Introduction to process algebra

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Actions & processes

Action

- elementary unit of behaviour that can execute itself atomically in time (no duration), after which it terminates successfully
- is a latency for interaction

$$\alpha \quad ::= \quad \tau \mid a \mid \alpha \mid \alpha$$

- $a \mid b \mid \cdots \mid z$ represent a collection of actions that occur at the same time instant
- τ is the empty action, which contains no actions and as such cannot be observed
- $\langle N, |, \tau \rangle$ forms a monoid

Actions & processes

Process

is a description of how the interaction capacities of a system evolve, *i.e.*, its behaviour
for example,

$$E \triangleq a.b + a.E$$

- analogy: regular expressions vs finite automata

The framework

Process

... abstract representation of a system's behaviour

Algebra

... a mathematical structure satisfying a particular set of axioms

Process Algebra

... a framework for the specification and manipulation of process terms as induced by a collection of operator symbols, encompassing an operational and an axiomatic theory

The framework

Transition systems : operational representation of system's behaviour through labelled graphs

Behavioural equivalences : to distinguished states in transition systems

Process terms : algebraic representation of transition systems (for the purpose of mathematical reasoning)

Structural operational semantics : inductive proof rules to provide each process term with its intended transition system

Equational theory Axiomatic theory of processes, expressed in an equational logic on process terms, that is sound and complete wrt bisimilarity.

Instantiating the framework

CCS: a prototypical process algebra

- *Calculus of Communicating Systems* [Milner, 1980]
- Actions:

$$Act \quad ::= \quad a \mid \overline{a} \mid \tau$$

  for $a \in N$, N denoting a set of names
- Processes:
  - No sequential composition: but action prefix $a.$
  - No distinction between termination and deadlock (why?)
  - Communication by binary handshake
    (of complementary actions)

Examples

Buffers

1-position buffer: $A(in, out) \triangleq in.\overline{out}.\mathbf{0}$

... non terminating: $B(in, out) \triangleq in.\overline{out}.B$

... with two output ports: $C(in, o_1, o_2) \triangleq in.(\overline{o_1}.C + \overline{o_2}.C)$

... non deterministic: $D(in, o_1, o_2) \triangleq in.\overline{o_1}.D + in.\overline{o_2}.D$

... with parameters: $B(in, out) \triangleq in(x).\overline{out}\langle x \rangle.B$

Examples

n-position buffers

1-position buffer:

$$S \triangleq \text{new}\,\{m\}\,(B\langle in, m\rangle \mid B\langle m, out\rangle)$$

n-position buffer:

$$Bn \triangleq \text{new}\,\{m_i | i < n\}\,(B\langle in, m_1\rangle \mid B\langle m_1, m_2\rangle \mid \cdots \mid B\langle m_{n-1}, out\rangle)$$

Examples

mutual exclusion

$$Sem \triangleq get.put.Sem$$

$$P_i \triangleq \overline{get}.c_i.\overline{put}.P_i$$

$$S \triangleq \mathrm{new}\,\{get, put\}\ (Sem \mid (|_{i \in I}\ P_i))$$

CCS Syntax

The set $\mathbb{P}$ of processes is the set of all terms generated by the following BNF:

$$E \ ::= \ A(x_1, \dots, x_n) \ | \ a.E \ | \ \sum_{i \in I} E_i \ | \ E_0 \,|\, E_1 \ | \ \text{new}\, K \ E$$

for $a \in Act$ and $K \subseteq L$

Abbreviatures

$$E_0 + E_1 \stackrel{\text{abv}}{=} \sum_{i \in \{0,1\}} E_i$$

$$\mathbf{0} \stackrel{\text{abv}}{=} \sum_{i \in \emptyset} E_i$$

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CCS Syntax

Process declaration

$$A(\tilde{x}) \triangleq E_A$$

with $\text{fn}(E_A) \subseteq \tilde{x}$ (where $\text{fn}(P)$ is the set of free variables of P).

- used as, *e.g.*, $A(a, b, c) \triangleq a.b.\mathbf{0} + c.A\langle d, e, f \rangle$

Process declaration: fixed point expression

$$\underline{fix}\,(X = E_X)$$

- syntactic substitution over $\mathbb{P}$, *cf.*,
  - $\{c/b\}\, a.b.\mathbf{0}$
  - (internal variables renaming) $\{x/y\}\,\text{new}\,\{x\}\; y.x.\mathbf{0} = \text{new}\,\{x'\}\; x.x'.\mathbf{0}$

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Sort

A sort of a process P is an interface for P

- minimal sort: $\mathcal{L}(P) = \bigcap\{K \subseteq L \mid K \text{ is a sort of } P\}$
- syntactic sort, *i.e.*, the set of free variables:

$$\begin{aligned}
\text{fn}(a.P) &= \{a\} \cup \text{fn}(P) \\
\text{fn}(\tau.P) &= \text{fn}(P) \\
\text{fn}(\sum_{i \in I} P_i) &= \bigcup_{i \in I} \text{fn}(P_i) \\
\text{fn}(P \mid Q) &= \text{fn}(P) \cup \text{fn}(Q) \\
\text{fn}(\text{new}\, K\ P) &= \text{fn}(P) - (K \cup \overline{K})
\end{aligned}$$

and, for each $P(\tilde{x}) \triangleq E$, $\text{fn}(E) \subseteq \text{fn}(P(\tilde{x})) = \tilde{x}$.

Sort

The minimal sort $\mathcal{L}(P)$ corresponds to the set of actions which effectively label valid transitions of P. All the other sorts are upper bounds of $\mathcal{L}(P)$ and represent interaction possibilities.

Warning

- $\mathsf{new}\,\{a\}\ (a.b.c.0)$ has no transitions, so its sort is $\emptyset$
- however: $\mathsf{fn}((\mathsf{new}\,\{a\}\ a.b.c.0)) = \{b, c\}$

Semantics

Two-level semantics

- arquitectural, expresses a notion of similar assembly configurations and is expressed through a structural congruence relation;
- behavioural given by transition rules which express how system's components interact

Semantics

Structural congruence

$\equiv$ over $\mathbb{P}$ is given by the closure of the following conditions:

- for all $A(\tilde{x}) \triangleq E_A$, $A(\tilde{y}) \equiv \{\tilde{x}/\tilde{y}\}\, E_A$,
  (*i.e.*, folding/unfolding preserve $\equiv$)
- α-conversion (*i.e.*, replacement of bounded variables).
- both $|$ and $+$ originate, with $\mathbf{0}$, abelian monoids
- forall $a \notin \mathrm{fn}(P)$ new $\{a\}\ (P \mid Q) \equiv P \mid$ new $\{a\}\ Q$
- new $\{a\}\ \mathbf{0} \equiv \mathbf{0}$

Semantics

$$\frac{}{a.p \xrightarrow{a} p}\ (prefix)$$

$$\frac{\{\tilde{k}/\tilde{x}\}\, p_A \xrightarrow{a} p'}{A(\tilde{k}) \xrightarrow{a} p'}\ (ident) \quad (\text{if } A(\tilde{x}) \triangleq p_A)$$

$$\frac{p \xrightarrow{a} p'}{p + q \xrightarrow{a} p'}\ (sum - l) \qquad \frac{q \xrightarrow{a} q'}{p + q \xrightarrow{a} q'}\ (sum - r)$$

Semantics

$$\frac{p \xrightarrow{a} p'}{p \mid q \xrightarrow{a} p' \mid q} \ (par - l) \qquad \frac{q \xrightarrow{a} q'}{p \mid q \xrightarrow{a} p \mid q'} \ (par - r)$$

$$\frac{p \xrightarrow{a} p' \quad q \xrightarrow{\overline{a}} q'}{p \mid q \xrightarrow{\tau} p' \mid q'} \ (react)$$

$$\frac{p \xrightarrow{a} p'}{\text{new } \{k\}\ p \xrightarrow{a} \text{new } \{k\}\ p} \ (res) \quad (\text{if } a \notin \{k, \overline{k}\})$$

Compatibility

Lemma

Structural congruence preserves transitions:

if $p \xrightarrow{a} p'$ and $p \equiv q$ there exists a process q' such that $q \xrightarrow{a} q'$ and $p' \equiv q'$.

Semantics

These rules define a LTS

$$\{\xrightarrow{a} \subseteq \mathbb{P} \times \mathbb{P} \mid a \in Act\}$$

Relation $\xrightarrow{a}$ is defined inductively over process structure entailing a semantic description which is

Structural *i.e.*, each process shape (defined by the most external combinator) has a type of transitions

Modular *i.e.*, a process trasition is defined from transitions in its sup-processes

Complete *i.e.*, all possible transitions are infered from these rules

static vs dynamic combinators

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Graphical representations

Synchronization diagram

- represent interfaces of processes
- static combinators are an algebra of synchronization diagrams

Transition graph

- derivative, n-derivative, transition tree
- folds into a transition graph

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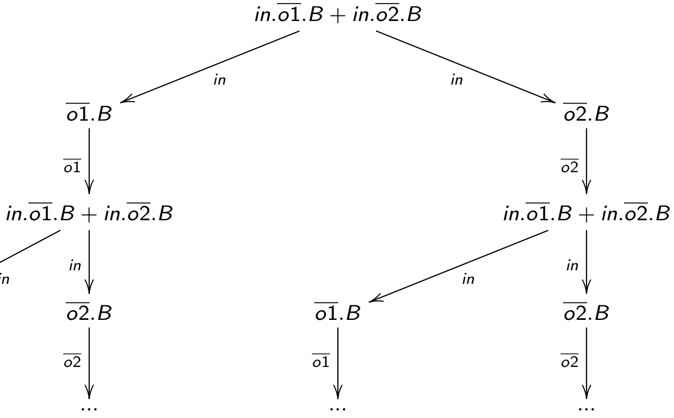
Transition tree

$B \triangleq in.\overline{o1}.B + in.\overline{o2}.B$



Transition graph

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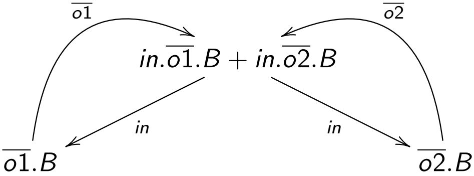

compare with $B' \triangleq in.(\overline{o1}.B' + \overline{o2}.B')$

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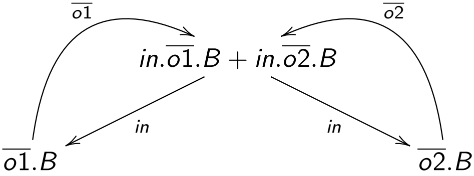

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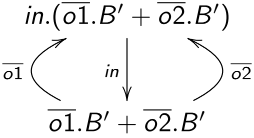

Data parameters

Language $\mathbb{P}$ is extended to $\mathbb{P}_V$ over a data universe V, a set V_e of expressions over V and a evaluation $Val : V_e \rightarrow V$

Example

$$B \triangleq in(x).B'_x$$
$$B'_v \triangleq \overline{out}\langle v\rangle.B$$

- Two prefix forms: $a(x).E$ and $\overline{a}\langle e\rangle.E$ (actions as ports)
- Data parameters: $A_S(x_1, ..., x_n) \triangleq E_A$, with $S \in V$ and each $x_i \in L$
- Conditional combinator: if b then P, if b then P_1 else P_2

Clearly

$$\text{if } b \text{ then } P_1 \text{ else } P_2 \stackrel{\text{abv}}{=} (\text{if } b \text{ then } P_1) + (\text{if } \neg b \text{ then } P_2)$$

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Data parameters

Additional semantic rules

$$\frac{}{a(x).E \xrightarrow{a(v)} \{v/x\}E}\ (prefix_i) \qquad \text{for } v \in V$$

$$\frac{}{\overline{a}\langle e\rangle.E \xrightarrow{\overline{a}\langle v\rangle} E}\ (prefix_o) \qquad \text{for } Val(e) = v$$

$$\frac{E_1 \xrightarrow{a} E'}{\text{if } b \text{ then } E_1 \text{ else } E_2 \xrightarrow{a} E'}\ (if_1) \qquad \text{for } Val(b) = \text{true}$$

$$\frac{E_2 \xrightarrow{a} E'}{\text{if } b \text{ then } E_1 \text{ else } E_2 \xrightarrow{a} E'}\ (if_2) \qquad \text{for } Val(b) = \text{false}$$

Back to *PP*

Encoding in the basic language: $\mathcal{T}(\ )\mathbb{P}_V\mathbb{P}$

$$\mathcal{T}(a(x).E) = \sum_{v\in V} a_v.\mathcal{T}(\{v/x\}E)$$
$$\mathcal{T}(\overline{a}\langle e\rangle.E) = \overline{a}_e.\mathcal{T}(E)$$
$$\mathcal{T}(\sum_{i\in I} E_i) = \sum_{i\in I} \mathcal{T}(E_i)$$
$$\mathcal{T}(E \mid F) = \mathcal{T}(E) \mid \mathcal{T}(F)$$
$$\mathcal{T}(\text{new}\, K\;\, E) = \text{new}\,\{a_v \mid a\in K, v\in V\}\;\mathcal{T}(E)$$

and

$$\mathcal{T}(\text{if}\, b\, \text{then}\, E) = \begin{cases} \mathcal{T}(E) & \text{if } Val(b) = \text{true} \\ \mathbf{0} & \text{if } Val(b) = \text{false} \end{cases}$$

EX1: Canonical concurrent form

$$P \triangleq \text{new } K \ (E_1 \mid E_2 \mid \ldots \mid E_n)$$

The chance machine

$$IO \triangleq m.\overline{bank}.(lost.\overline{loss}.IO + rel(x).\overline{win}\langle x\rangle.IO)$$

$$B_n \triangleq bank.\overline{max}\langle n+1\rangle.left(x).B_x$$

$$Dc \triangleq max(z).(\overline{lost}.\overline{left}\langle z\rangle.Dc + \sum_{1\leq x\leq z} \overline{rel}\langle x\rangle.\overline{left}\langle z-x\rangle.Dc)$$

$$M_n \triangleq \text{new } \{bank, max, left, lost, rel\} \ (IO \mid B_n \mid Dc)$$

EX2: Sequential patterns

1. List all states (configurations of variable assignments)
2. Define an order to capture systems's evolution
3. Specify an expression in $\mathbb{P}$ to define it

A 3-bit converter

$$
\begin{aligned}
A &\triangleq rq.B \\
B &\triangleq out0.C + out1.\overline{odd}.A \\
C &\triangleq out0.D + out1.\overline{even}.A \\
D &\triangleq out0.\overline{zero}.A + out1.\overline{even}.A
\end{aligned}
$$

Processes are 'prototypical' transition systems

... hence all definitions apply:

$E \sim F$

- Processes E, F are bisimilar if there exist a bisimulation S st $\{\langle E, F\rangle\} \in S$.
- A binary relation S in $\mathbb{P}$ is a (strict) bisimulation iff, whenever $(E, F) \in S$ and $a \in Act$,

$$\begin{array}{ll} \text{i)} & E \xrightarrow{a} E' \;\Rightarrow\; F \xrightarrow{a} F' \wedge (E', F') \in S \\ \text{ii)} & F \xrightarrow{a} F' \;\Rightarrow\; E \xrightarrow{a} E' \wedge (E', F') \in S \end{array}$$

I.e.,

$$\sim \;=\; \bigcup\{S \subseteq \mathbb{P} \times \mathbb{P} \mid S \quad \text{is a (strict) bisimulation}\}$$

Processes are 'prototipycal' transition systems

Example: $S \sim M$

$$T \triangleq i.\overline{k}.T$$
$$R \triangleq k.j.R$$
$$S \triangleq \text{new}\,\{k\}\,(T \mid R)$$

$$M \triangleq i.\tau.N$$
$$N \triangleq j.i.\tau.N + i.j.\tau.N$$

through bisimulation

$$R = \{\langle S, M\rangle\rangle, \langle \text{new}\,\{k\}\,(\overline{k}.T \mid R), \tau.N\rangle, \langle \text{new}\,\{k\}\,(T \mid j.R), N\rangle, \langle \text{new}\,\{k\}\,(\overline{k}.T \mid j.R), j.\tau.N\rangle\}$$

Example: Semaphores

A semaphore

$$Sem \triangleq get.put.Sem$$

n-semaphores

$$\begin{aligned} Sem_n &\triangleq Sem_{n,0} \\ Sem_{n,0} &\triangleq get.Sem_{n,1} \\ Sem_{n,i} &\triangleq get.Sem_{n,i+1} + put.Sem_{n,i-1} \\ &\quad (\text{for } 0 < i < n) \\ Sem_{n,n} &\triangleq put.Sem_{n,n-1} \end{aligned}$$

Sem_n can also be implemented by the parallel composition of n Sem processes:

$$Sem^n \triangleq Sem \mid Sem \mid \ldots \mid Sem$$

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Example: Semaphores

Is $Sem_n \sim Sem^n$?

For $n = 2$:

$$\{\langle Sem_{2,0}, Sem \mid Sem\rangle, \langle Sem_{2,1}, Sem \mid put.Sem\rangle,$$
$$\langle Sem_{2,1}, put.Sem \mid Sem\rangle \langle Sem_{2,2}, put.Sem \mid put.Sem\rangle\}$$

is a bisimulation.

- but can we get rid of structurally congruent pairs?

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Bisimulation up to $\equiv$

Definition

A binary relation S in $\mathbb{P}$ is a (strict) bisimulation up to $\equiv$ iff, whenever $(E, F) \in S$ and $a \in Act$,

$$\text{i) } E \xrightarrow{a} E' \;\Rightarrow\; F \xrightarrow{a} F' \wedge (E', F') \in \equiv \cdot S \cdot \equiv$$

$$\text{ii) } F \xrightarrow{a} F' \;\Rightarrow\; E \xrightarrow{a} E' \wedge (E', F') \in \equiv \cdot S \cdot \equiv$$

Lemma

If S is a (strict) bisimulation up to $\equiv$, then $S \subseteq \sim$

- To prove $Sem_n \sim Sem^n$ a bisimulation will contain 2^n pairs, while a bisimulation up to $\equiv$ only requires $n+1$ pairs.

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A ∼-calculus

Lemma

$$E \equiv F \;\Rightarrow\; E \sim F$$

- proof idea: show that $\{(E+E, E) \mid E \in \mathbb{P}\} \cup Id_{\mathbb{P}}$ is a bisimulation

Lemma

$$\text{new } K' \; (\text{new } K \;\, E) \sim \text{new } (K \cup K') \; E$$

$$\text{new } K \;\, E \sim E \qquad \text{if } \mathbb{L}(E) \cap (K \cup \overline{K}) = \emptyset$$

$$\text{new } K \; (E \mid F) \sim \text{new } K \;\, E \mid \text{new } K \;\, F \qquad \text{if } \mathbb{L}(E) \cap \overline{\mathbb{L}(F)} \cap (K \cup \overline{K}) = \emptyset$$

- proof idea: discuss whether S is a bisimulation:

$$S \;=\; \{(\text{new } K \;\, E, E) \mid E \in \mathbb{P} \wedge \mathbb{L}(E) \cap (K \cup \overline{K}) = \emptyset\}$$

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$\sim$ is a congruence

congruence is the name of modularity in Mathematics

- process combinators preserve $\sim$

Lemma

Assume $E \sim F$. Then,

$$a.E \sim a.F$$
$$E + P \sim F + P$$
$$E \mid P \sim F \mid P$$
$$\text{new } K \; E \sim \text{new } K \; F$$

- recursive definition preserves $\sim$

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$\sim$ is a congruence

- First $\sim$ is extended to processes with variables:

$$E \sim F \;\equiv\; \forall_{\tilde{P}} \,.\; \{\tilde{P}/\tilde{X}\}\, E \sim \{\tilde{P}/\tilde{X}\}\, F$$

- Then prove:

Lemma

i) $\tilde{P} \triangleq \tilde{E} \;\Rightarrow\; \tilde{P} \sim \tilde{E}$
where $\tilde{E}$ is a family of process expressions and $\tilde{P}$ a family of process identifiers.

ii) Let $\tilde{E} \sim \tilde{F}$, where $\tilde{E}$ and $\tilde{F}$ are families of recursive process expressions over a family of process variables $\tilde{X}$, and define:

$$\tilde{A} \triangleq \{\tilde{A}/\tilde{X}\}\, \tilde{E} \;\text{ and }\; \tilde{B} \triangleq \{\tilde{B}/\tilde{X}\}\, \tilde{F}$$

Then

$$\tilde{A} \sim \tilde{B}$$

The expansion theorem

Every process is equivalent to the sum of its derivatives

$$E \sim \sum\{a.E' \mid E \xrightarrow{a} E'\}$$

understood?

$$E \sim \sum\{a.E' \mid E \xrightarrow{a} E'\}$$

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The expansion theorem

The usual definition (based on the concurrent canonical form):

$$E \sim \sum \{ f_i(a).\text{new}\, K\, (\{f_1\}\, E_1 \mid \ldots \mid \{f_i\}\, E_i' \mid \ldots \mid \{f_n\}\, E_n) \mid E_i \xrightarrow{a} E_i' \wedge f_i(a) \notin K \cup \overline{K} \}$$
$$+$$
$$\sum \{ \tau.\text{new}\, K\, (\{f_1\}\, E_1 \mid \ldots \mid \{f_i\}\, E_i' \mid \ldots \mid \{f_j\}\, E_j' \mid \ldots \mid \{f_n\}\, E_n) \mid E_i \xrightarrow{a} E_i' \wedge E_j \xrightarrow{b} E_j' \wedge f_i(a) = \overline{f_j(b)} \}$$

for $E \triangleq \text{new}\, K\, (\{f_1\}\, E_1 \mid \ldots \mid \{f_n\}\, E_n)$, with $n \geq 1$

The expansion theorem

Corollary (for $n = 1$ and $f_1 = \mathrm{id}$)

$$\text{new } K\ (E + F) \sim \text{new } K\ E + \text{new } K\ F$$

$$\text{new } K\ (a.E) \sim \begin{cases} \mathbf{0} & \text{if } a \in (K \cup \overline{K}) \\ a.(\text{new } K\ E) & \text{otherwise} \end{cases}$$

Example

$S \sim M$

$$
\begin{aligned}
S &\sim \text{new}\,\{k\}\;(T \mid R) \\
&\sim i.\text{new}\,\{k\}\;(\overline{k}.T \mid R) \\
&\sim i.\tau.\text{new}\,\{k\}\;(T \mid j.R) \\
&\sim i.\tau.(i.\text{new}\,\{k\}\;(\overline{k}.T \mid j.R) + j.\text{new}\,\{k\}\;(T \mid R)) \\
&\sim i.\tau.(i.j.\text{new}\,\{k\}\;(\overline{k}.T \mid R) + j.i.\text{new}\,\{k\}\;(\overline{k}.T \mid R)) \\
&\sim i.\tau.(i.j.\tau.\text{new}\,\{k\}\;(T \mid j.R) + j.i.\tau.\text{new}\,\{k\}\;(T \mid j.R))
\end{aligned}
$$

Let $N' = \text{new}\,\{k\}\;(T \mid j.R)$.
This expands into $N' \sim i.j.\tau.\text{new}\,\{k\}\;(T \mid j.R) + j.i.\tau.\text{new}\,\{k\}\;(T \mid j.R)$,
Therefore $N' \sim N$ and $S \sim i.\tau.N \sim M$

- requires result on unique solutions for recursive process equations

Observable transitions

$$\stackrel{a}{\Longrightarrow} \subseteq \mathbb{P} \times \mathbb{P}$$

- $L \cup \{\epsilon\}$
- A $\stackrel{\epsilon}{\Longrightarrow}$-transition corresponds to zero or more non observable transitions
- inference rules for $\stackrel{a}{\Longrightarrow}$:

$$\frac{}{E \stackrel{\epsilon}{\Longrightarrow} E}\,(O_1)$$

$$\frac{E \stackrel{\tau}{\longrightarrow} E' \quad E' \stackrel{\epsilon}{\Longrightarrow} F}{E \stackrel{\epsilon}{\Longrightarrow} F}\,(O_2)$$

$$\frac{E \stackrel{\epsilon}{\Longrightarrow} E' \quad E' \stackrel{a}{\longrightarrow} F' \quad F' \stackrel{\epsilon}{\Longrightarrow} F}{E \stackrel{a}{\Longrightarrow} F}\,(O_3) \quad \text{for } a \in L$$

Example

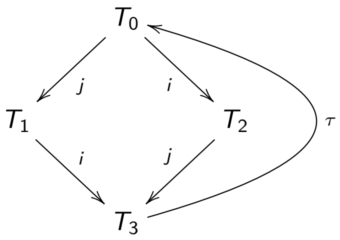
$$T_0 \triangleq j.T_1 + i.T_2$$
$$T_1 \triangleq i.T_3$$
$$T_2 \triangleq j.T_3$$
$$T_3 \triangleq \tau.T_0$$

and

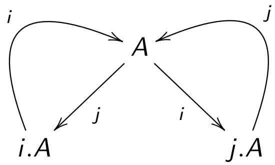
$$A \triangleq i.j.A + j.i.A$$

Example

From their graphs,



and



we conclude that $T_0 \approx A$ (why?).

Observational equivalence

$E \approx F$

- Processes E, F are observationally equivalent if there exists a weak bisimulation S st $\{\langle E, F\rangle\} \in S$.
- A binary relation S in $\mathbb{P}$ is a weak bisimulation iff, whenever $(E, F) \in S$ and $a \in L \cup \{\epsilon\}$,

$$\begin{array}{rl} \text{i)} & E \stackrel{a}{\Longrightarrow} E' \;\Rightarrow\; F \stackrel{a}{\Longrightarrow} F' \;\wedge\; (E', F') \in S \\ \text{ii)} & F \stackrel{a}{\Longrightarrow} F' \;\Rightarrow\; E \stackrel{a}{\Longrightarrow} E' \;\wedge\; (E', F') \in S \end{array}$$

I.e.,

$$\approx \;=\; \bigcup\{S \subseteq \mathbb{P} \times \mathbb{P} \mid S \text{ is a weak bisimulation}\}$$

Observational equivalence

Properties

- as expected: $\approx$ is an equivalence relation
- basic property: for any $E \in \mathbb{P}$,

$$E \;\approx\; \tau.E$$

  (proof idea: $\mathrm{id}_{\mathbb{P}} \cup \{(E, \tau.E) \mid E \in \mathbb{P}\}$ is a weak bisimulation

- weak vs. strict:

$$\sim \subseteq \approx$$

Is $\approx$ a congruence?

Lemma

Let $E \approx F$. Then, for any $P \in \mathbb{P}$ and $K \subseteq L$,

$$a.E \approx a.F$$
$$E \mid P \approx F \mid P$$
$$\text{new}\, K\ E \approx \text{new}\, K\ F$$

but

$$E + P \approx F + P$$

does not hold, in general.

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Is ≈ a congruence?

Example (initial τ restricts options 'menu')

$$i.\mathbf{0} \approx \tau.i.\mathbf{0}$$

However

$$j.\mathbf{0} + i.\mathbf{0} \not\approx j.\mathbf{0} + \tau.i.\mathbf{0}$$

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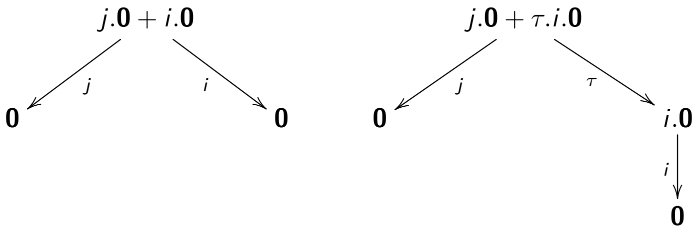

Forcing a congruence: $E = F$

Solution: force any initial τ to be matched by another τ

Process equality

Two processes E and F are equal (or observationally congruent) iff

i) $E \approx F$

ii) $E \xrightarrow{\tau} E' \;\Rightarrow\; F \xrightarrow{\tau} X \overset{\epsilon}{\Longrightarrow} F'$ and $E' \approx F'$

iii) $F \xrightarrow{\tau} F' \;\Rightarrow\; E \xrightarrow{\tau} X \overset{\epsilon}{\Longrightarrow} E'$ and $E' \approx F'$

- note that $E \neq \tau.E$, but $\tau.E = \tau.\tau.E$

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Forcing a congruence: $E = F$

$=$ can be regarded as a restriction of $\approx$ to all pairs of processes which preserve it in additive contexts

Lemma

Let E and F be processes st the union of their sorts is distinct of L. Then,

$$E = F \;\equiv\; \forall_{G \in \mathbb{P}} \,.\, (E + G \approx F + G)$$

Properties of $=$

Lemma

$$E = F \;\equiv\; (E = F) \vee (E = \tau.F) \vee (\tau.E = F)$$

- note that $E \neq \tau.E$, but $\tau.E = \tau.\tau.E$

Properties of $=$

Lemma

$$\sim \;\subseteq\; = \;\subseteq\; \approx$$

So,

the whole $\sim$ theory remains valid

Additionally,

Lemma (additional laws)

$$a.\tau.E = a.E$$
$$E + \tau.E = \tau.E$$
$$a.(E + \tau.F) = a.(E + \tau.F) + a.F$$

Solving equations

Have equations over $(\mathbb{P}, \sim)$ or $(\mathbb{P}, =)$ (unique) solutions?

Lemma

Recursive equations $\tilde{X} = \tilde{E}(\tilde{X})$ or $\tilde{X} \sim \tilde{E}(\tilde{X})$, over $\mathbb{P}$, have unique solutions (up to $=$ or $\sim$, respectively). Formally,

i) Let $\tilde{E} = \{E_i \mid i \in I\}$ be a family of expressions with a maximum of I free variables ($\{X_i \mid i \in I\}$) such that any variable free in E_i is weakly guarded. Then

$$\tilde{P} \sim \{\tilde{P}/\tilde{X}\}\tilde{E} \wedge \tilde{Q} \sim \{\tilde{Q}/\tilde{X}\}\tilde{E} \;\Rightarrow\; \tilde{P} \sim \tilde{Q}$$

ii) Let $\tilde{E} = \{E_i \mid i \in I\}$ be a family of expressions with a maximum of I free variables ($\{X_i \mid i \in I\}$) such that any variable free in E_i is guarded and sequential. Then

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Conditions on variables

guarded :
: X occurs in a sub-expression of type $a.E'$ for $a \in Act - \{\tau\}$

weakly guarded :
: X occurs in a sub-expression of type $a.E'$ for $a \in Act$

in both cases assures that, until a guard is reached, behaviour does not depends on the process that instantiates the variable

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Conditions on variables

sequential :

X is sequential in E if every strict sub-expression in which X occurs is either $a.E'$, for $a \in Act$, or $\Sigma\tilde{E}$.

avoids X to become guarded by a τ as a result of an interaction

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Example (1)

Consider

$$\begin{aligned} Sem &\triangleq get.put.Sem \\ P_1 &\triangleq \overline{get}.c_1.\overline{put}.P_1 \\ P_2 &\triangleq \overline{get}.c_2.\overline{put}.P_2 \\ S &\triangleq \text{new}\,\{get, put\}\;(Sem \mid P_1 \mid P_2) \end{aligned}$$

and

$$S' \triangleq \tau.c_1.S' + \tau.c_2.S'$$

to prove $S \sim S'$, show both are solutions of

$$X \;=\; \tau.c_1.X + \tau.c_2.X$$

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$$\begin{aligned} Sem &\triangleq get.put.Sem \\ P_1 &\triangleq \overline{get}.c_1.\overline{put}.P_1 \\ P_2 &\triangleq \overline{get}.c_2.\overline{put}.P_2 \\ S &\triangleq \text{new}\,\{get, put\}\;(Sem \mid P_1 \mid P_2) \end{aligned}$$

and

$$S' \triangleq \tau.c_1.S' + \tau.c_2.S'$$

to prove $S \sim S'$, show both are solutions of

$$X \;=\; \tau.c_1.X + \tau.c_2.X$$

Example (1)

proof

$$
\begin{aligned}
S = \; & \tau.\text{new}\, K\ (c_1.\overline{put}.P_1 \mid P_2 \mid put.Sem) + \tau.\text{new}\, K\ (P_1 \mid c_2.\overline{put}.P_2 \mid put.Sem) \\
= \; & \tau.c_1.\text{new}\, K\ (\overline{put}.P_1 \mid P_2 \mid put.Sem) + \tau.c_2.\text{new}\, K\ (P_1 \mid \overline{put}.P_2 \mid put.Sem) \\
= \; & \tau.c_1.\tau.\text{new}\, K\ (P_1 \mid P_2 \mid Sem) + \tau.c_2.\tau.\text{new}\, K\ (P_1 \mid P_2 \mid Sem) \\
= \; & \tau.c_1.\tau.S + \tau.c_2.\tau.S \\
= \; & \tau.c_1.S + \tau.c_2.S \\
= \; & \{S/X\}E
\end{aligned}
$$

for S' is immediate

Example (2)

Consider,

$$B \triangleq in.B_1 \qquad B' \triangleq \text{new } m\ (C_1 \mid C_2)$$

$$B_1 \triangleq in.B_2 + \overline{out}.B \qquad C_1 \triangleq in.\overline{m}.C_1$$

$$C_2 \triangleq m.\overline{out}.C_2$$

B' is a solution of

$$X = E(X, Y, Z) = in.Y$$

$$Y = E_1(X, Y, Z) = in.Z + \overline{out}.X$$

$$Z = E_3(X, Y, Z) = \overline{out}.Y$$

through $\sigma = \{B/X, B_1/Y, B_2/Z\}$

Example (2)

To prove B =B'

$$\begin{aligned} B' &= \text{new}\, m\, (C_1 \mid C_2) \\ &= in.\text{new}\, m\, (\overline{m}.C_1 \mid C_2) \\ &= in.\tau.\text{new}\, m\, (C_1 \mid \overline{out}.C_2) \\ &= in.\text{new}\, m\, (C_1 \mid \overline{out}.C_2) \end{aligned}$$

Let $S_1 = \text{new}\, m\, (C_1 \mid \overline{out}.C_2)$ to proceed:

$$\begin{aligned} S_1 &= \text{new}\, m\, (C_1 \mid \overline{out}.C_2) \\ &= in.\text{new}\, m\, (\overline{m}.C_1 \mid \overline{out}.C_2) + \overline{out}.\text{new}\, m\, (C_1 \mid C_2) \\ &= in.\text{new}\, m\, (\overline{m}.C_1 \mid \overline{out}.C_2) + \overline{out}.B' \end{aligned}$$

Example (2)

Finally, let, $S_2 = \text{new } m\ (\overline{m}.C_1 \mid \overline{out}.C_2)$. Then,

$$
\begin{aligned}
S_2 &= \text{new } m\ (\overline{m}.C_1 \mid \overline{out}.C_2) \\
&= \overline{out}.\text{new } m\ (\overline{m}.C_1 \mid C_2) \\
&= \overline{out}.\tau.\text{new } m\ (C_1 \mid \overline{out}.C_2) \\
&= \overline{out}.\tau.S_1 \\
&= \overline{out}.S_1
\end{aligned}
$$

Example (2)

Note the same problem can be solved with a system of 2 equations:

$$\begin{aligned} X &= E(X,Y) = in.Y \\ Y &= E'(X,Y) = in.\overline{out}.Y + \overline{out}.in.Y \end{aligned}$$

Clearly, by substitution,

$$\begin{aligned} B &= in.B_1 \\ B_1 &= in.\overline{out}.B_1 + \overline{out}.in.B_1 \end{aligned}$$

Example (2)

On the other hand, it's already proved that $B' = ... = in.S_1$.
so,

$$\begin{aligned} S_1 &= \text{new}\, m\, (C_1 \mid \overline{out}.C_2) \\ &= in.\text{new}\, m\, (\overline{m}.C_1 \mid \overline{out}.C_2) + \overline{out}.B' \\ &= in.\overline{out}.\text{new}\, m\, (\overline{m}.C_1 \mid C_2) + \overline{out}.B' \\ &= in.\overline{out}.\tau.\text{new}\, m\, (C_1 \mid \overline{out}.C_2) + \overline{out}.B' \\ &= in.\overline{out}.\tau.S_1 + \overline{out}.B' \\ &= in.\overline{out}.S_1 + \overline{out}.B' \\ &= in.\overline{out}.S_1 + \overline{out}.in.S_1 \end{aligned}$$

Hence, $B' = \{B'/X, S_1/Y\}E$ and $S_1 = \{B'/X, S_1/Y\}E'$