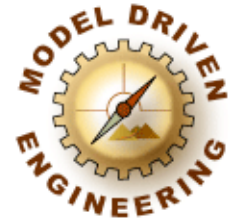
Graph Transformation Step

- **Operational description**
 - prepare transformation by
 - selecting rule $p: L \rightarrow R$
 - selecting match $m: L \rightarrow G$
 - create new graph H by
 - removing from G the occurrence of $L \setminus R$
 - adding to result a copy of $R \setminus L$



Graph Transformation Step

Introduction
GT theory
Experiments
Conclusion



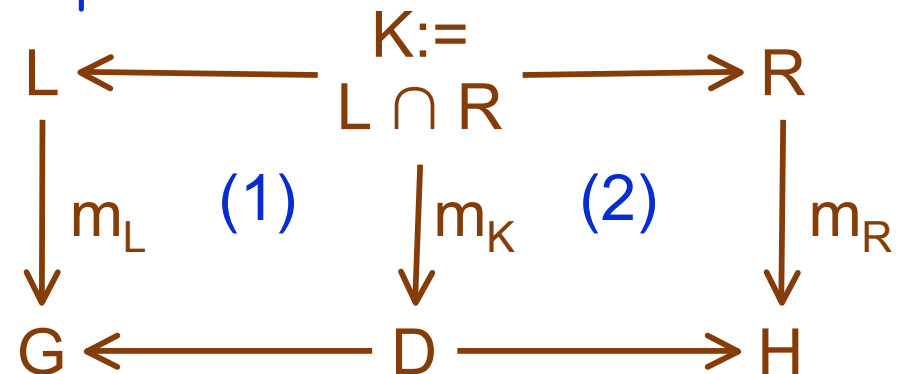
- **Declarative** description

- Set-theoretic

- $G \Rightarrow_{p(m_L)} H$ iff there exists a homomorphism $m: L \cup R \rightarrow G \cup H$ such that $m(L \setminus R) = G \setminus H$ and $m(R \setminus L) = H \setminus G$

- Category-theoretic (DPO):

- $G \Rightarrow_{p(m_L)} H$ iff (1) and (2) are pushouts

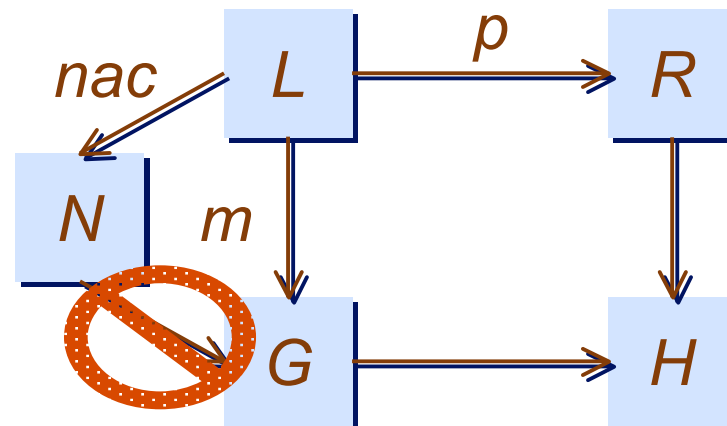


Advanced GT features

- Negative application conditions
- Set nodes (multi objects)
- Attributed graphs and GTs
- Programmed graph transformation
- Graph grammars

Negative application condition

- A negative application condition $nac: L \rightarrow N$ of a graph production $p: L \rightarrow R$ represents a forbidden context. In a graph transformation $t: G \Rightarrow H$, no match $N \rightarrow G$ must be found



Negative application condition

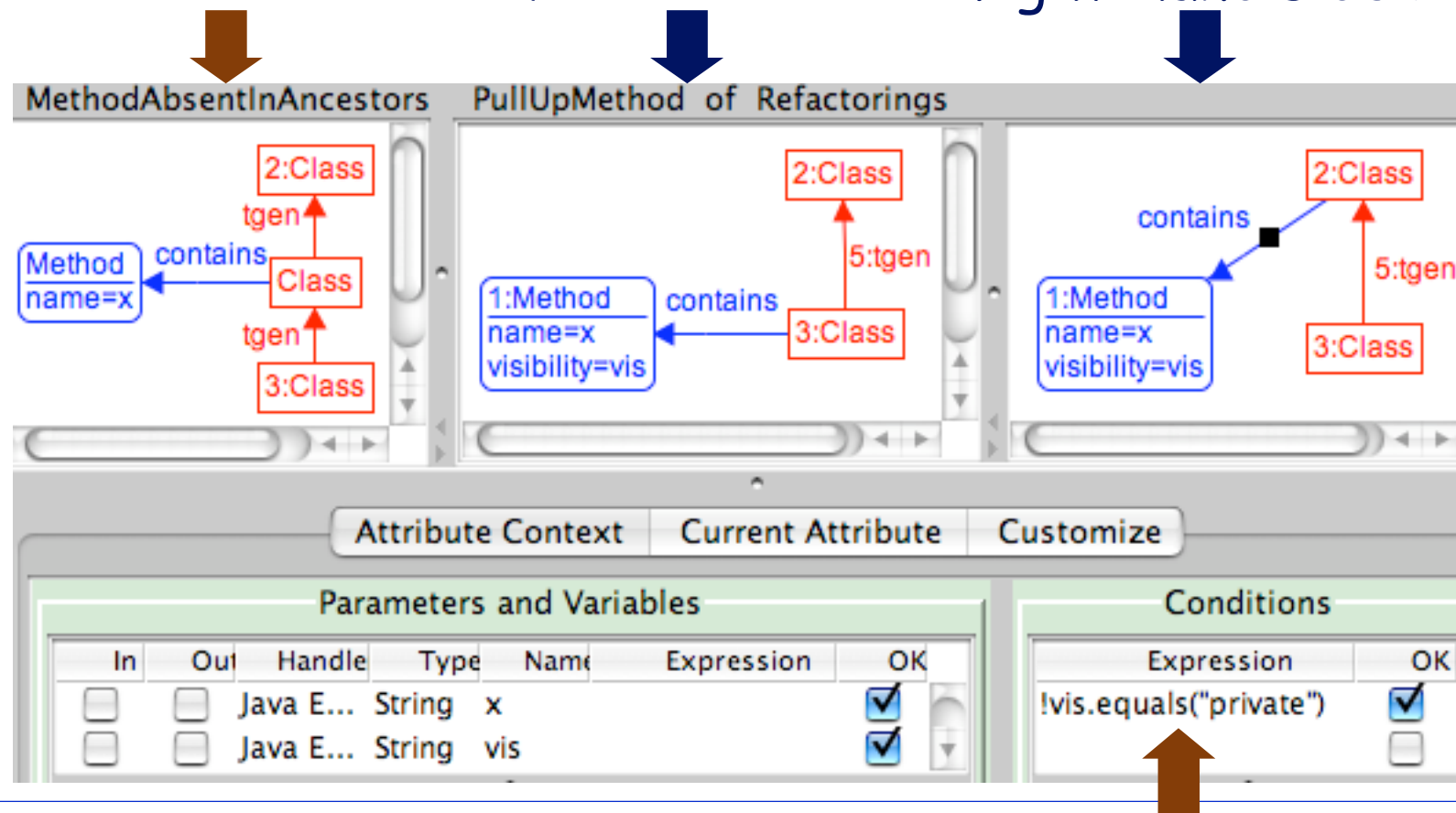
- Example: Pull Up Method refactoring
 - **node attributes** needed
 - to express name and visibility of methods
 - to *constrain* (or *modify*) visibility of methods
 - **negative application conditions (NAC)** needed
 - Method signature of method to be pulled up should be absent in ancestors
 - Method to be pulled up should not be private

Negative application condition

- Example: Pull Up Method refactoring part 1 (P_1)

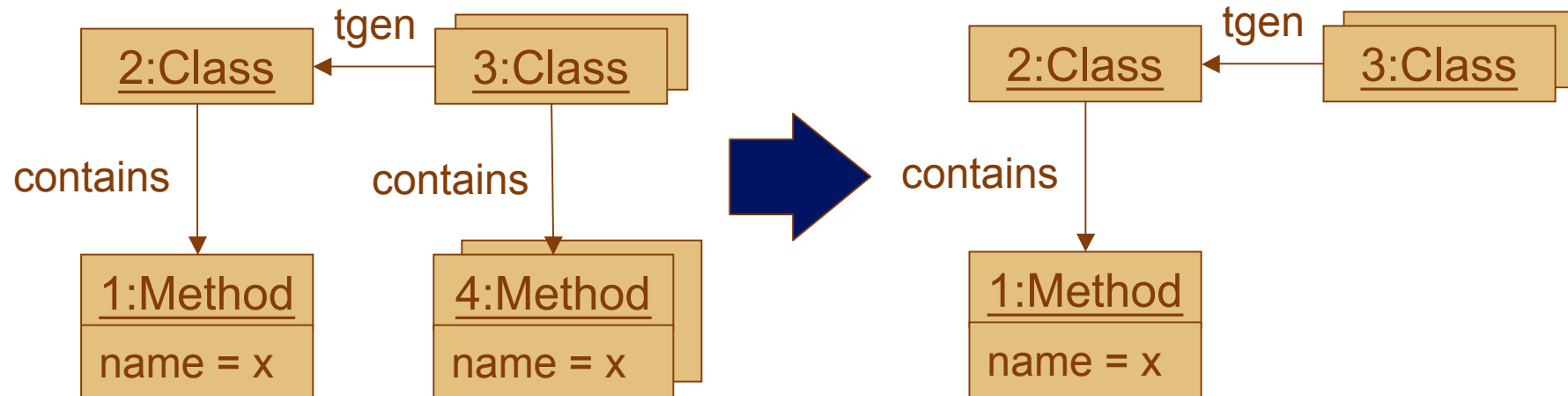
NAC

left-hand side L right-hand side R



Set nodes

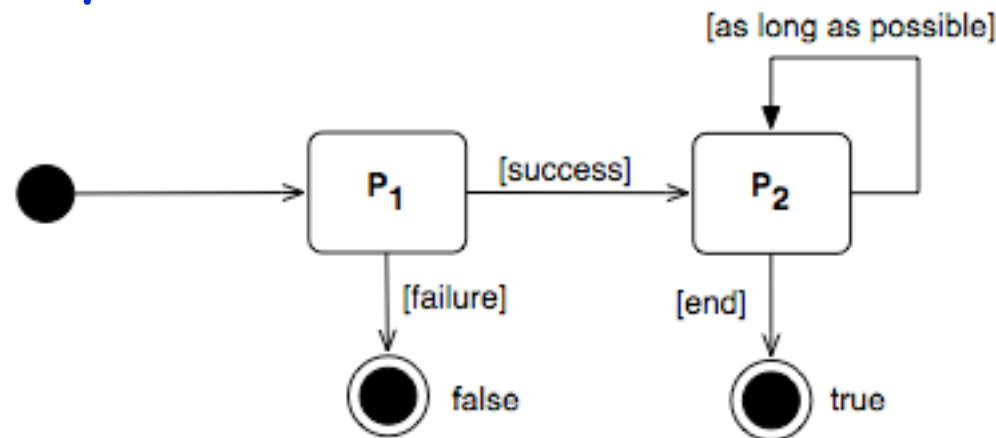
- Example: Pull Up Method refactoring part 2 (P_2)
 - Remove *all* remaining methods with same name in *all* subclasses
 - requires set nodes ...



- ... or programmed GT

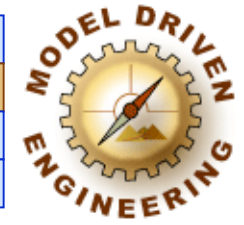
Programmed GT

- For composite graph productions, we need to control their order of application by means of
 - sequencing
 - branching
 - looping
- For example



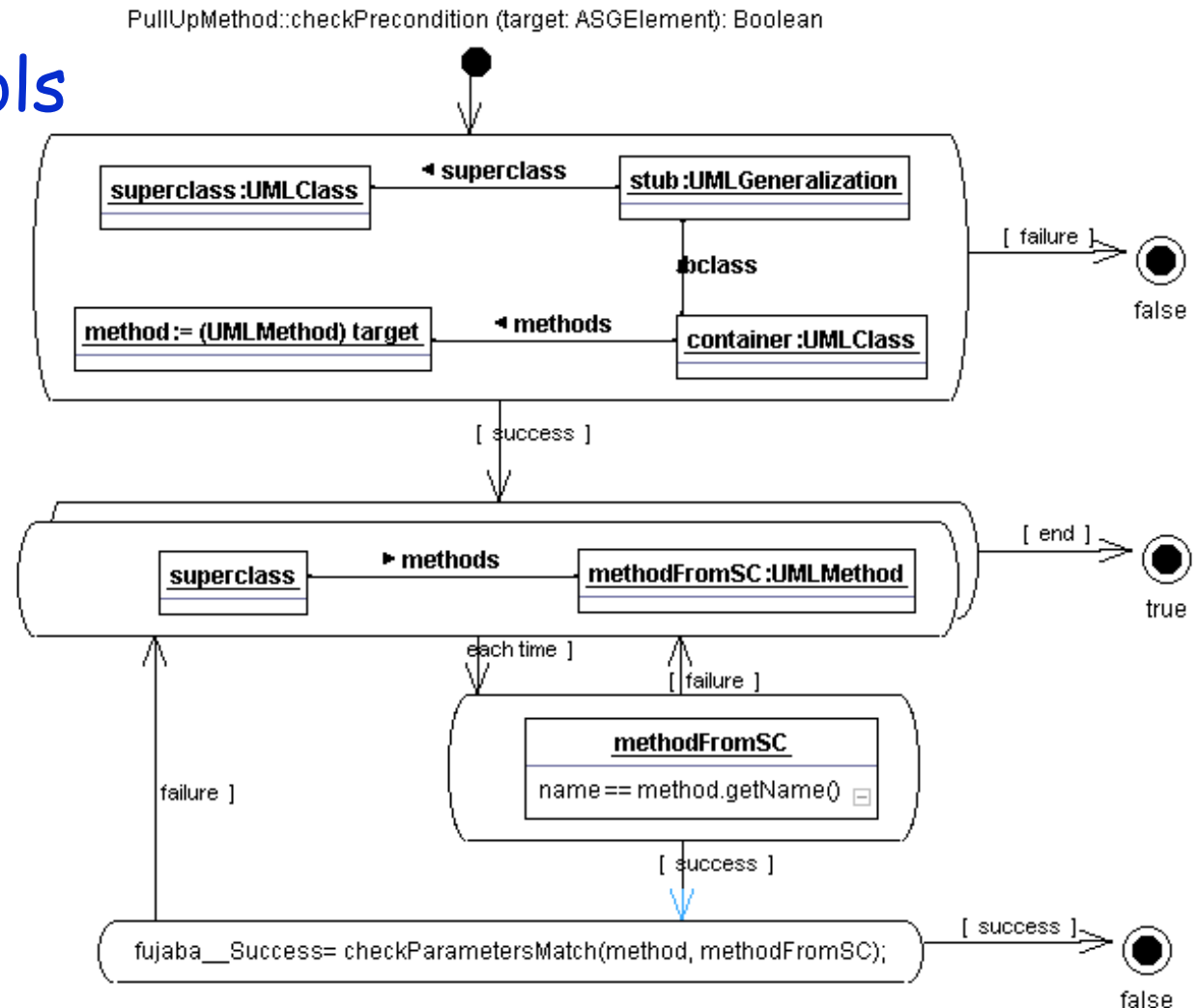
Programmed GT

Introduction
GT theory
Experiments
Conclusion



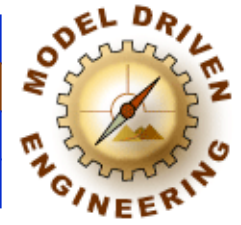
Supported by tools
like Fujaba

using the story
diagram notation



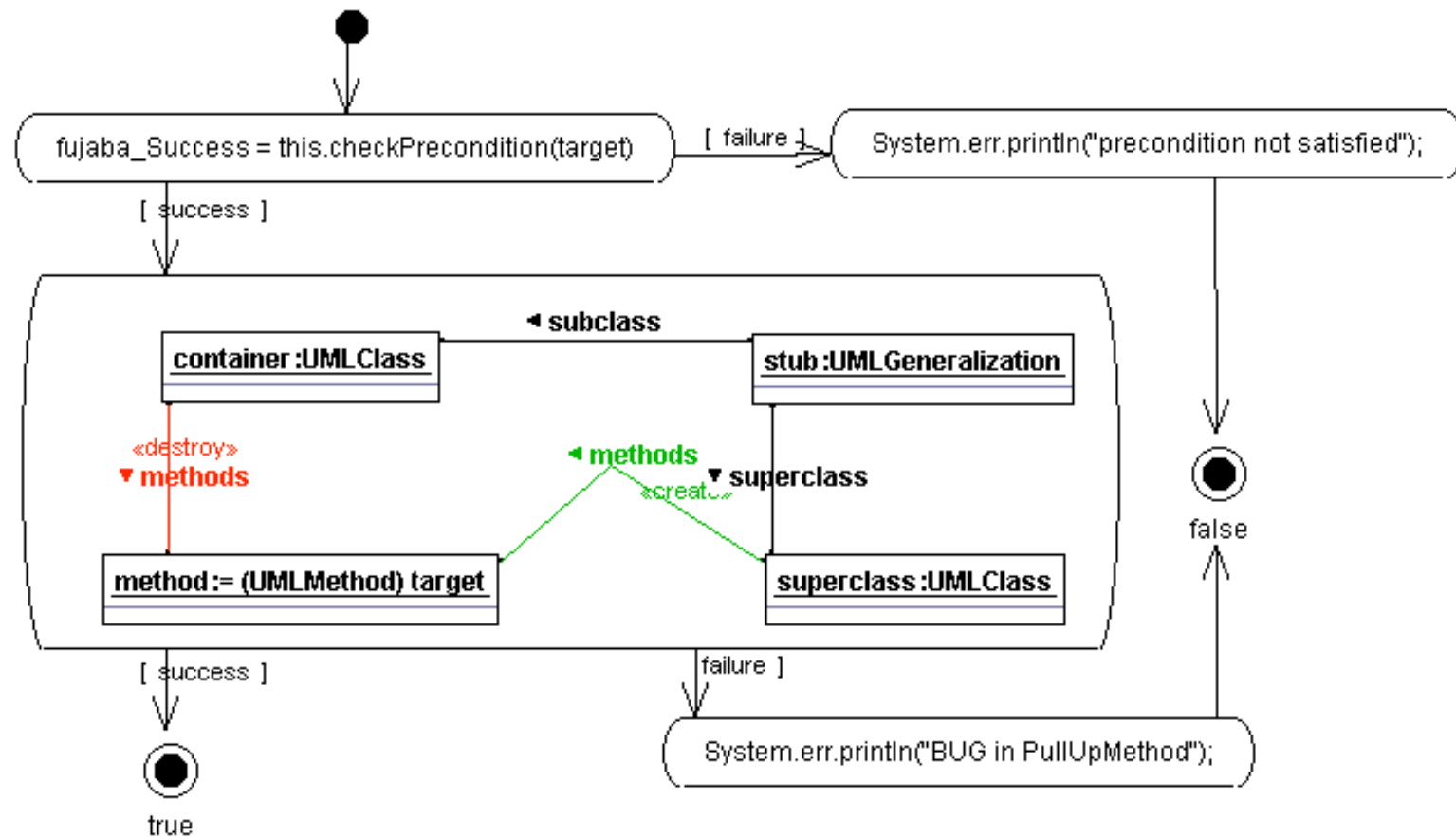
Programmed GT

Introduction
GT theory
Experiments
Conclusion



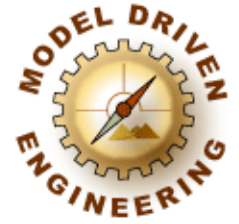
- Supported by tools like Fujaba

PullUpMethod::execute (target: ASGElement): Boolean

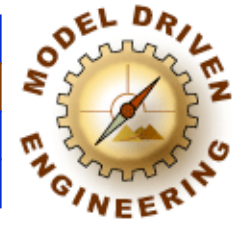


Programmed GT

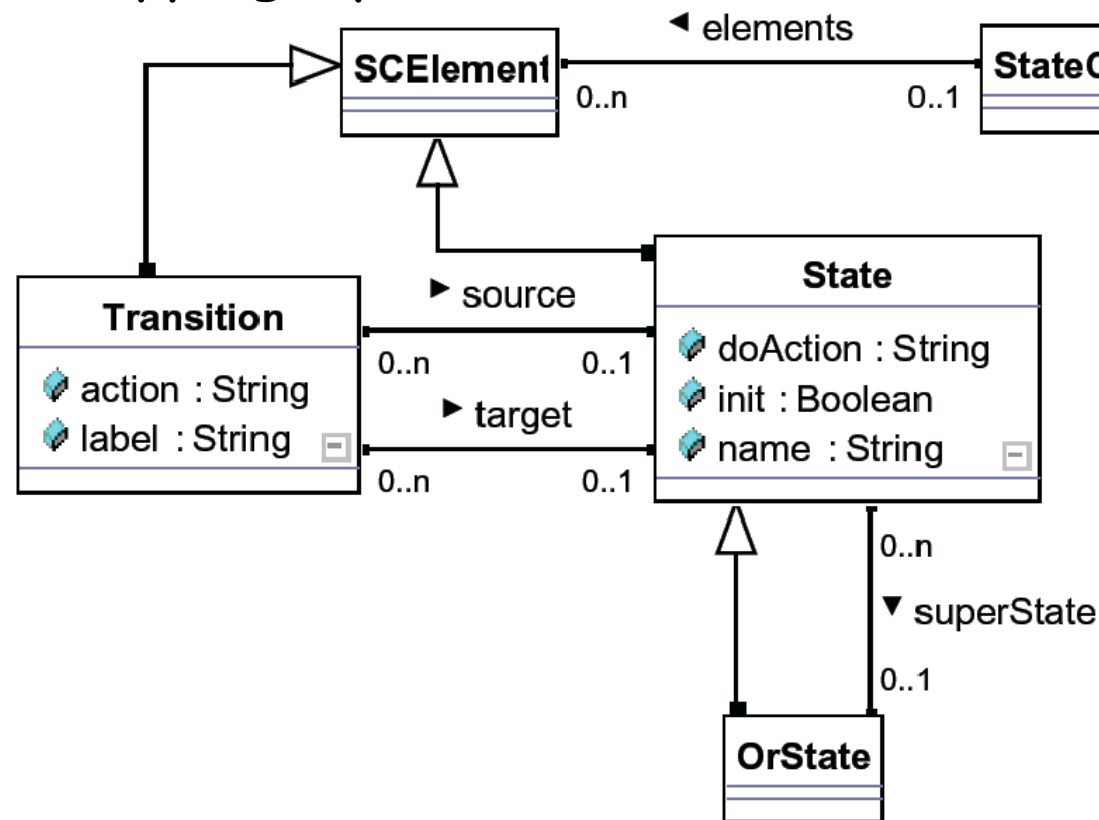
Introduction
GT theory
Experiments
Conclusion



- Supported by tools like *Fujaba*
 - using the story diagram notation
- Example
 - Statechart flattening revisited

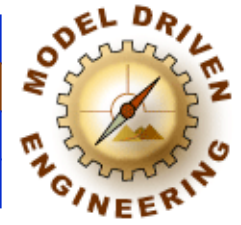


- Example: Statechart flattening revisited
 - Step 1: type graph for statecharts

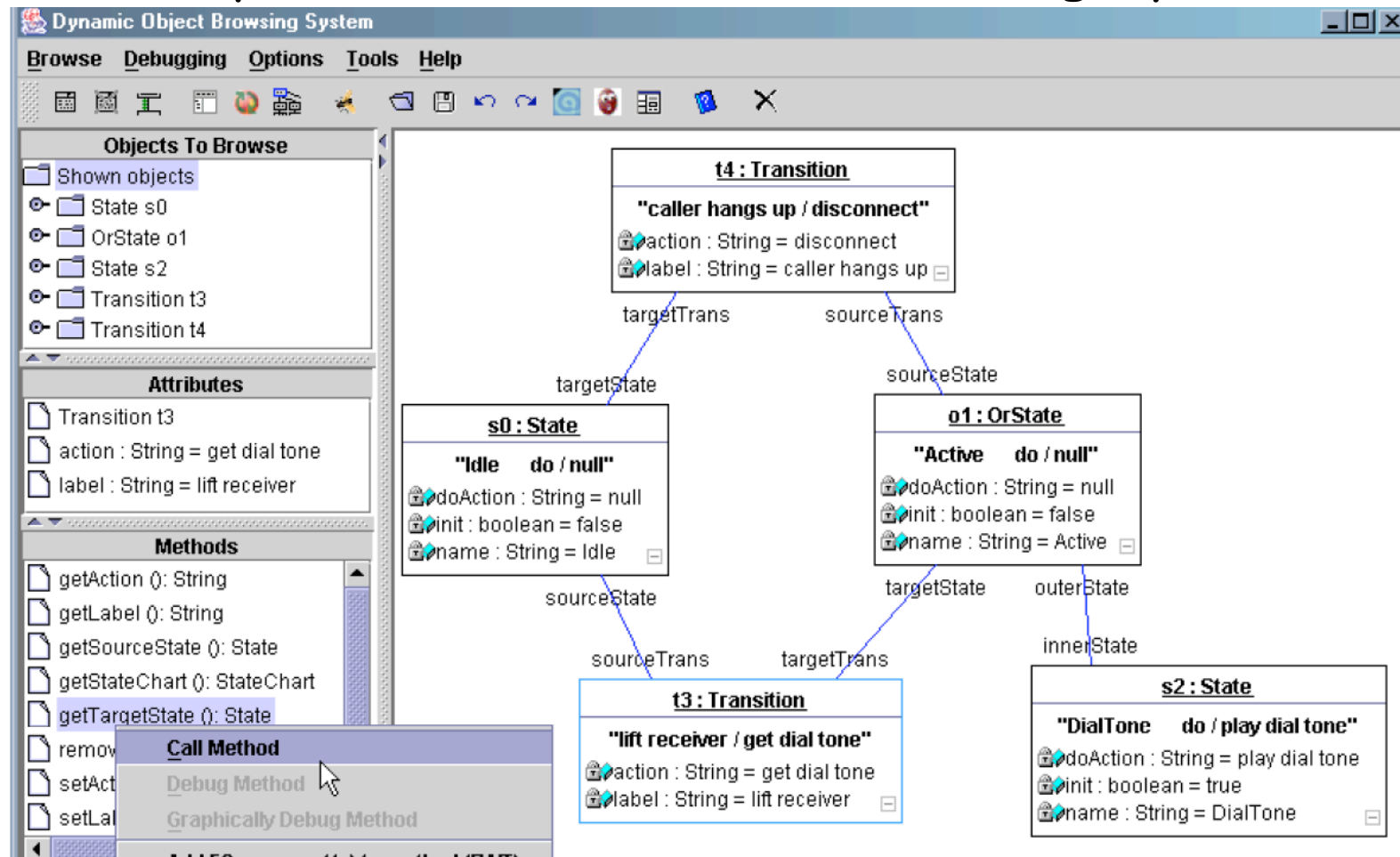


Programmed GT

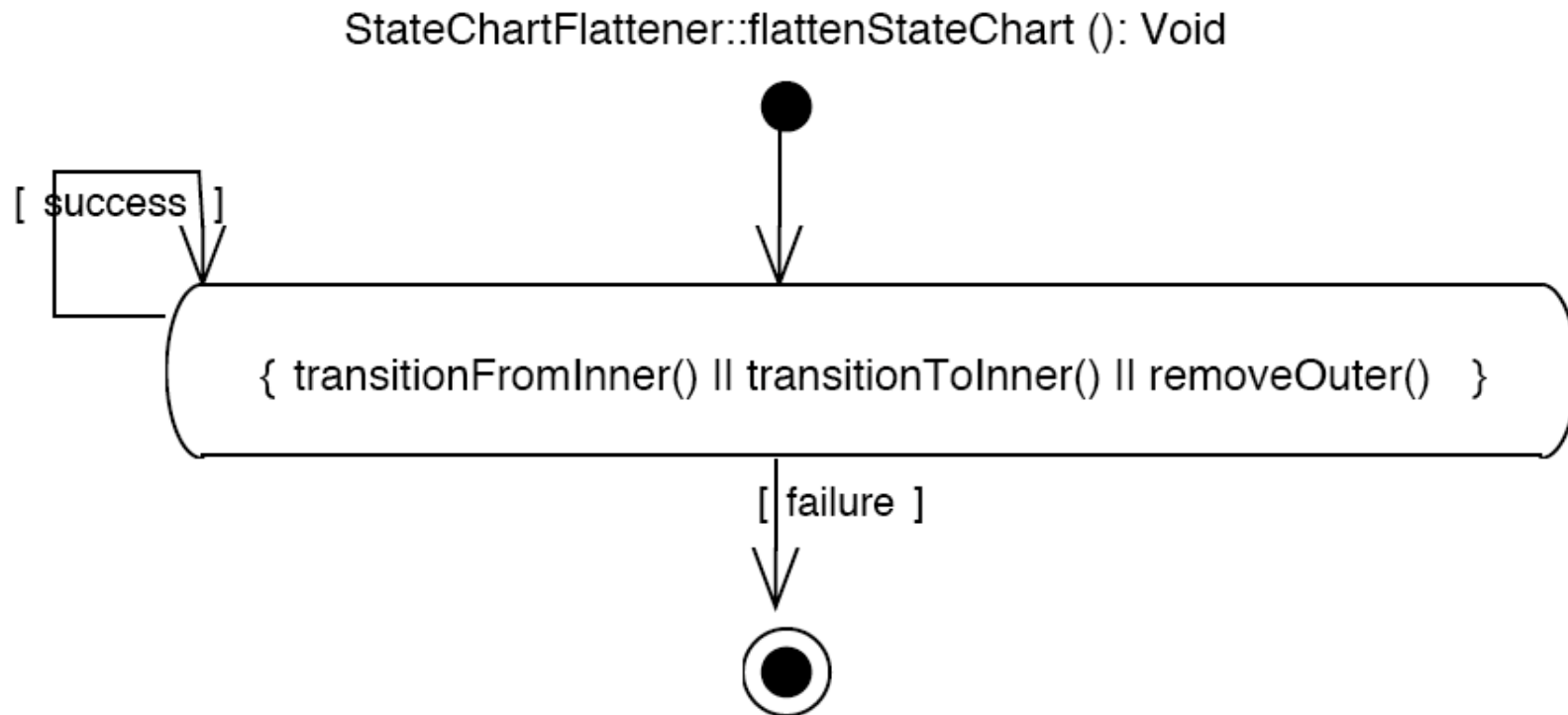
Introduction
GT theory
Experiments
Conclusion



- Example: Statechart flattening revisited
 - Step 2: statecharts as executable graphs

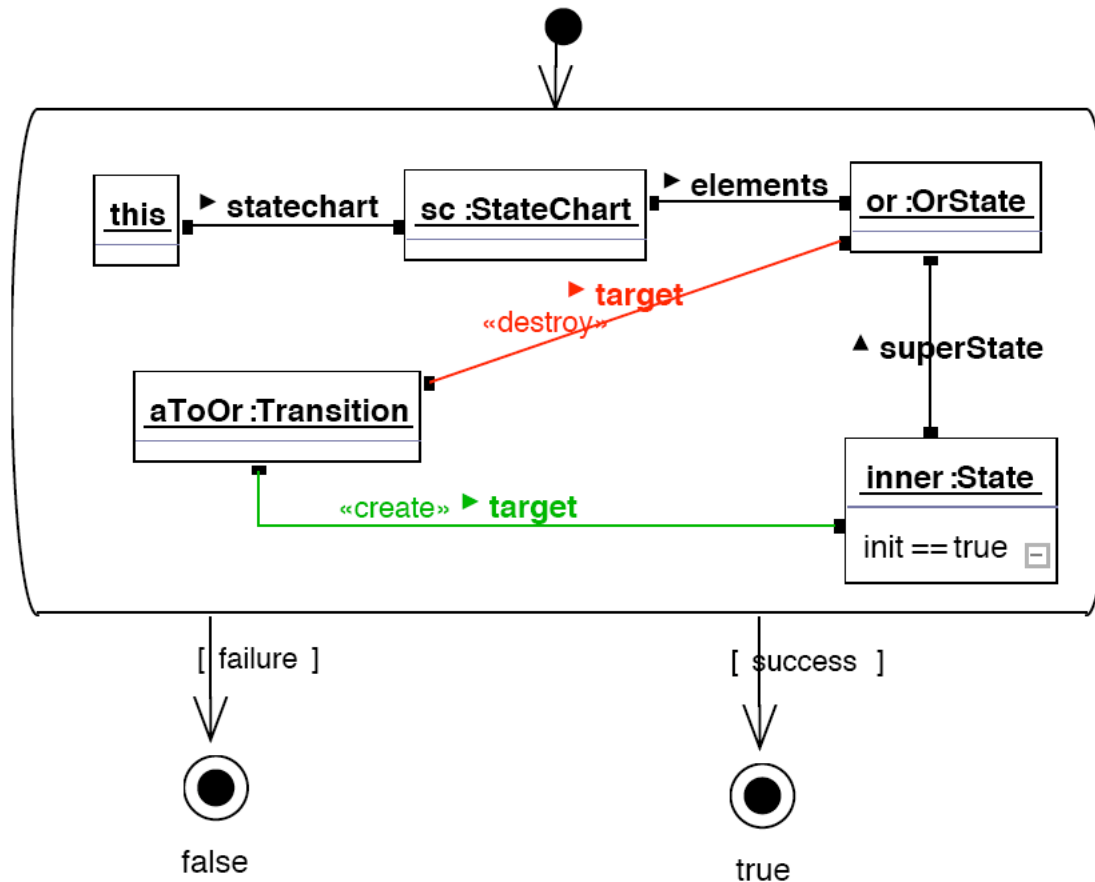


- Example: Statechart flattening revisited
 - Step 3: specifying the statechart flattening



- Example: Statechart flattening revisited
 - Step 3: specifying the statechart flattening

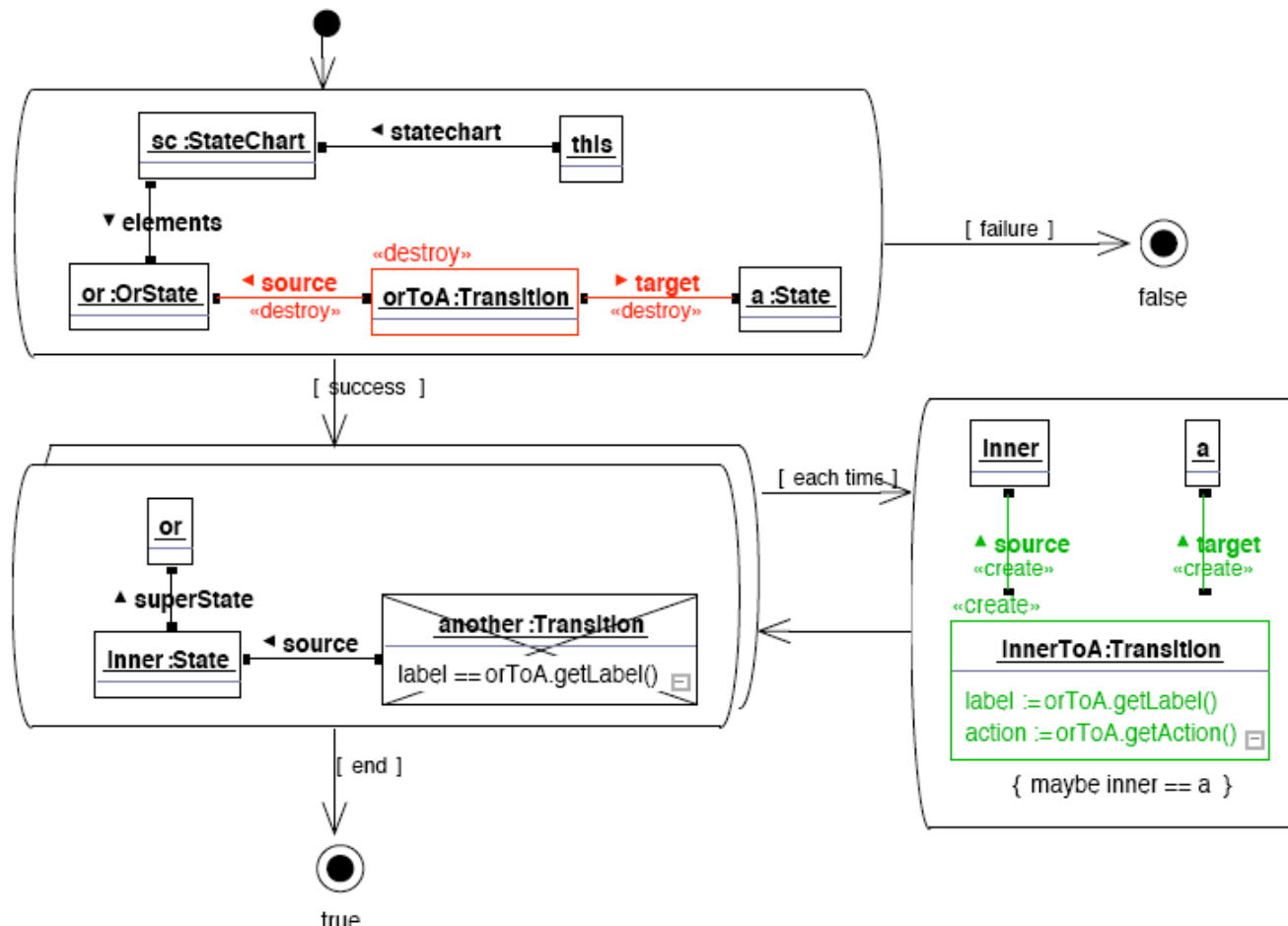
StateChartFlattener::transitionToInner (): Boolean



Programmed GT

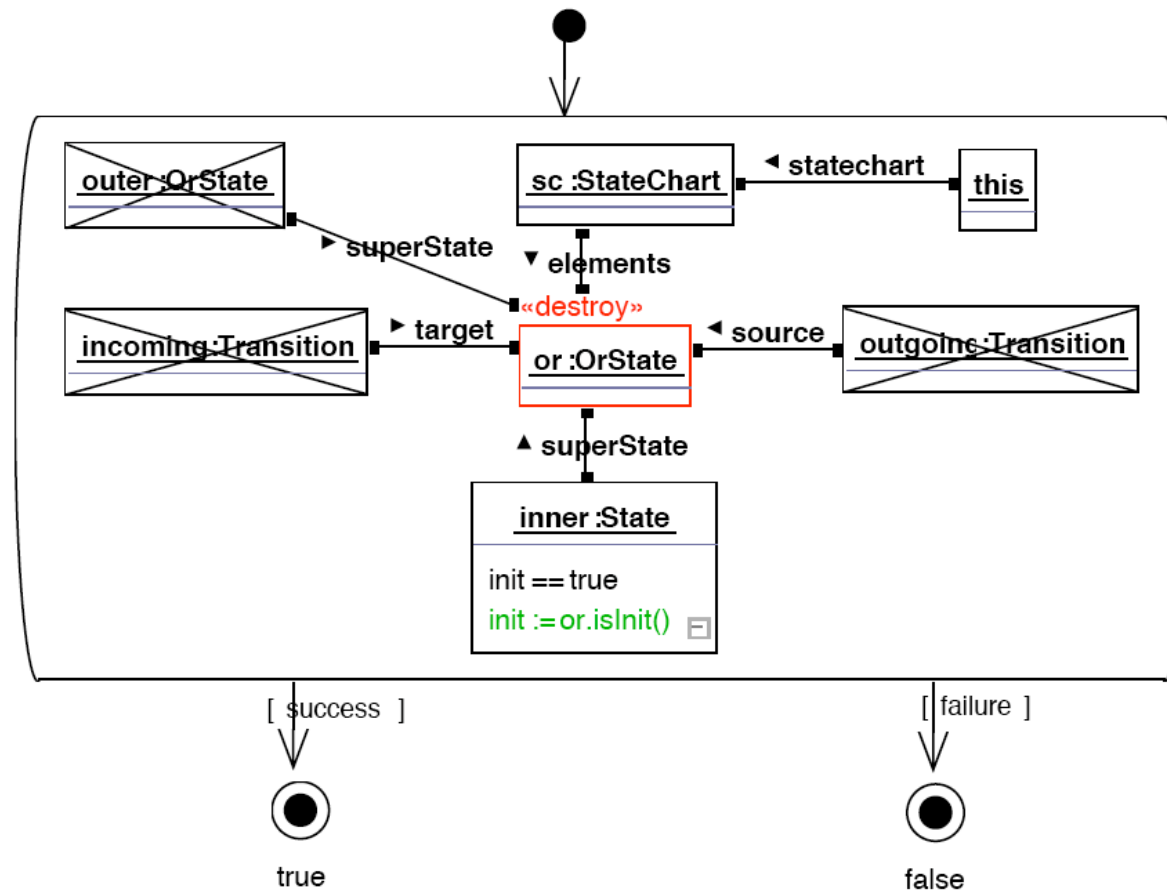
- Example: Statechart flattening revisited
 - Step 3: specifying the statechart flattening

StateChartFlattener::transitionFromInner (): Boolean



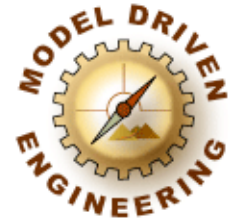
- Example: Statechart flattening revisited
 - Step 3: specifying the statechart flattening

StateChartFlattener::removeOuter (): Boolean



Graph Grammars

Introduction
GT theory
Experiments
Conclusion



- A **graph grammar** is a set of graph productions
- No control structure is imposed on the graph productions to be applied
 - Productions are applied at random, whenever a match is found
- Graph grammars are supported by AGG tool

Where it comes from ...

Chomsky
Grammars



Term
Rewriting



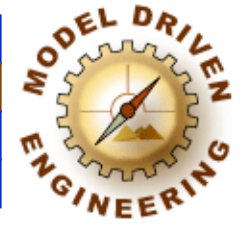
Petri
Nets



Graph Transformation and Graph Grammars

What it is good for ...

Introduction
GT theory
Experiments
Conclusion



Chomsky
Grammars



Term
Rewriting



Petri
Nets



Graph Transformation and Graph Grammars



Diagram
Languages



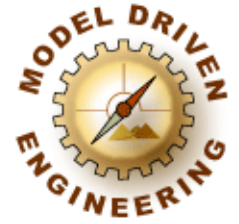
Models of
Computation



Visual
Programming

What it is good for ...

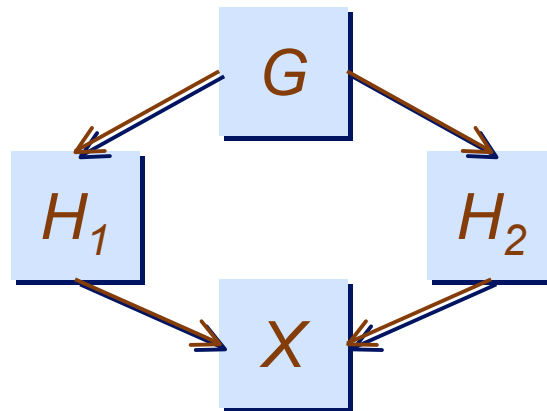
Introduction
GT theory
Experiments
Conclusion



- Behaviour modelling
 - Detecting dependencies and conflicts in functional behaviour
 - Crucial in collaborative/parallel software development

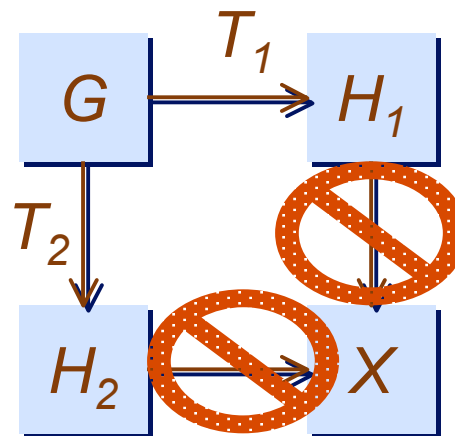
Confluence

- A graph grammar is **confluent** if it has functional behaviour
 - The end result does not depend on the order in which graph productions are applied



Conflicts

- Two graph transformations T_1 and T_2 are in conflict if
 - it is not possible to apply T_1 after T_2 , or T_2 after T_1 , or both, via the same match

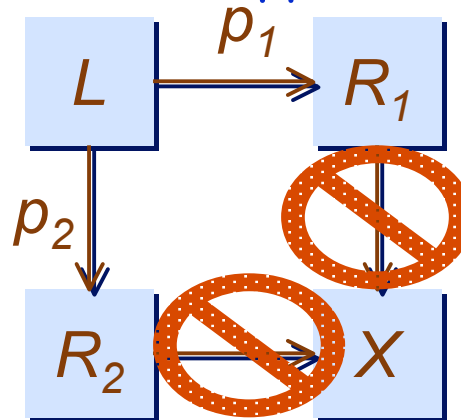


Parallel independence

- Two graph transformations T_1 and T_2 are *parallel independent* if
 - they are not in conflict
- They are *parallel dependent* if they are in conflict
 - In that case they may or may not be sequentialisable

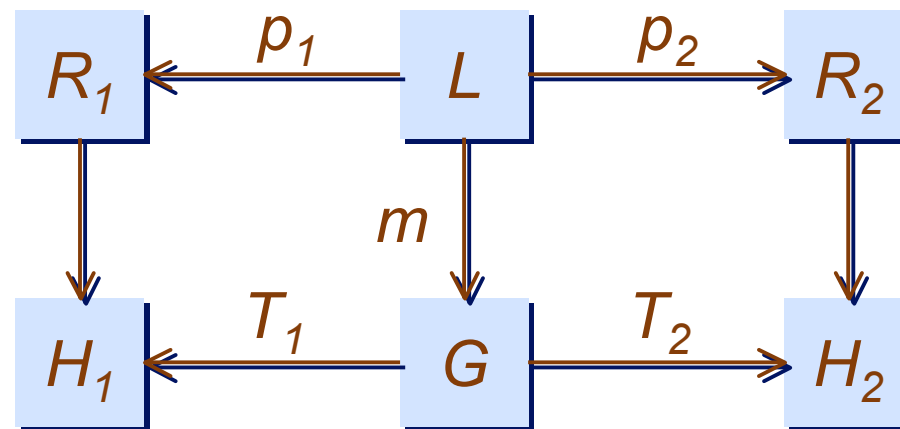
Critical pair analysis

- Needed to detect whether a graph grammar has the confluence property
 - p_1 and p_2 form a *critical pair* if they can both be applied to the same minimal context graph L but applying p_1 prohibits application of p_2 and/or vice versa



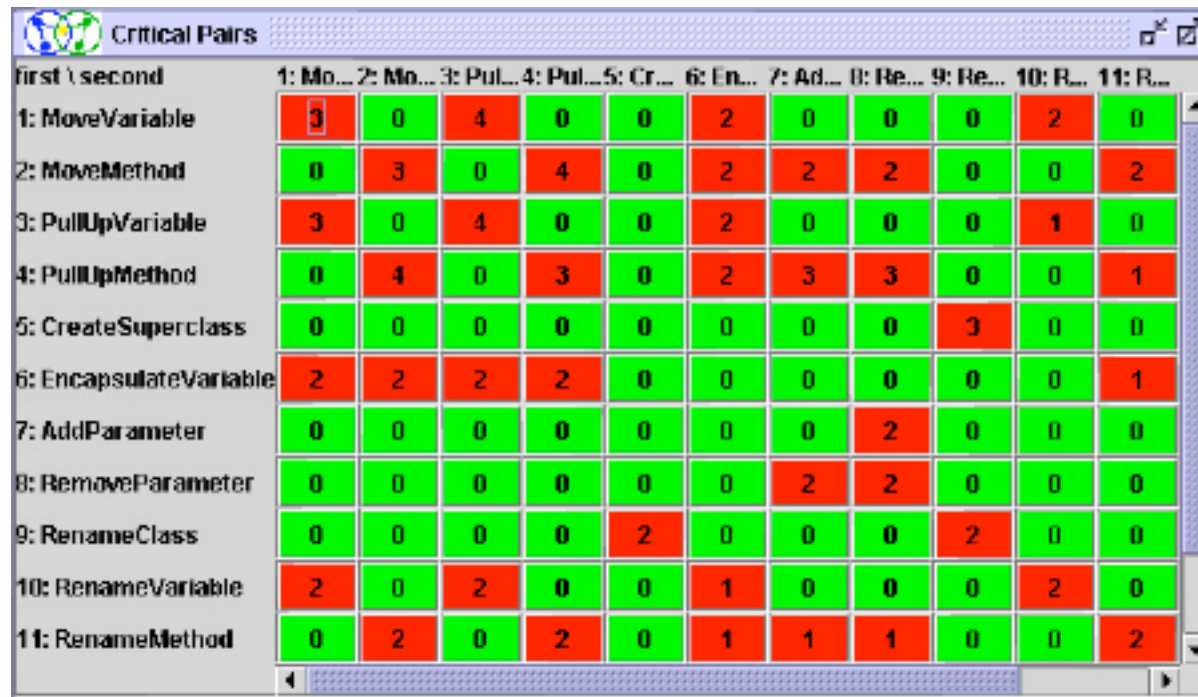
Critical pair analysis

- Conflicts between graph transformations can be found by detecting critical pairs of graph productions
- **Critical pair lemma**
 - For each pair $H_1 \Leftarrow G \Rightarrow H_2$ of graph transformations in conflict, there is a critical pair $R_1 \Leftarrow L \Rightarrow R_2$ expressing the same conflict in a minimal context



Critical pair analysis

- Supported by *AGG* tool
 - potential conflicts can be detected statically
 - e.g. critical pairs between refactoring productions



The screenshot shows the 'Critical Pairs' window of the AGG tool. It displays a matrix of critical pairs between 11 refactoring productions. The rows and columns are labeled as follows:

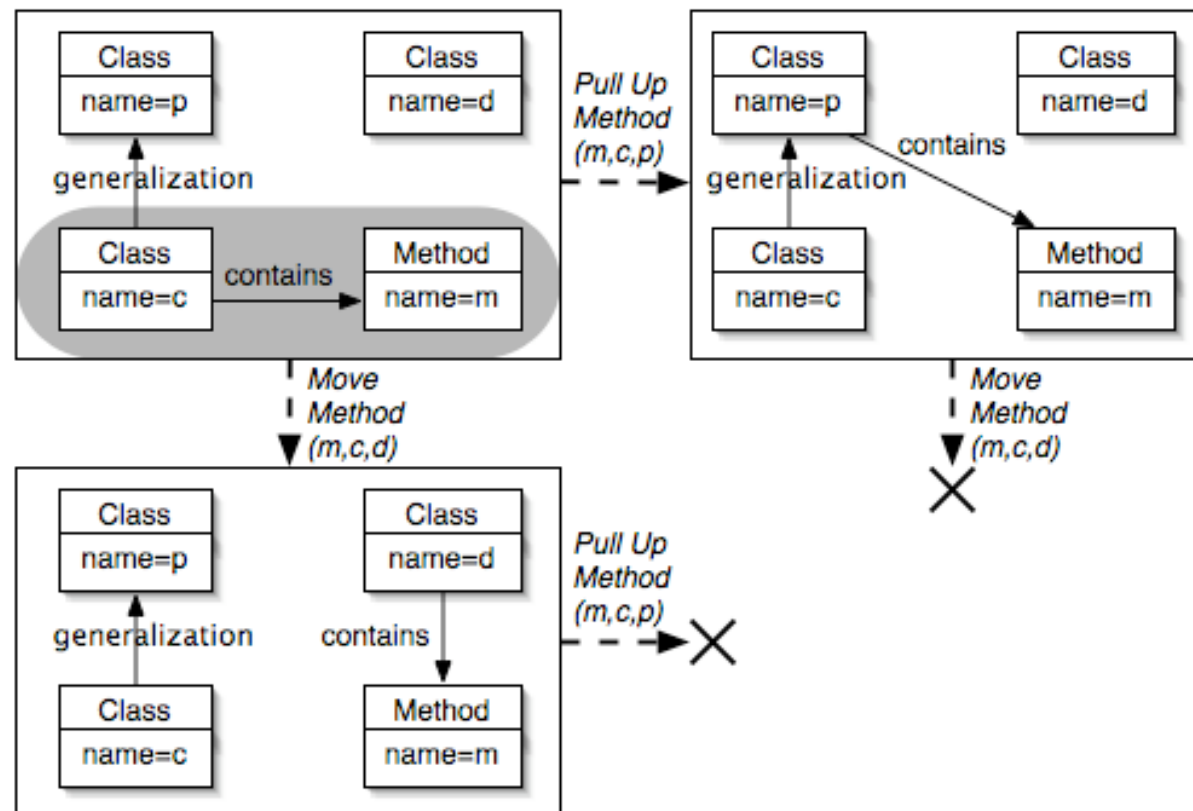
- 1: MoveVariable
- 2: MoveMethod
- 3: PullUpVariable
- 4: PullUpMethod
- 5: CreateSuperclass
- 6: EncapsulateVariable
- 7: AddParameter
- 8: RemoveParameter
- 9: RenameClass
- 10: RenameVariable
- 11: RenameMethod

The matrix cells contain numerical values representing the number of critical pairs. The diagonal elements are all 0. The off-diagonal elements are as follows:

	1: MoveVariable	2: MoveMethod	3: PullUpVariable	4: PullUpMethod	5: CreateSuperclass	6: EncapsulateVariable	7: AddParameter	8: RemoveParameter	9: RenameClass	10: RenameVariable	11: RenameMethod
1: MoveVariable	0	0	0	0	0	0	0	0	0	0	0
2: MoveMethod	0	0	0	0	0	0	0	0	0	0	0
3: PullUpVariable	0	0	0	0	0	0	0	0	0	0	0
4: PullUpMethod	0	0	0	0	0	0	0	0	0	0	0
5: CreateSuperclass	0	0	0	0	0	0	0	0	0	0	0
6: EncapsulateVariable	0	0	0	0	0	0	0	0	0	0	0
7: AddParameter	0	0	0	0	0	0	0	0	0	0	0
8: RemoveParameter	0	0	0	0	0	0	0	0	0	0	0
9: RenameClass	0	0	0	0	0	0	0	0	0	0	0
10: RenameVariable	0	0	0	0	0	0	0	0	0	0	0
11: RenameMethod	0	0	0	0	0	0	0	0	0	0	0

Critical pair analysis

- Concrete example of a critical pair
 - PullUpMethod* versus *MoveMethod*



Sequential dependencies

- Given two graph productions p_1 and p_2 (not necessarily applicable to the same initial graph), are there sequential dependencies between them?
 - Negative dependency
 - p_1 cannot be applied after p_2
 - p_1 violates the truth of p_2 's precondition
 - Corresponds to a critical pair
 - Positive dependency
 - p_1 can only be applied after p_2
 - p_1 enables the truth of part of p_2 's precondition
 - Not (yet) supported by AGG tool