

## Formulário

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### Sintaxe

$$E ::= A(v_1, \dots, v_n) \mid \mathbf{0} \mid a.E \mid E_0 + E_1 \mid E_0 \mid E_1 \mid \text{new } K E$$

### Congruência estrutural

- $\equiv$  é uma relação de *equivalência* fechada para conversão- $\alpha$  e preservada pelos operadores
- *composição paralela* e *escolha* formam, com  $\mathbf{0}$ , monoides abelianos
- (para  $a \notin \text{free}(P)$ )  $\text{new } \{a\} (P \mid Q) \equiv P \mid \text{new } \{a\} Q$  e  $\text{new } \{a\} \mathbf{0} \equiv \mathbf{0}$
- (para  $A(\tilde{x}) \triangleq P_A$ )  $A(\tilde{y}) \equiv \{\tilde{y}/\tilde{x}\} P_A$

### Semântica

$$\begin{array}{c} \frac{}{\alpha.E \xrightarrow{\alpha} E} (\text{prefix}) \quad \frac{E \xrightarrow{\alpha} E' \quad F \xrightarrow{\bar{\alpha}} F'}{E \mid F \xrightarrow{\tau} E' \mid F'} (\text{react}) \\ \frac{E \xrightarrow{\alpha} E'}{E \mid F \xrightarrow{\alpha} E' \mid F} (\text{par} - l) \quad \frac{F \xrightarrow{\alpha} F'}{E \mid F \xrightarrow{\alpha} E \mid F'} (\text{par} - r) \\ \frac{E \xrightarrow{\alpha} E'}{E + F \xrightarrow{\alpha} E'} (\text{sum} - l) \quad \frac{F \xrightarrow{\alpha} F'}{E + F \xrightarrow{\alpha} F'} (\text{sum} - r) \\ \frac{E \xrightarrow{\alpha} E'}{\text{new } \{\beta\} E \xrightarrow{\alpha} \text{new } \{\beta\} E'} (\text{res}) \quad (\text{se } \alpha \notin \{\beta, \bar{\beta}\}) \\ \frac{\{\tilde{\beta}/\tilde{\delta}\} E_A \xrightarrow{\alpha} E'}{A(\tilde{\beta}) \xrightarrow{\alpha} E'} (\text{ident}) \quad (\text{se } A(\tilde{\delta}) \triangleq E_A) \end{array}$$

### Teoria equivalência estrita

$$\begin{array}{c} (E + F) + G \sim E + (F + G) \\ E + F \sim F + E \\ E + \mathbf{0} \sim E \\ E + E \sim E \\ (E \mid F) \mid G \sim E \mid (F \mid G) \\ E \mid F \sim F \mid E \\ E \mid \mathbf{0} \sim E \end{array}$$

$$\begin{array}{ll}
\text{new } K \ E \sim E & \text{se } \mathbb{L}(E) \cap (K \cup \overline{K}) = \emptyset \\
\text{new } K' \ (\text{new } K \ E) \sim \text{new } K \cup K' \ E & \\
\text{new } K \ (E \mid F) \sim \text{new } K \ E \mid \text{new } K \ F & \text{se } \mathbb{L}(E) \cap \overline{\mathbb{L}(F)} \cap (K \cup \overline{K}) = \emptyset
\end{array}$$

$$\begin{aligned}
E \sim & \sum \{a.\text{new } K \ (E_1 \mid \dots \mid E'_i \mid \dots \mid E_n) \mid E_i \xrightarrow{a} E'_i \wedge a \notin K \cup \overline{K}\} \\
& + \\
& \sum \{\tau.\text{new } K \ (E_1 \mid \dots \mid E'_i \mid \dots \mid E'_j \mid \dots \mid E_n) \mid E_i \xrightarrow{a} E'_i \wedge E_j \xrightarrow{\bar{a}} E'_j\}
\end{aligned}$$

(para  $E \triangleq \text{new } K \ (E_1 \mid \dots \mid E_n)$ , com  $n \geq 0$ )

$$\begin{array}{l}
\text{new } K \ (E + F) \sim \text{new } K \ E + \text{new } K \ F \\
\text{new } K \ (a.E) \sim \begin{cases} \mathbf{0} & \text{se } a \in (K \cup \overline{K}) \\ a.(\text{new } K \ E) & \text{caso contrário} \end{cases}
\end{array}$$

## Equivalência Observacional

$$\frac{}{E \xRightarrow{\epsilon} E} (O_1)$$

$$\frac{E \xrightarrow{\tau} E' \quad E' \xRightarrow{\epsilon} F}{E \xRightarrow{\epsilon} F} (O_2)$$

$$\frac{E \xRightarrow{\epsilon} E' \quad E' \xrightarrow{a} F' \quad F' \xRightarrow{\epsilon} F}{E \xRightarrow{a} F} (O_3) \quad \text{para } a \in L$$

- $\approx$  é uma congruência, excepto para contextos aditivos arbitrários
- $E \approx \tau.E$

## Igualdade

- $a.\tau.E = a.E$  ,  $E + \tau.E = \tau.E$  ,  $a.(E + \tau.F) = a.(E + \tau.F) + a.F$
- $E = F$  sse, para todo o  $G \in \mathbb{P}$ ,  $E + G \approx F + G$
- $E \approx F$  sse  $E = F$  ou  $E = \tau.F$  ou  $\tau.E = F$
- $\sim \subseteq = \subseteq \approx$