

Coq Exercises

MAPi – PSVC

Porto, 2008

Please solve the following exercises *without* using automatic tactics such as `auto`, `tauto`, `omega`, etc.

1 Logical Reasoning

1. Prove the following propositional formulas ($P, Q : \text{Prop}$):

- (a) $(P \rightarrow Q) \rightarrow \sim Q \rightarrow \sim P$
- (b) $(P \vee Q \rightarrow R) \rightarrow (P \rightarrow R)$
- (c) $(P \vee Q) \wedge \sim P \rightarrow Q$
- (d) $P \wedge Q \rightarrow R \leftrightarrow P \rightarrow Q \rightarrow R$
- (e) $\sim(P \vee Q) \leftrightarrow \sim P \wedge \sim Q$
- (f) $\sim\sim(P \vee \sim P)$

2. Prove the following theorem:

Theorem $ex4 : \forall (X : \text{Set}) (P : X \rightarrow \text{Prop}), \sim(\forall x, \sim(P x)) \rightarrow (\text{exists } x, P x)$.

3. Assuming the *Excluded Middle* axiom, prove:

- (a) **Theorem** *Pierce* : $\forall P Q, ((P \rightarrow Q) \rightarrow P) \rightarrow P$.
- (b) **Theorem** *NNE* : $\forall P, \sim\sim P \rightarrow P$.

2 Natural Numbers

Recall the definition of natural numbers as an inductive type (command `Print nat`).

1. Define a function $add : nat \rightarrow nat \rightarrow nat$ that computes the sum of two natural numbers.
2. Prove commutativity of add (Hint: use *lemmas* to isolate interesting proof obligations).
3. Prove the following fact: *for a natural number n , $n * n$ can be computed as the sum of the first n odd numbers (e.g. $3 * 3 = 1 + 3 + 5$).* (Hint: define first a function that sums the odd numbers. You should also use definitions/facts/theorems from the library *Arith* — try also the `omega` tactic).

3 Reasoning about lists

The following exercises require the library `Lists`. You can load that library by executing the command `Require Import Lists`.

1. Consider the following inductive relation:

Inductive *last* (*A* : **Set**) (*x* : *A*) : *list A* → **Prop** :=
| *last_base* : *last x (cons x nil)*
| *last_step* : $\forall l y, \text{last } x \ l \rightarrow \text{last } x \ (\text{cons } y \ l)$.

- (a) Use `inversion` to prove that $\forall x, \sim(\text{last } x \ \text{nil})$.
 - (b) (Difficult) Try to avoid using that tactic.
2. Consider the following definition for the `Even` predicate:

Inductive *Even* : *nat* → **Prop** :=
| *Even_base* : *Even 0*
| *Even_step* : $\forall n, \text{Even } n \rightarrow \text{Even } (S \ (S \ n))$.

- (a) Define, accordingly, the `Odd` predicate. Prove that, for every number *n*, *Even n* → *Odd (S n)*.
 - (b) Define the function `rev` that reverses a list.
 - (c) Prove that, for every list *l*, *rev (rev l) = l*.
 - (d) Recall the definition for the function `app` (concatenation of lists). Prove that for every lists *l*₁ and *l*₂, *rev (app l1 l2) = app (rev l2) (rev l1)*.
3. Consider the inductive predicate:

Inductive *InL* (*A* : **Type**) (*a* : *A*) : *list A* → **Prop** :=
| *InHead* : $\forall (xs : \text{list } A), \text{InL } a \ (\text{cons } a \ xs)$
| *InTail* : $\forall (x : A) (xs : \text{list } A), \text{InL } a \ xs \rightarrow \text{InL } a \ (\text{cons } x \ xs)$.

Prove the following properties:

- (a) $\forall (A : \text{Type}) (a : A) (l1 \ l2 : \text{list } A), \text{InL } a \ l1 \vee \text{InL } a \ l2 \rightarrow \text{InL } a \ (\text{app } l1 \ l2)$.
 - (b) $\forall (A : \text{Type}) (a : A) (l1 \ l2 : \text{list } A), \text{InL } a \ (\text{app } l1 \ l2) \rightarrow \text{InL } a \ l1 \vee \text{InL } a \ l2$.
4. Define the function `elem` that checks if an element belongs a list of integers (Hint: import `ZArith` module in order to use the `Z_eq_dec` that tests for integer equality).
 5. Prove the correctness of `elem`, that is, $\forall (a : Z) (l1 \ l2 : \text{list } Z), \text{elem } a \ (\text{app } l1 \ l2) = \text{orb } (\text{elem } a \ l1) \ (\text{elem } a \ l2)$ (the function `orb` is the boolean-or function defined in Library `Bool`).