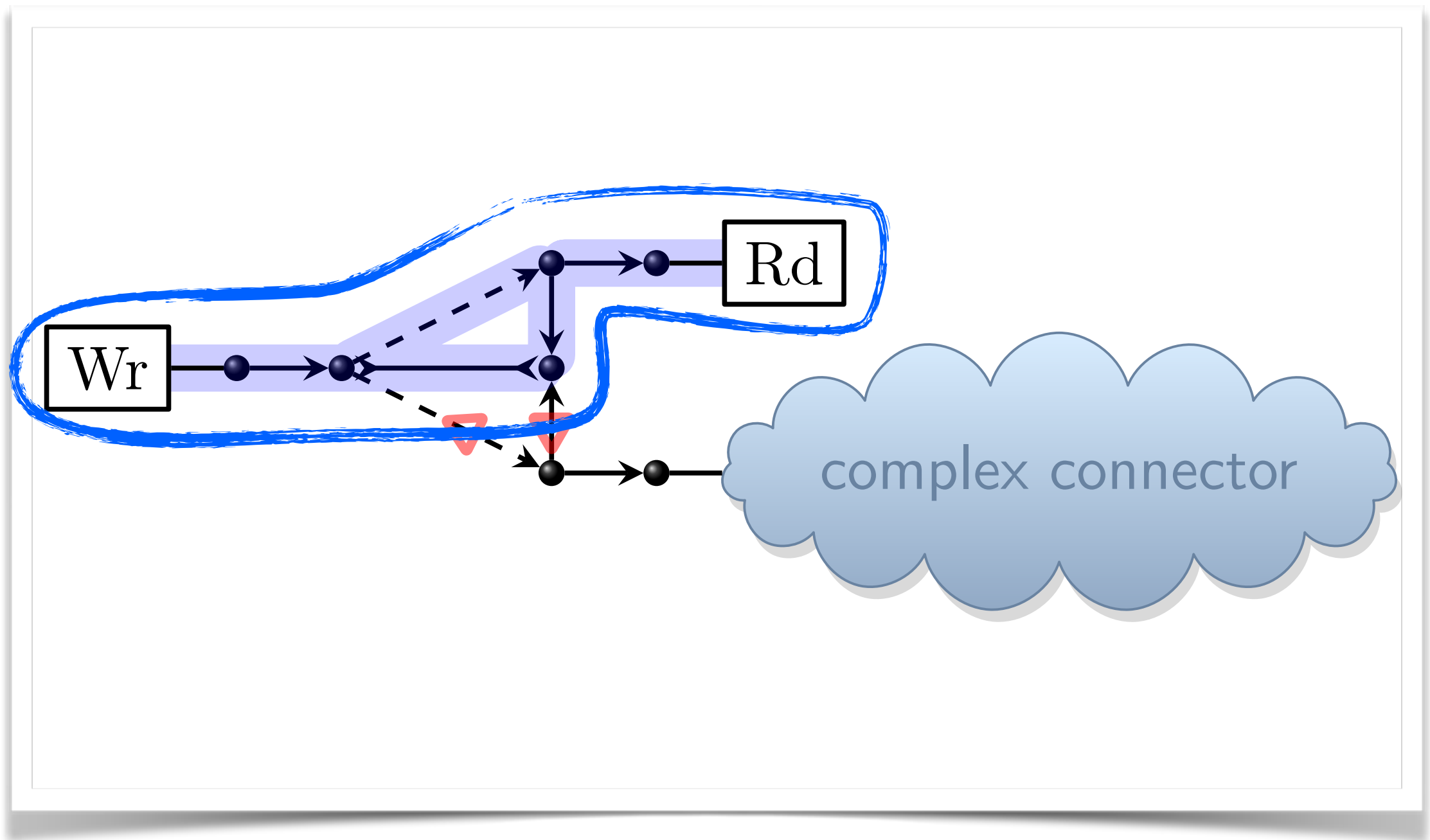


Architectural design: the coordination perspective

José Proença
HASLab - INESC TEC & UM
Arquitectura e Cálculo 2015-16





Locality (concurrency)

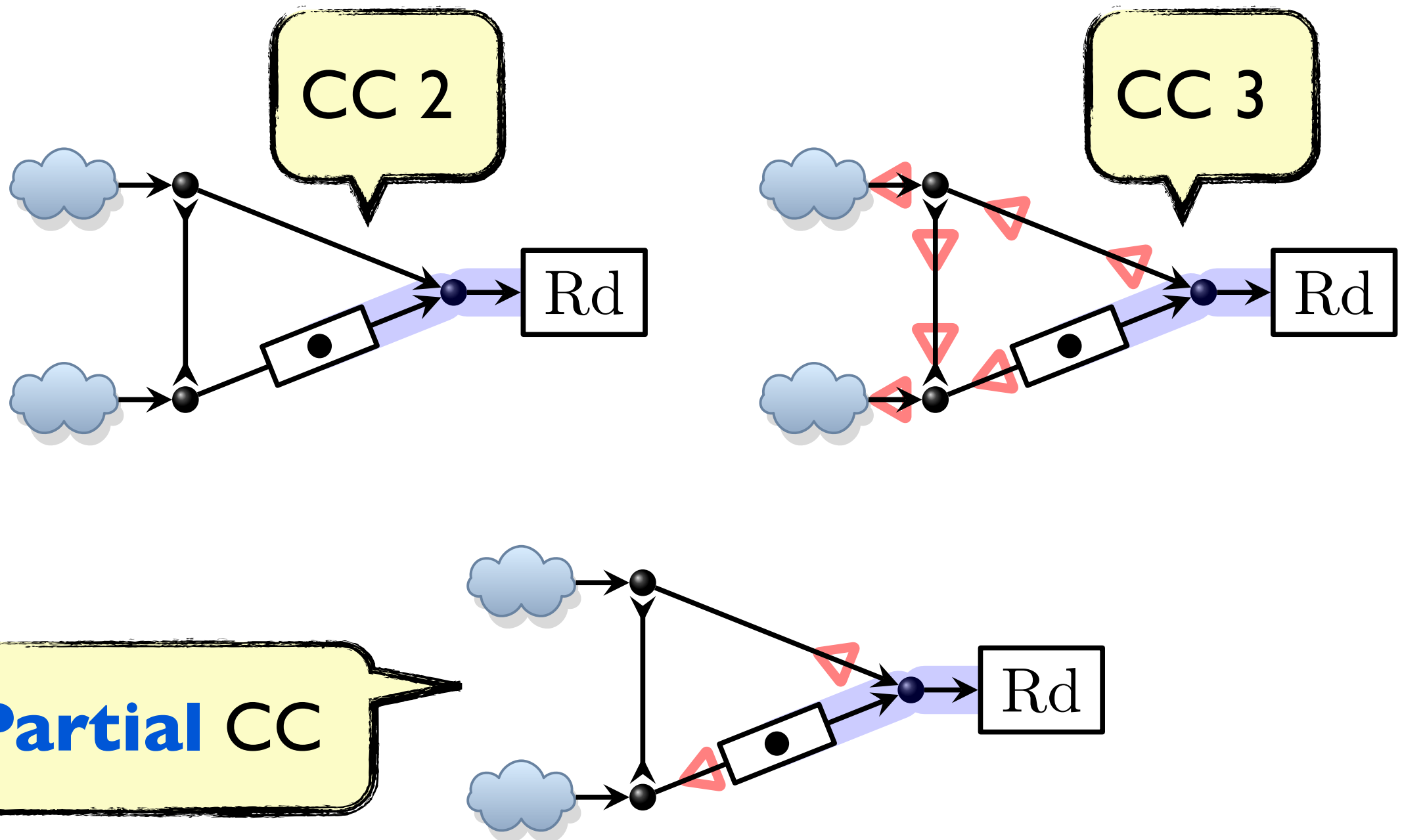
Motivation

- Connector colouring is not optimal for **distributed** systems.
- **All-or-nothing** – all channels are needed to decide where data goes.
- Need to identify **local flows** that are not composed with the full connector.
- Model **context** dependency

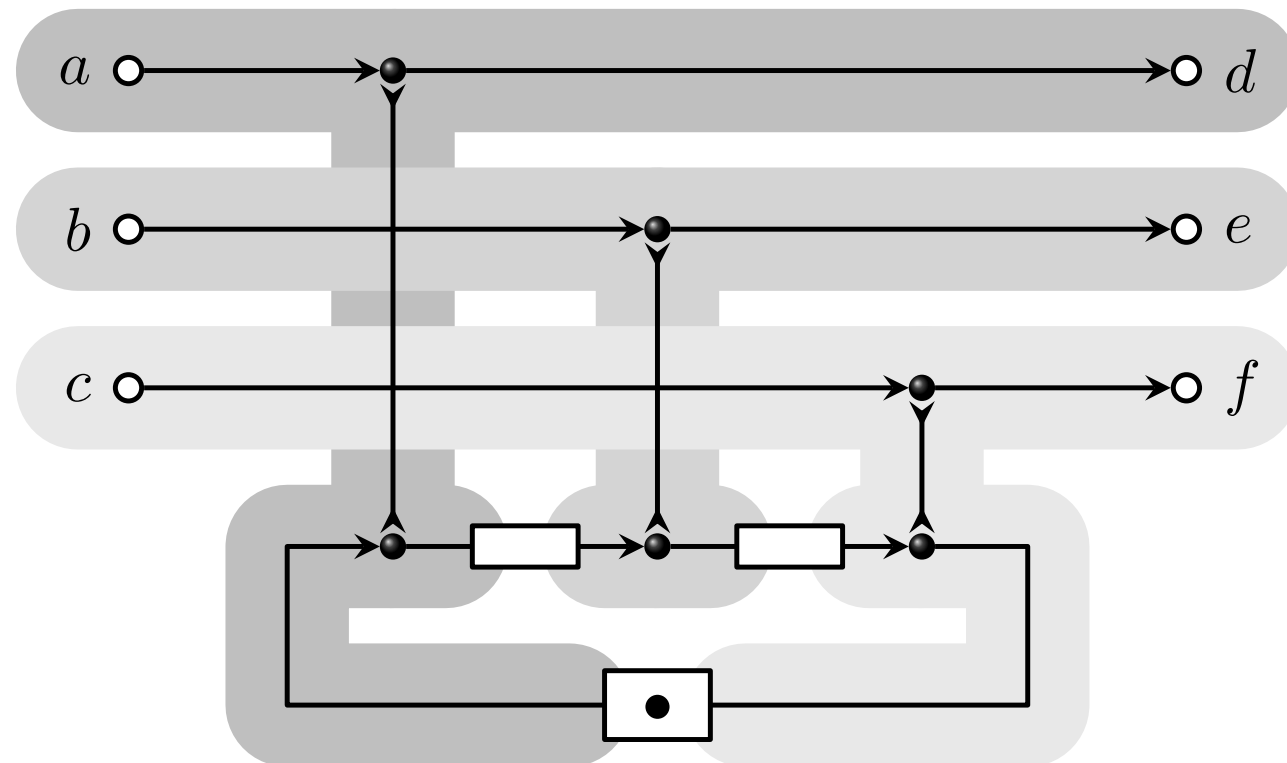
Problems

- 2 colours (or constraint automata):
 - ▶ **assume** primitives can make a *no-flow* step
- 3 colours:
 - ▶ **cannot assume** primitives have a no-flow colour – which direction would it be?
 - ▶ **Idea?**: add another no-flow colour, without direction, and assume all primitives have it...

Example

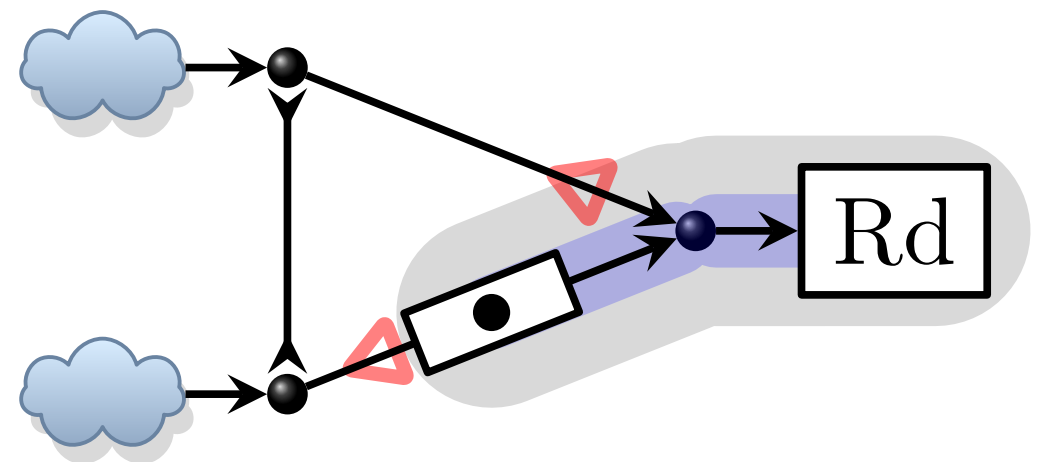


Synchronous regions



Static regions:
boundaries
=
FIFO's

Dynamic regions:
boundaries
=
GiveReason

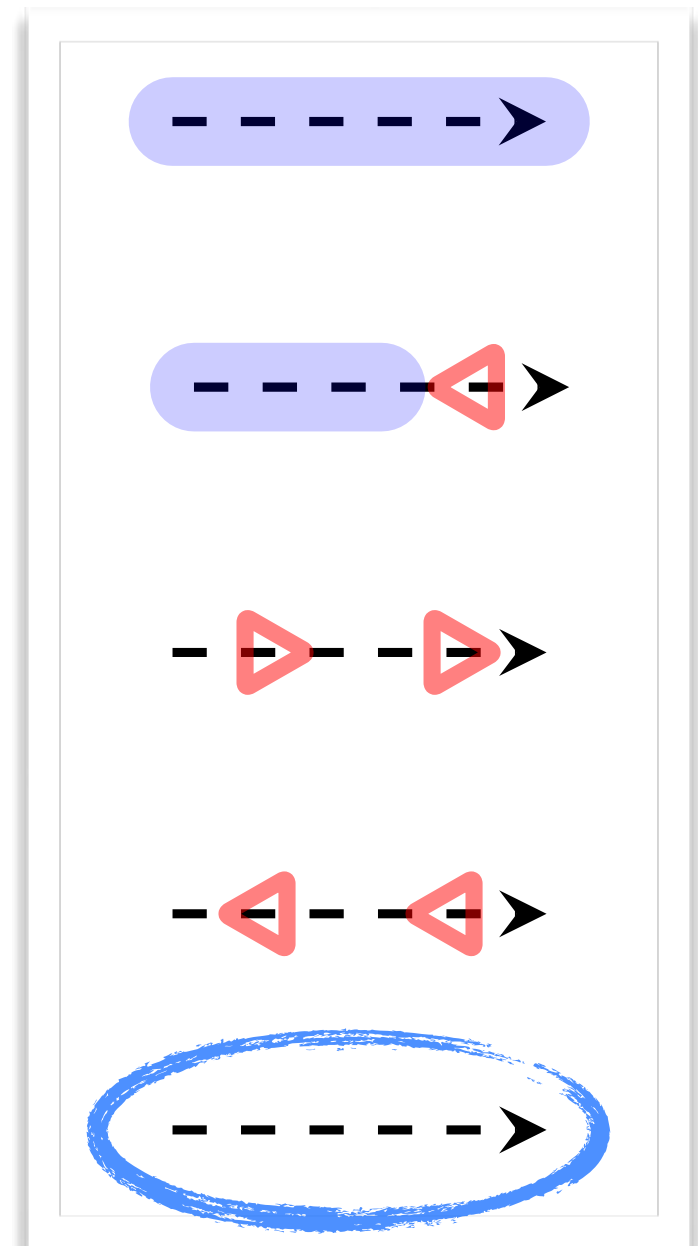
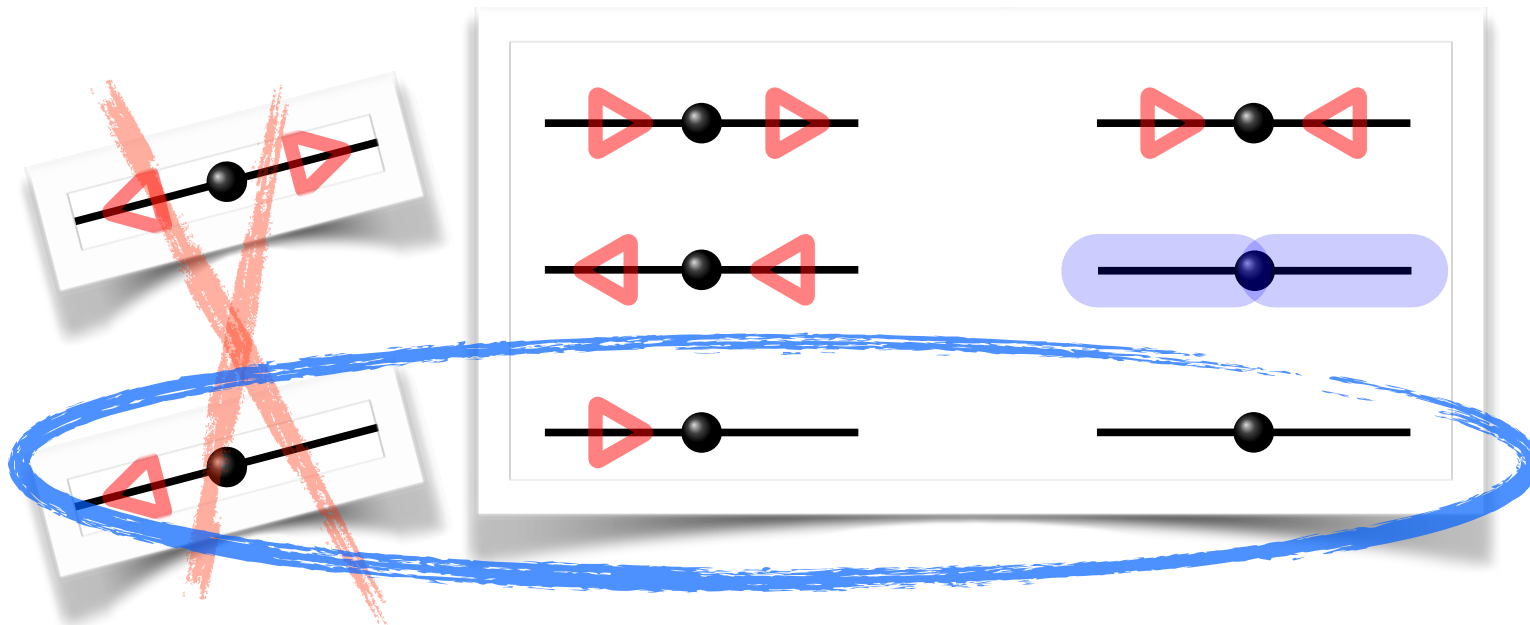


Partial connector colouring

- *Colouring:*

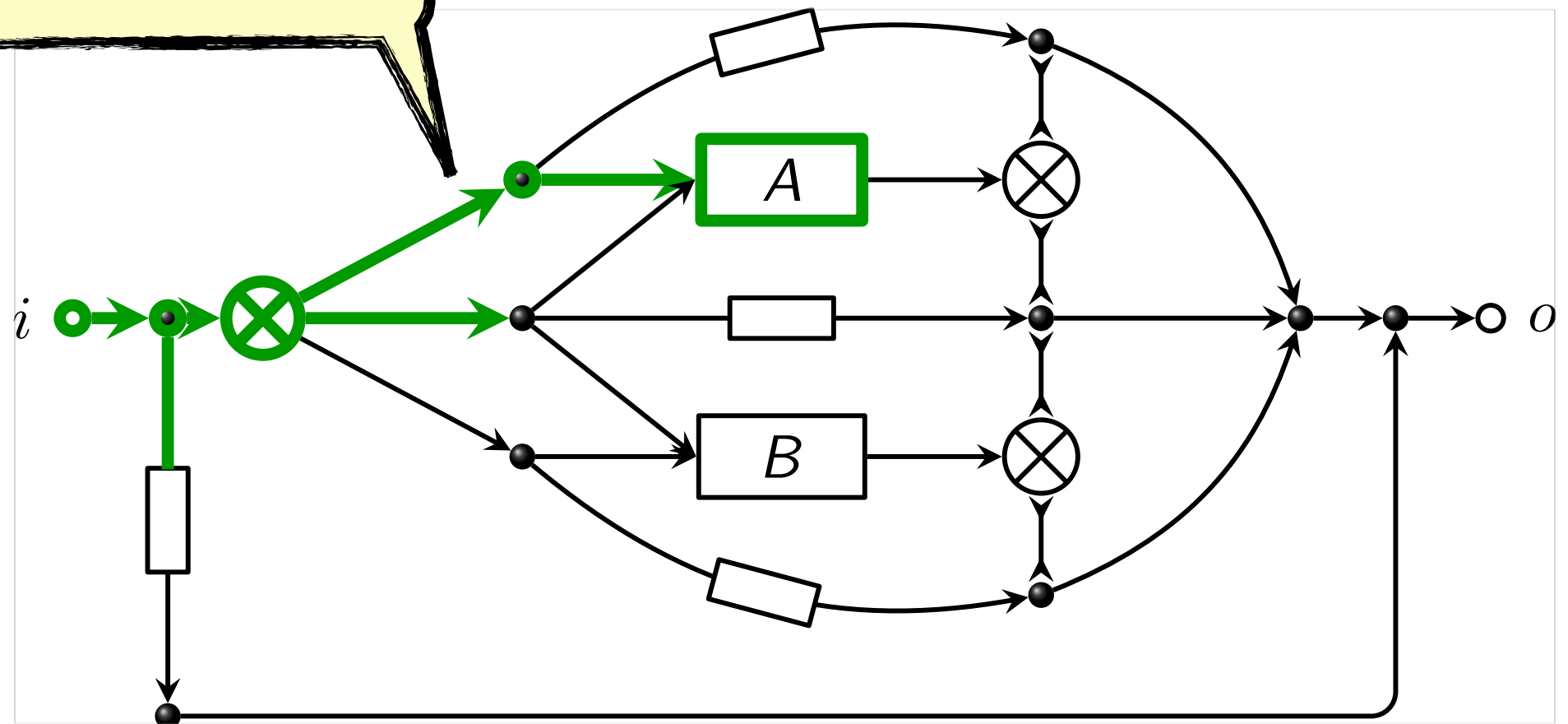
End $\rightarrow$ {Flow, GiveReason, GetReason}

- *Composition* = matching colours:



In practice

Shall I search now
for a colouring?



Outline

1. Visual semantics for Reo

- ▶ Connector colouring (CC)¹

2. Locality (concurrency)

- ▶ partial connector colouring (PCC)²

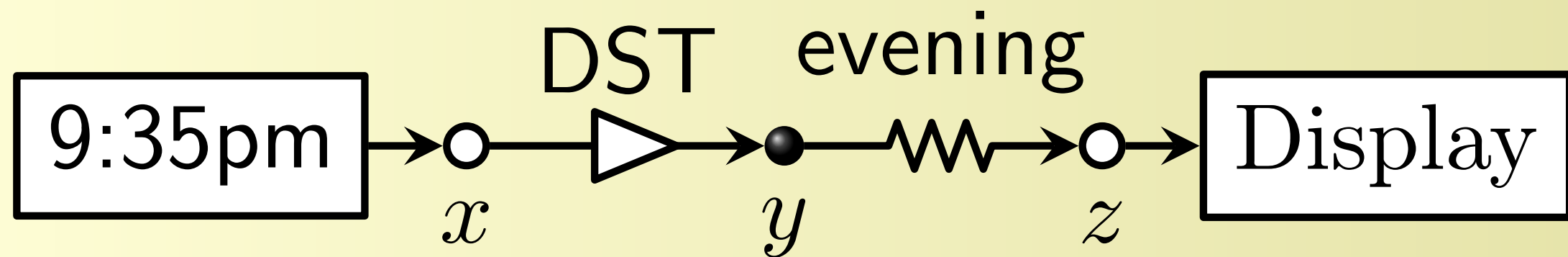
3. Constraints

- ▶ SAT solving with data for Reo³

¹ Dave Clarke, David Costa, and Farhad Arbab. Connector colouring I: Synchronisation and context dependency

² Dave Clarke and José Proença. Partial connector colouring

³ Dave Clarke, José Proença, Alexander Lazovik, and Farhad Arbab, Channel-based coordination via constraint satisfaction
José Proença, Dave Clarke, Interactive interaction constraints

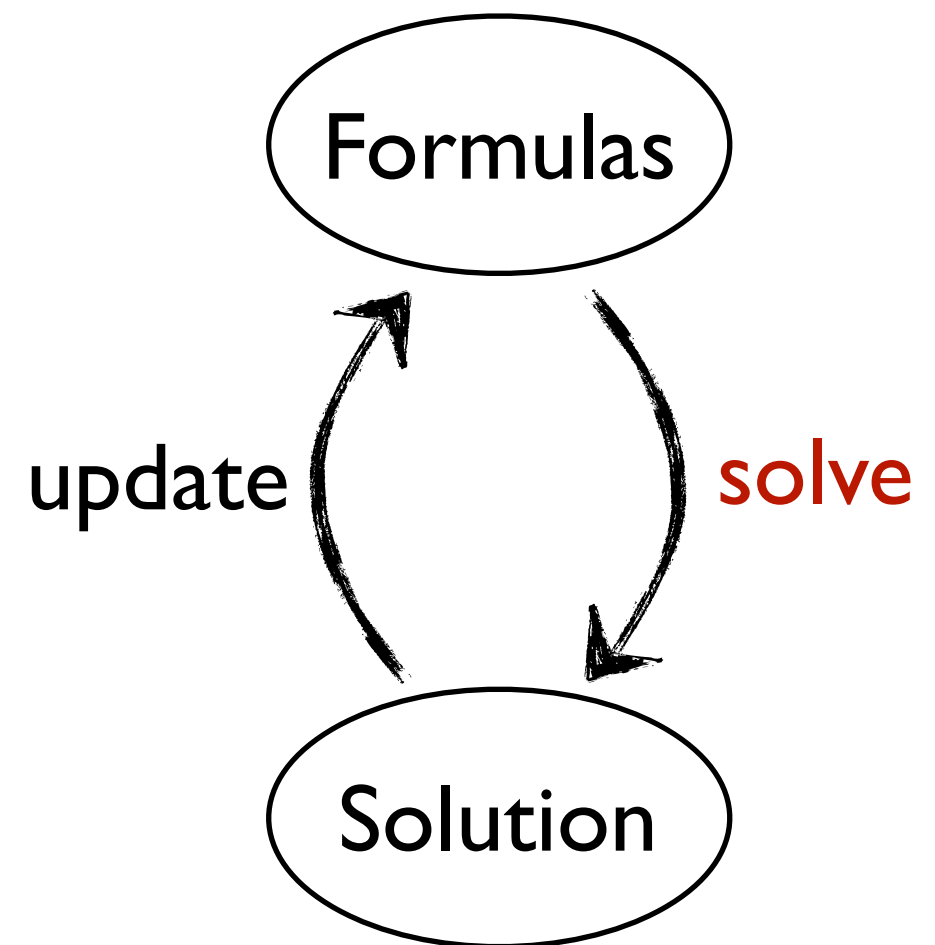


$$\begin{array}{lll}
 x \rightarrow \hat{x} := 9:35\text{pm} & x \leftrightarrow y & y \rightarrow \hat{y} := \text{DST}(\hat{x}) \\
 (y \wedge \text{evening}(\hat{y})) \leftrightarrow z & & z \rightarrow \hat{z} := \hat{y}
 \end{array}$$

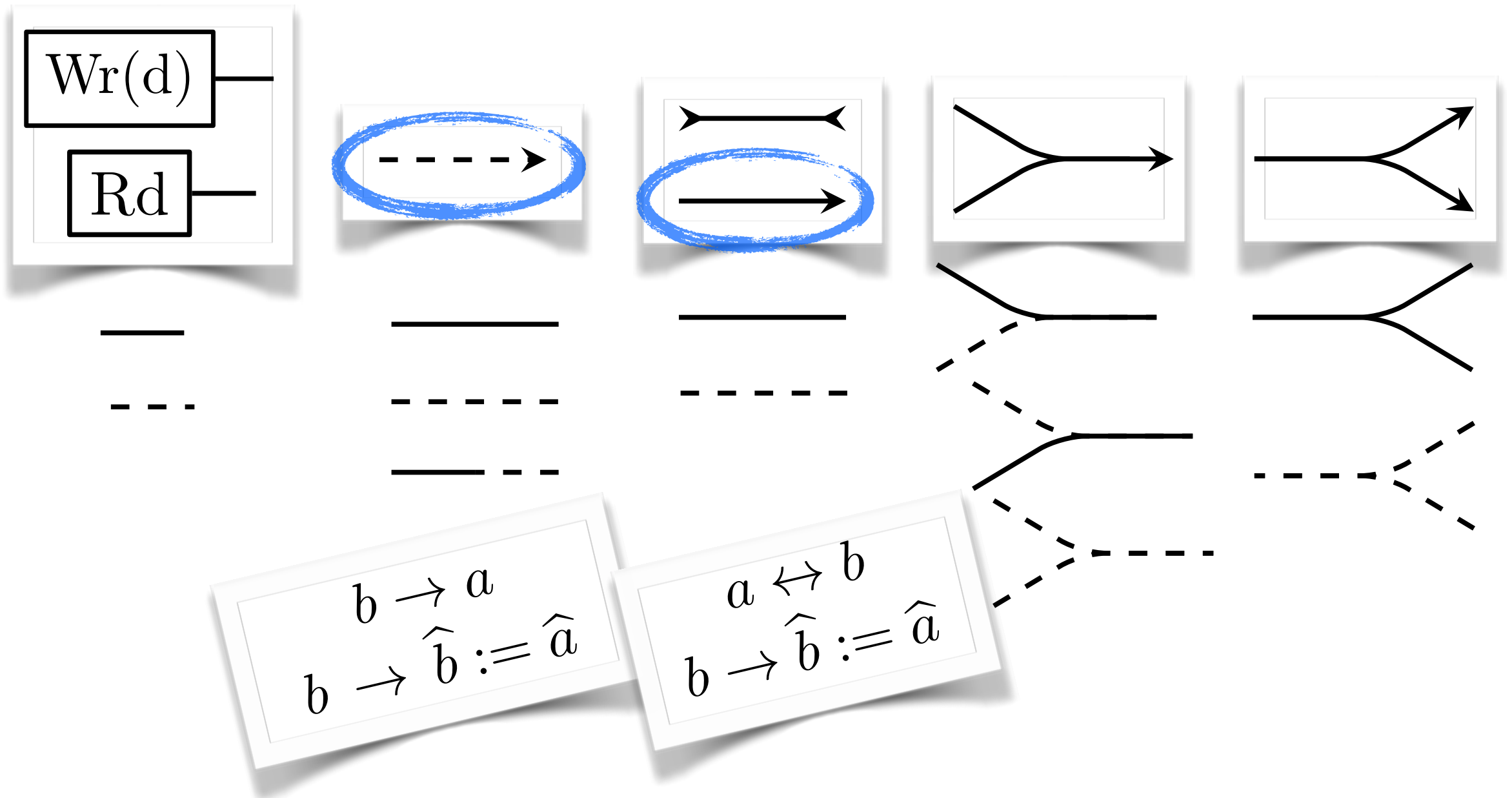
Reo as (Interactive) Constraints

Coordination as constraint satisfaction

$a \longrightarrow b$	$a \leftrightarrow b$ $b \rightarrow \hat{b} := \hat{a}$
$a \dashrightarrow b$	$b \rightarrow a$ $b \rightarrow \hat{b} := \hat{a}$
$a \overset{P}{\rightsquigarrow} b$	$b \rightarrow \hat{b} := \hat{a}$ $(a \wedge P(\hat{a})) \leftrightarrow b$
$a \overset{f}{\rightarrowtriangle} b$	$a \leftrightarrow b$ $b \rightarrow \hat{b} := f(\hat{a})$



Exercise: write constraints

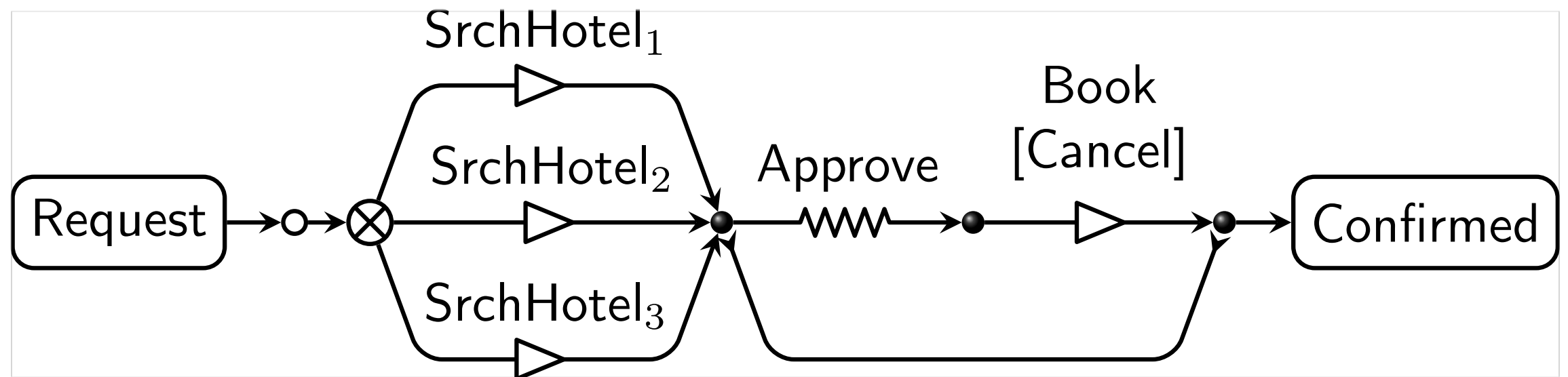


Building up constraints

- Connector colouring as constraints
- Data constraints
- Interactive constraints

Context
(need for extra interaction)

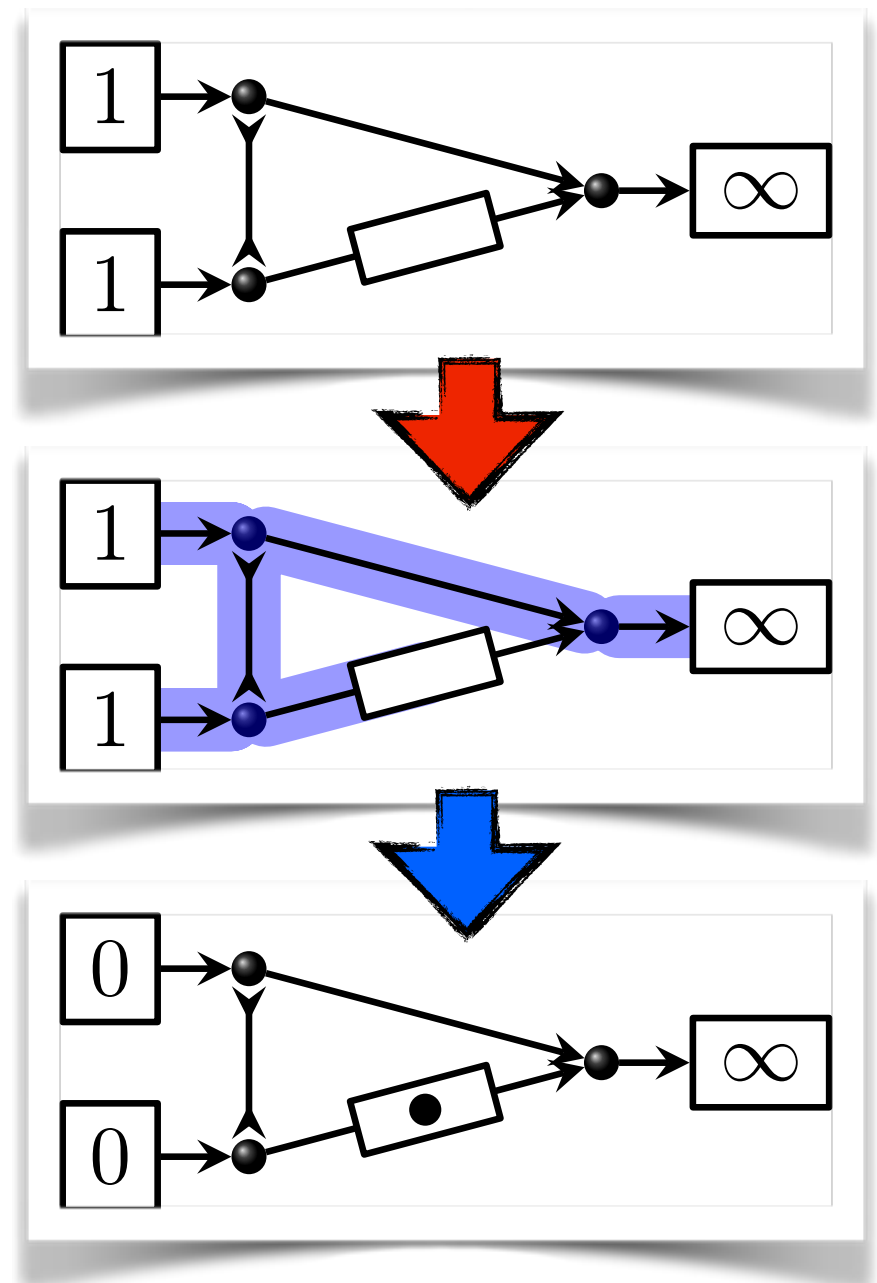
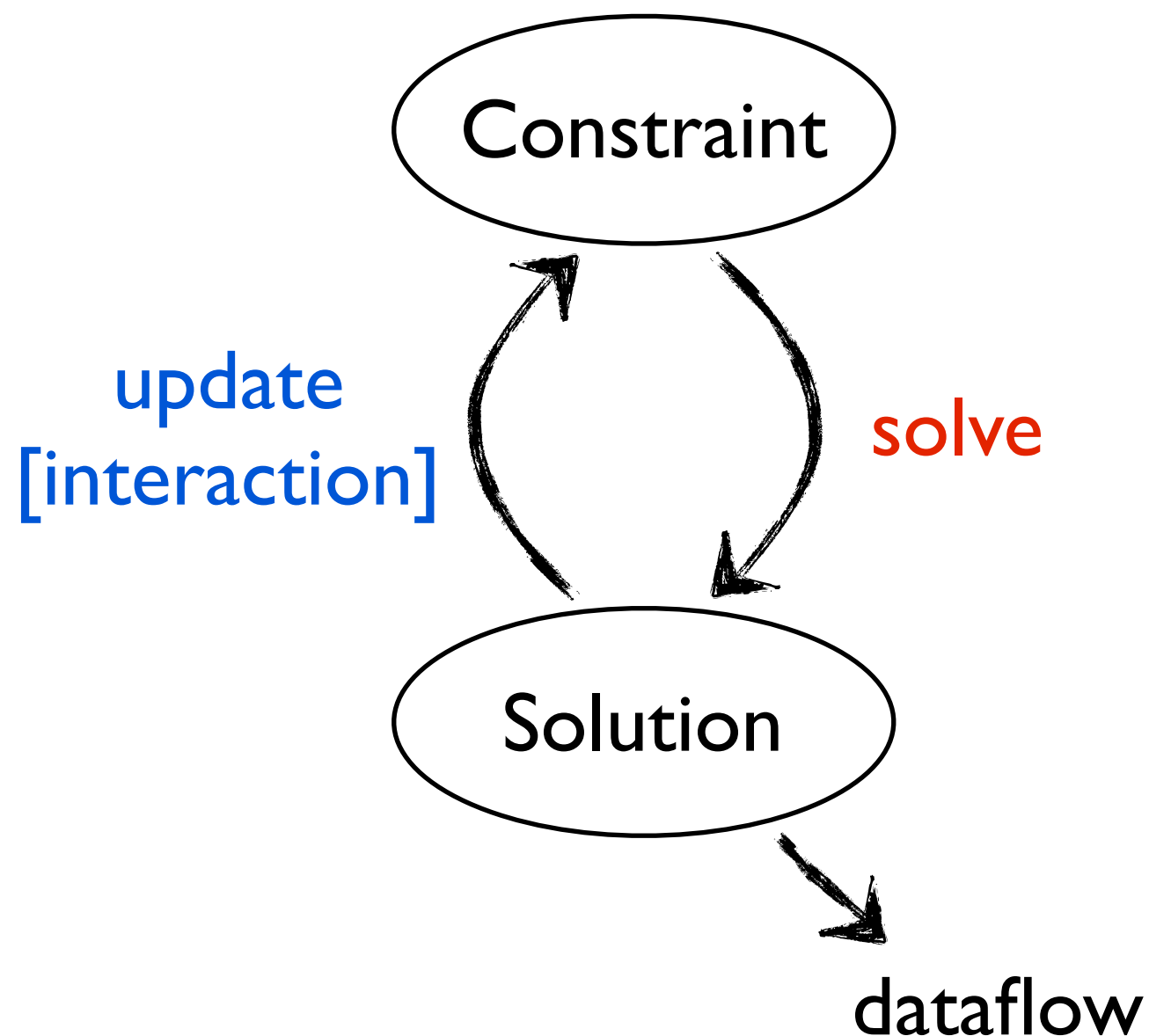
Hotel booking



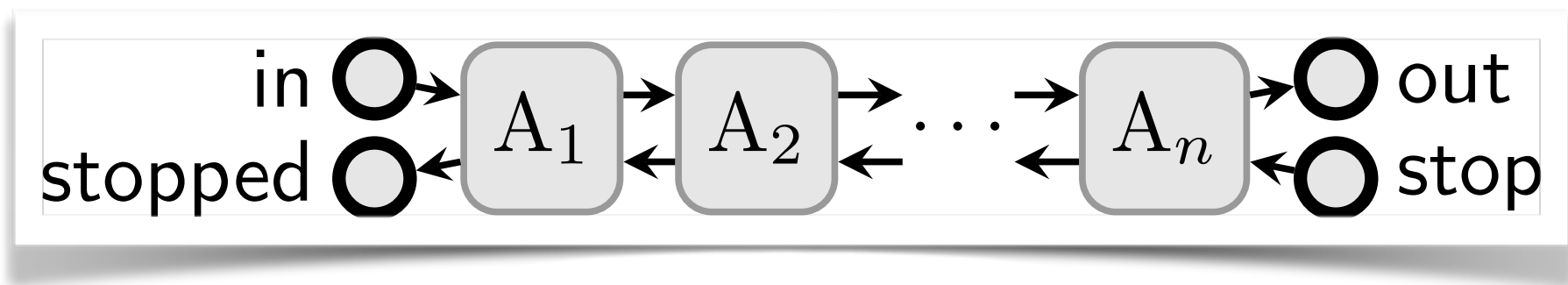
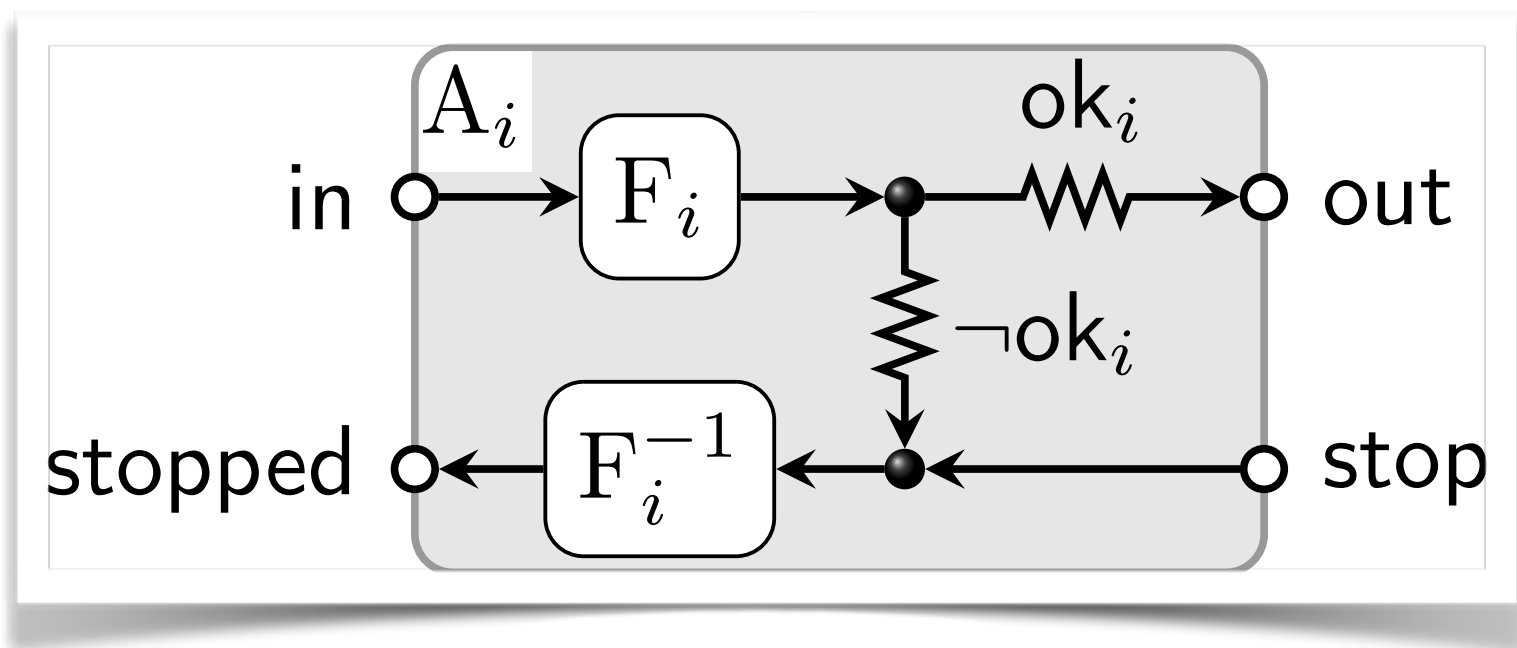
- Interaction with hotel repositories
- Interaction with users
- Interaction with hotels (availability & payment)

Problem of Interaction

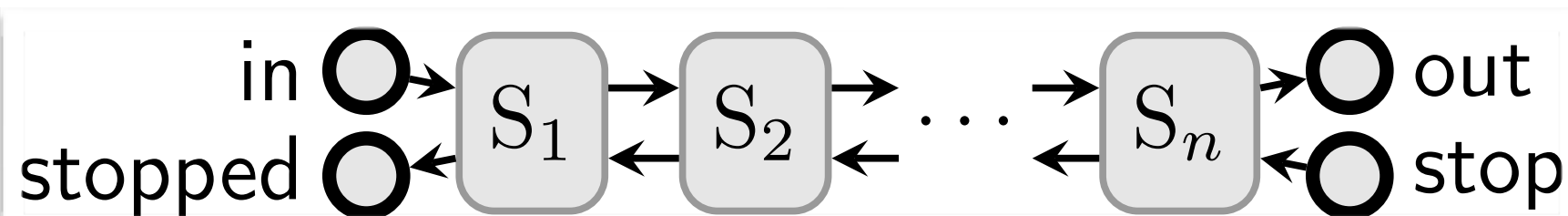
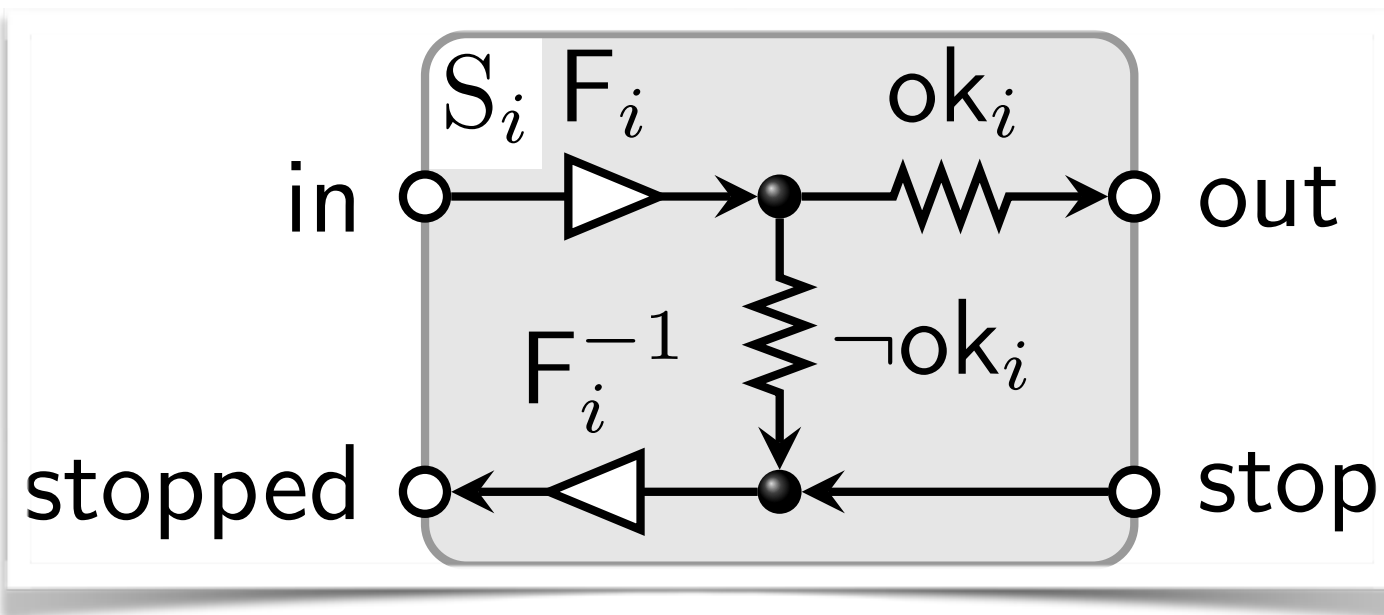
Coordination as constraint satisfaction



Asynchronous transactions



Synchronous transactions



Problem

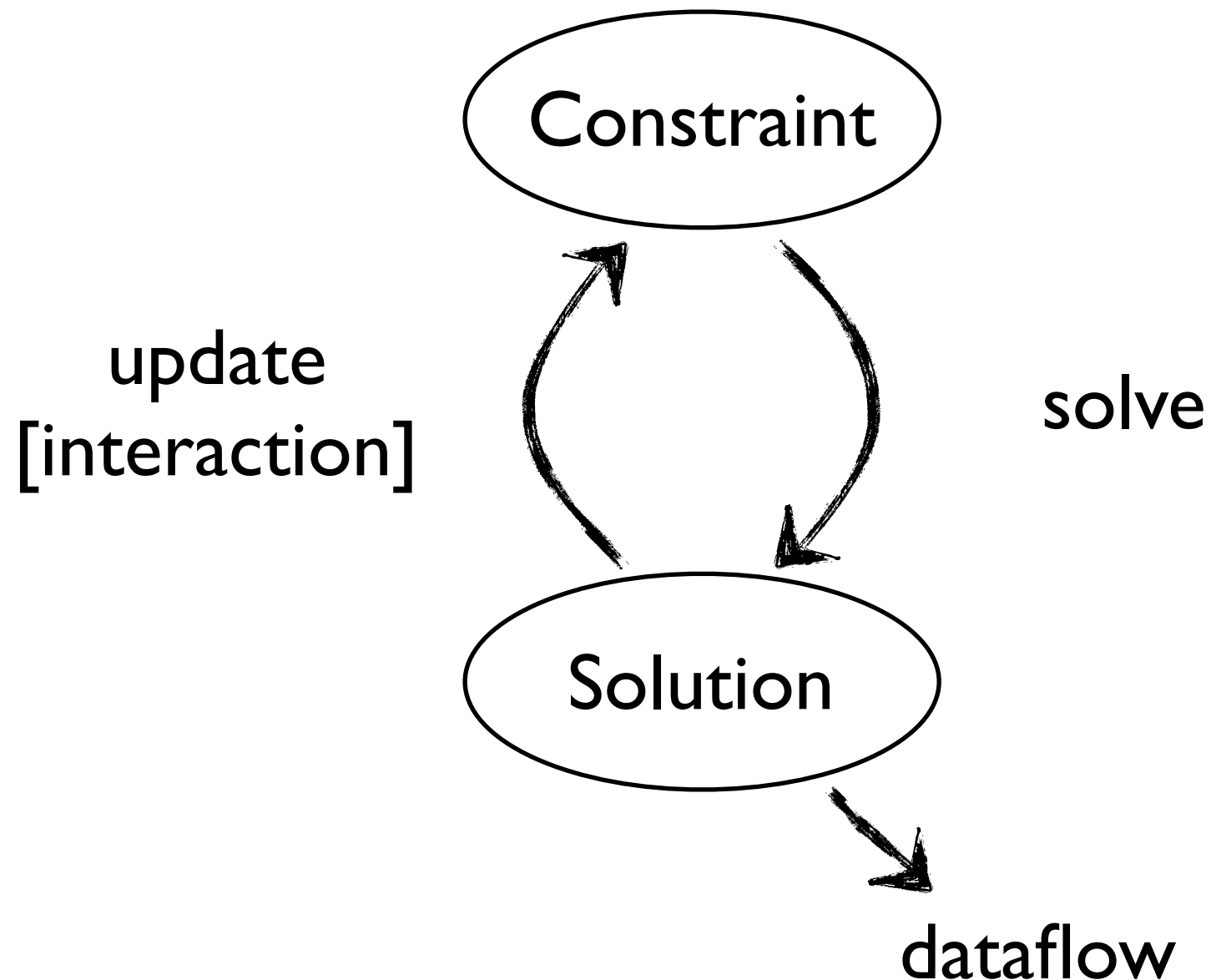
Interaction or Synchrony

but not both

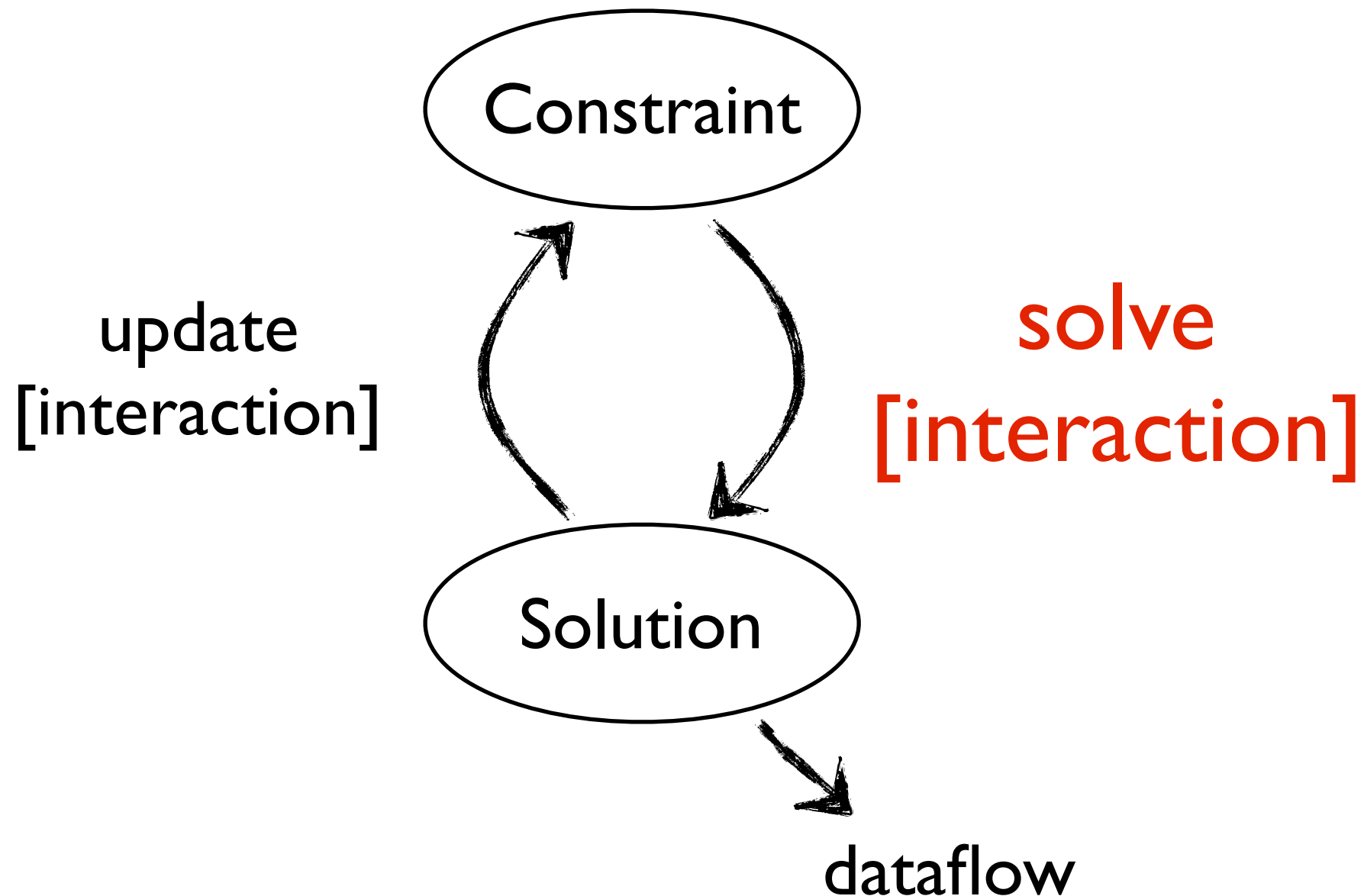
Solution

Interactive Interaction Constraints

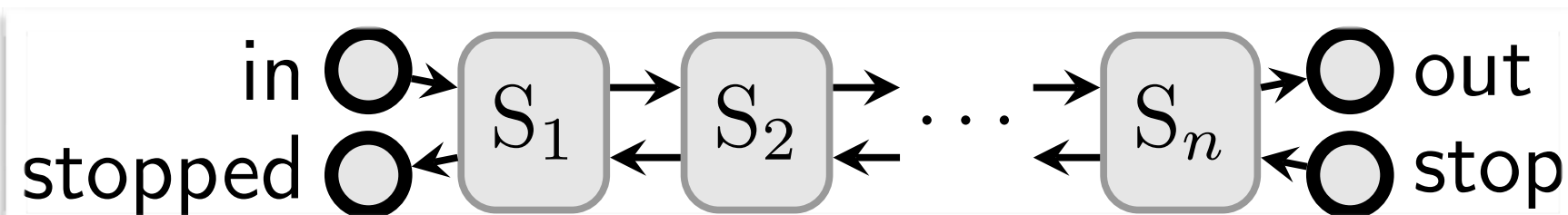
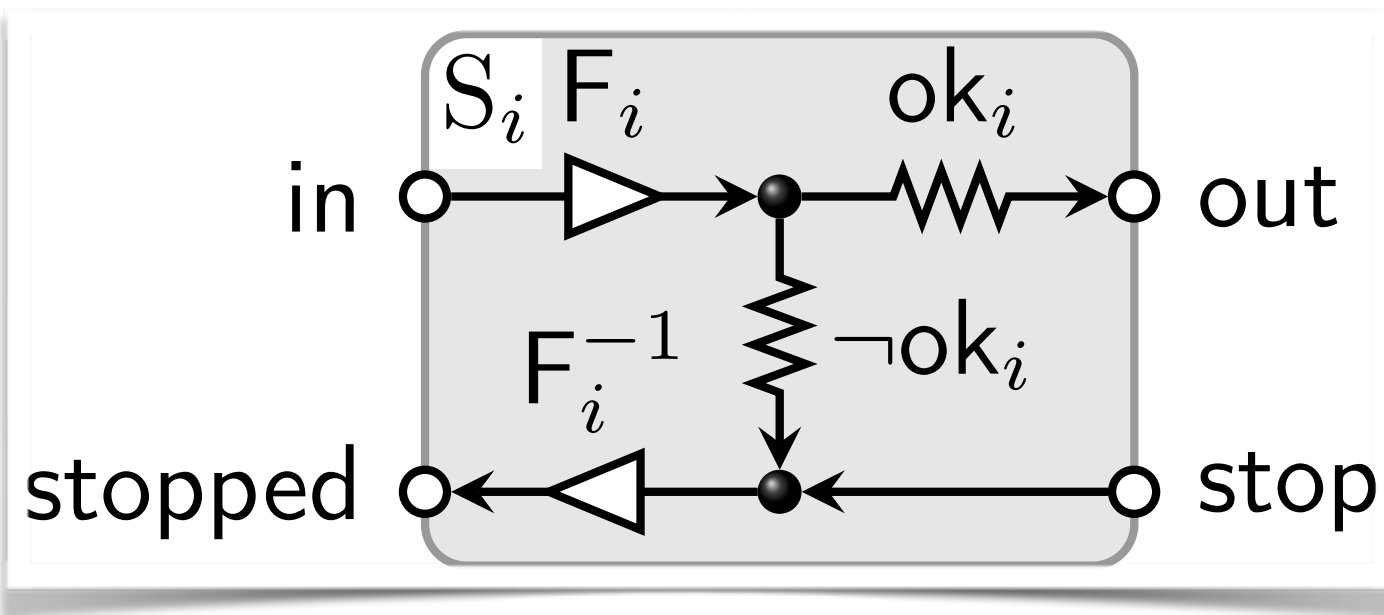
Coordination as constraint satisfaction



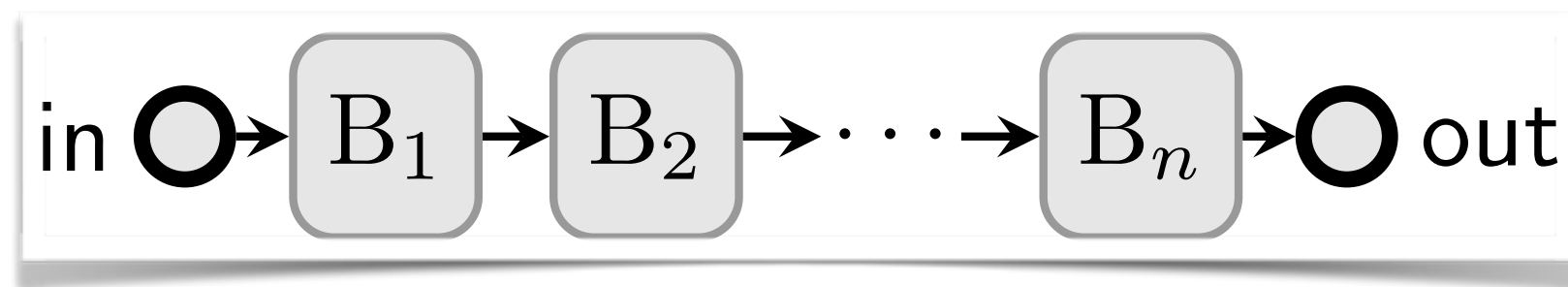
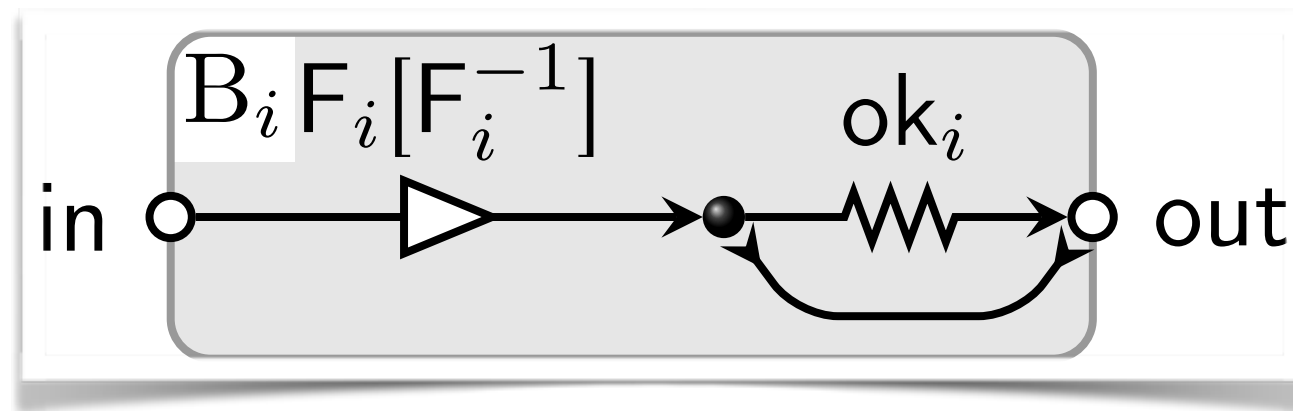
Coordination as **interactive** constraint satisfaction



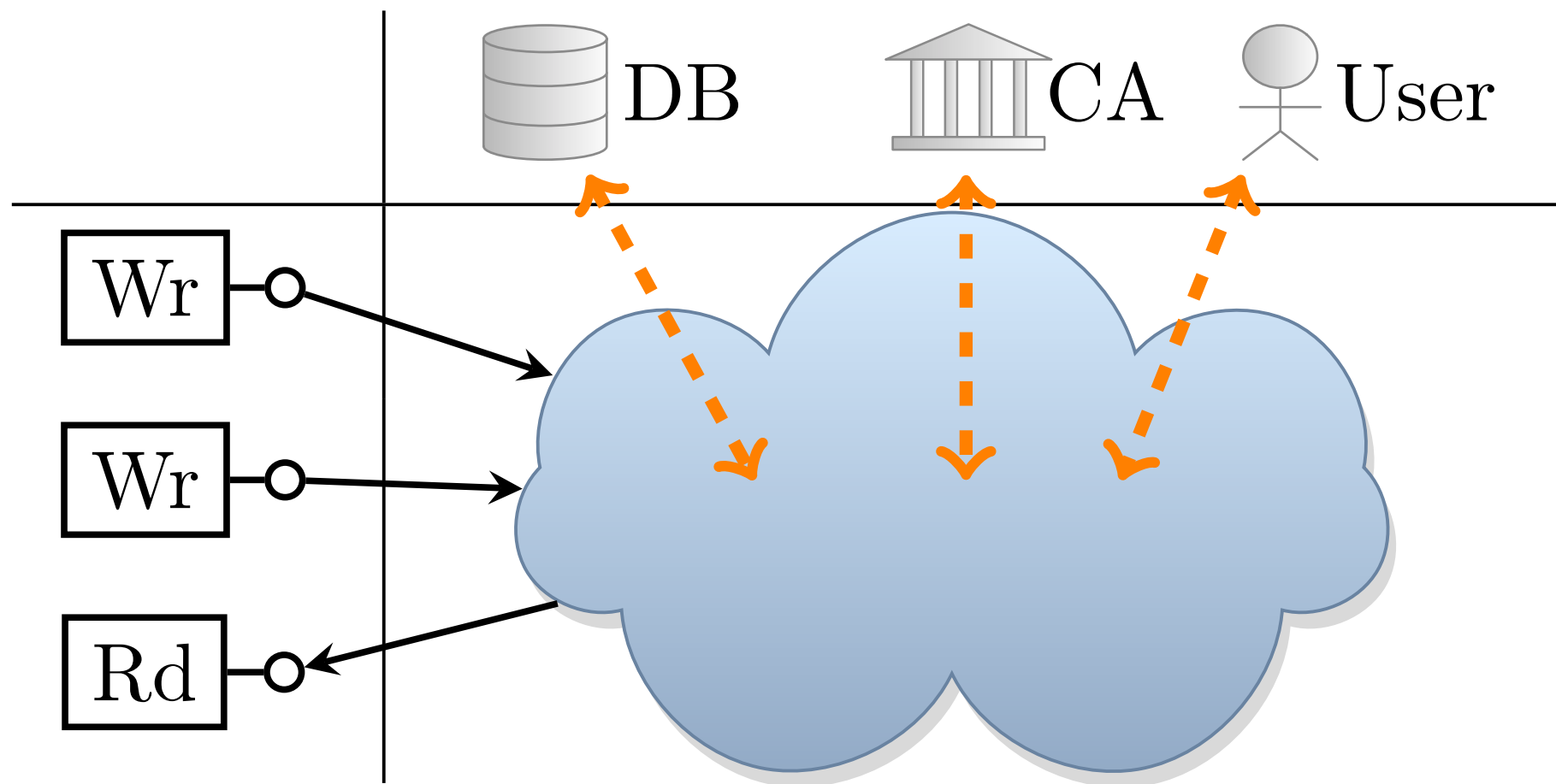
Synchronous transactions



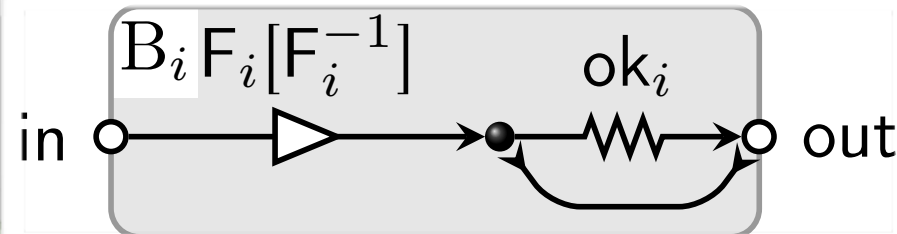
Synchronous transactions



Interactive Interaction Constraints



Scala / Java implementation



```
val f = Function("f") {
  case s: String => /* do something */
}

val finv = Function("f^-1") {
  case s: String => /* do something */
}

val ok = Predicate("ok") {
  case s: String => /* do something */
}

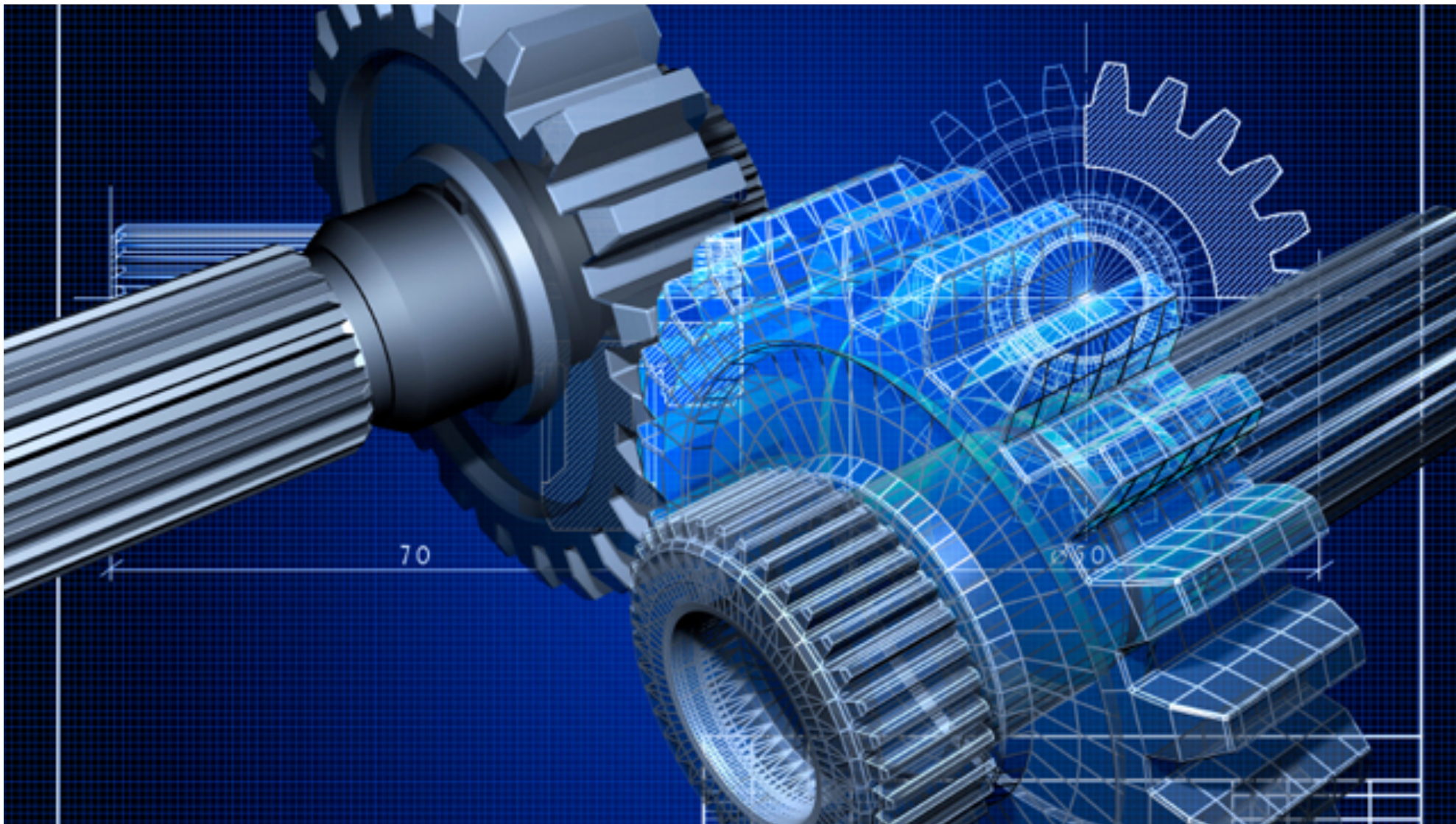
val connector =
  writer("in", List("a", "b")) ++
  transf("in", "x", f, finv) ++
  filter("x", "out", ok) ++
  sdrain("x", "out") ++
  reader("out", 2)

connector.run()
```

```
class Filter(a: String, b: String,
            p: Predicate)
  extends ... {

  def getConstraint = Constraint(
    b --> a,
    b --> (b := a),
    b --> (a <= p),
    (a /\ (a <= p)) --> b
  )

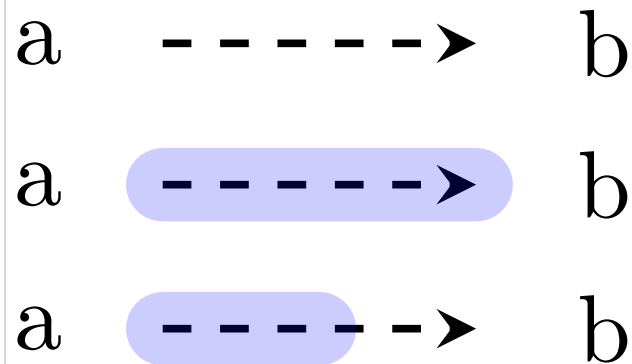
  /* override def update(s:
    Option[Solution]) = ... */
}
```



More about constraints

CC2 as constraints

- [*Colouring*: $\text{End} \rightarrow \{\text{Flow}, \text{NoFlow}\}$]
- *Formula*: Boolean over End
- *Composition* = conjunction

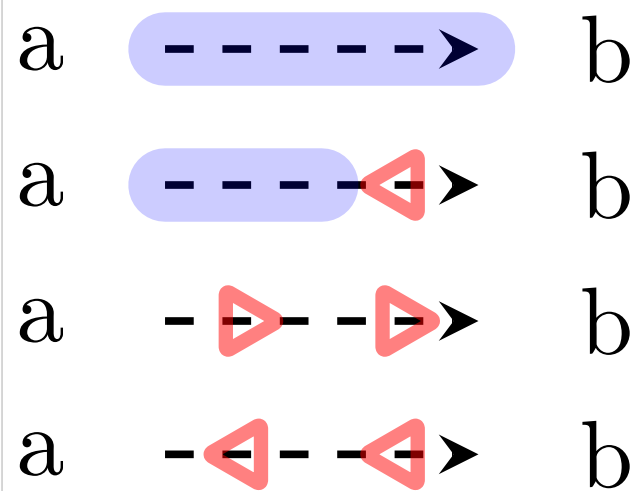


$$\begin{aligned} & (\neg a \wedge \neg b) \vee \\ & (a \wedge b) \vee \\ & (a \wedge \neg b) \end{aligned}$$

$$b \rightarrow a$$

CC3 as constraints

- [*Colouring*: $\text{End} \rightarrow \{\text{Flow}, \text{GiveReason}, \text{GetReason}\}$]
- *Formula*: Boolean over $\text{End}, \text{End}_{\text{src}}, \text{End}_{\text{snk}}$
- a = flow on a ; a_{src} = give reason ; $\neg b_{\text{snk}}$ = get reason

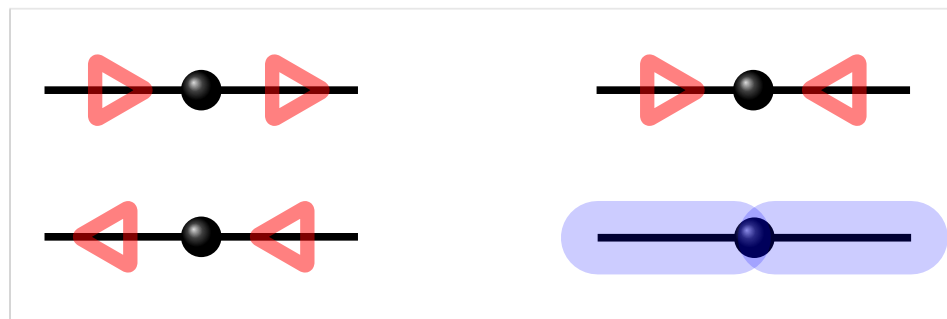


$$\begin{aligned}
 &b \rightarrow a \wedge \\
 &(a \wedge \neg b) \rightarrow (a \wedge \neg b_{\text{snk}}) \wedge \\
 &\neg a \rightarrow (\neg b \wedge \neg a_{\text{src}} \wedge b_{\text{snk}})
 \end{aligned}$$

CC3 as constraints

Composition

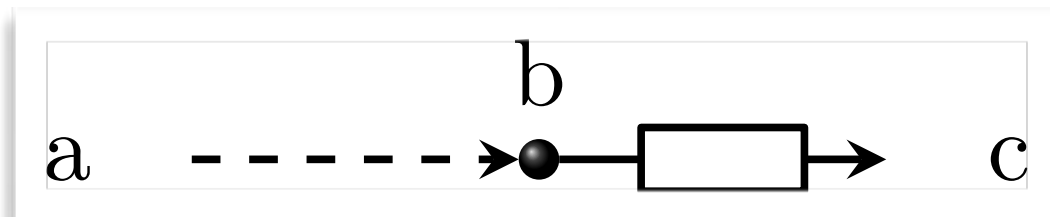
- a = flow on a ; a_{src} = give reason ; $\neg b_{snk}$ = get reason
- one-to-one composition: source to sink ends



$$\forall x \quad \cdot \quad x_{snk} \quad \vee \quad x_{src}$$

CC3 as constraints

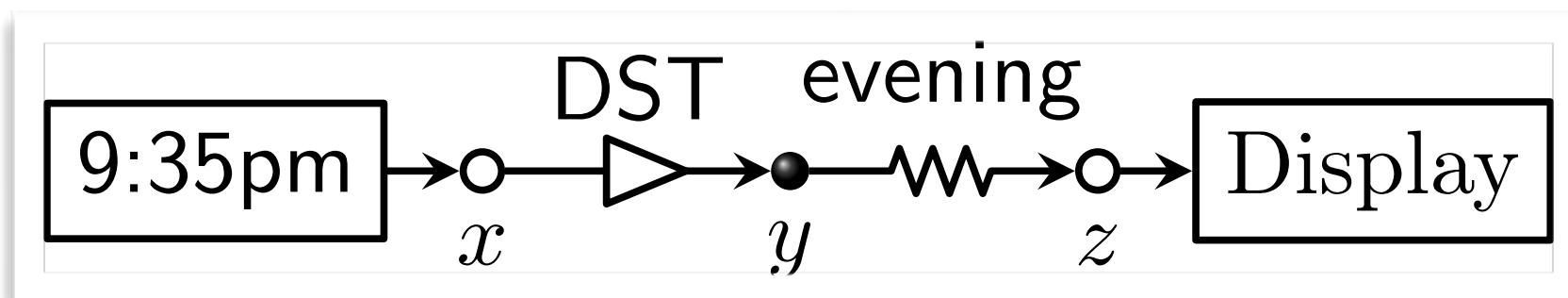
example



$$\begin{aligned}
 \phi &= b \rightarrow a \wedge \neg c \\
 &\quad (a \wedge \neg b) \rightarrow (a \wedge \neg b_{snk}) \wedge \\
 &\quad \neg a \rightarrow (\neg b \wedge \neg a_{src} \wedge b_{snk}) \wedge \\
 &\quad (\neg b \rightarrow \neg b_{src}) \wedge c_{snk} \wedge \\
 &\quad b_{src} \vee b_{snk}
 \end{aligned}$$

$$\begin{aligned}
 \{a \wedge b \wedge \neg c \wedge c_{snk}\} &\models \phi \\
 \{\neg a \wedge \neg b \wedge \neg c \wedge \neg a_{src} \wedge b_{snk} \wedge \neg b_{src} \wedge c_{snk}\} &\models \phi
 \end{aligned}$$

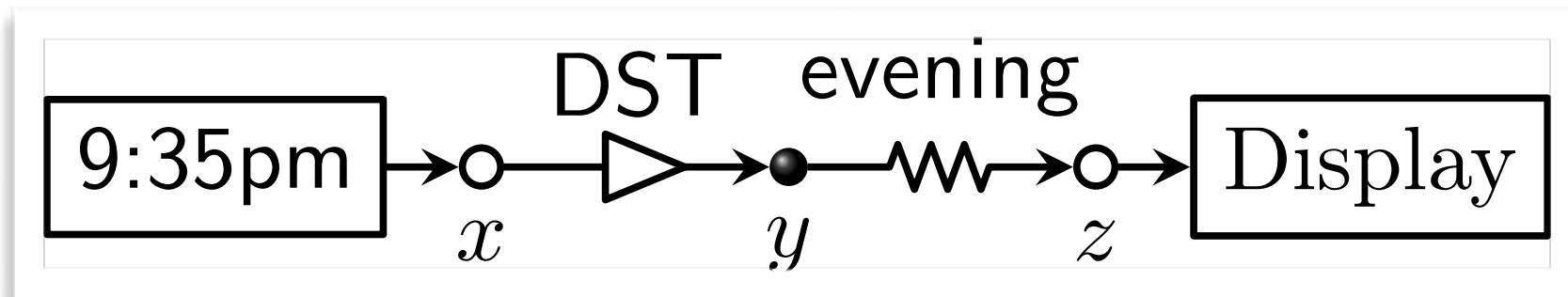
Data constraints



$$\begin{array}{lll} x \rightarrow \hat{x} := 9:35\text{pm} & x \leftrightarrow y & y \rightarrow \hat{y} := \text{DST}(\hat{x}) \\ (y \wedge \text{evening}(\hat{y})) \leftrightarrow z & & z \rightarrow \hat{z} := \hat{y} \end{array}$$

How to solve this?

Predicate abstraction



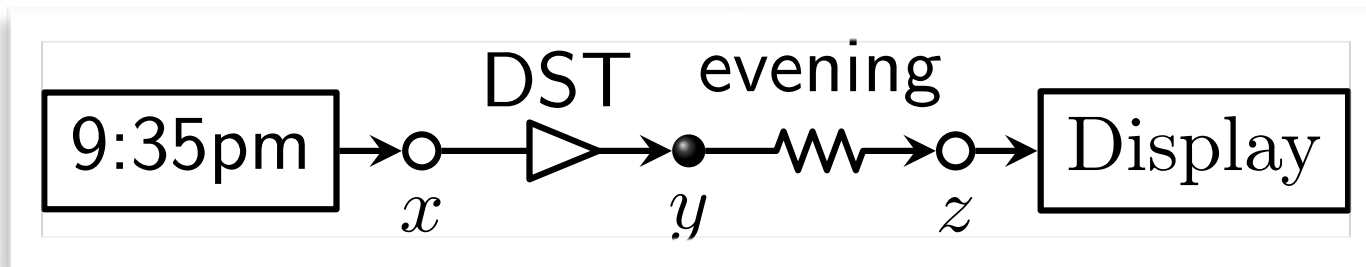
original

$$\begin{array}{lll}
 x \rightarrow \hat{x} := 9:35\text{pm} & x \leftrightarrow y & y \rightarrow \hat{y} := \text{DST}(\hat{x}) \\
 (y \wedge \text{evening}(\hat{y})) \leftrightarrow z & & z \rightarrow \hat{z} := \hat{y}
 \end{array}$$

boolean

$$\begin{array}{ll}
 x \rightarrow \hat{x}_{\text{ev.dst}} := [\text{evening}(\text{DST}(9:35\text{pm}))] & x \leftrightarrow y \\
 y \rightarrow \hat{y}_{\text{ev}} := \hat{x}_{\text{ev.dst}} & (y \wedge \hat{y}_{\text{ev}}) \leftrightarrow z
 \end{array}$$

Interaction via Choco



boolean

$$\begin{aligned}
 x \rightarrow \hat{x}_{\text{ev.dst}} &:= [\text{evening}(\text{DST}(9:35\text{pm}))] & x \leftrightarrow y \\
 y \rightarrow \hat{y}_{\text{ev}} &:= \hat{x}_{\text{ev.dst}} & (y \wedge \hat{y}_{\text{ev}}) \leftrightarrow z
 \end{aligned}$$

interactive

$$\begin{aligned}
 x \rightarrow \text{XPred}(\text{ev.dst}, x, 9:35\text{pm}) & & x \leftrightarrow y \\
 y \rightarrow \hat{y}_{\text{ev}} &:= \hat{x}_{\text{ev.dst}} & (y \wedge \hat{y}_{\text{ev}}) \leftrightarrow z
 \end{aligned}$$

Interaction via Choco

9:35pm

- instance of a Choco constraint
- reacts when x or $\hat{x}_{\text{ev.dst}}$ is instantiated
- $\neg x \Rightarrow \hat{x}_{\text{ev.dst}}$ can be anything
- $x \Rightarrow \hat{x}_{\text{ev.dst}} = \text{ev}(\text{dst}(9:35\text{pm}))$

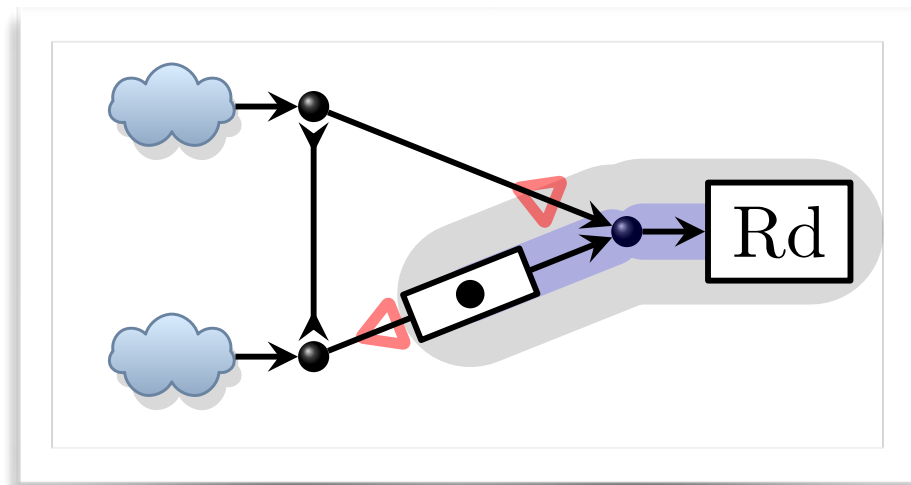
boolean

$$\begin{array}{ll}
 x \rightarrow \hat{x}_{\text{ev.dst}} = [\text{evening}(\text{DST}(9:35\text{pm}))] & x \leftrightarrow y \\
 y \rightarrow \hat{y}_{\text{ev}} := \hat{x}_{\text{ev.dst}} & (y \wedge \hat{y}_{\text{ev}}) \leftrightarrow z
 \end{array}$$

interactive

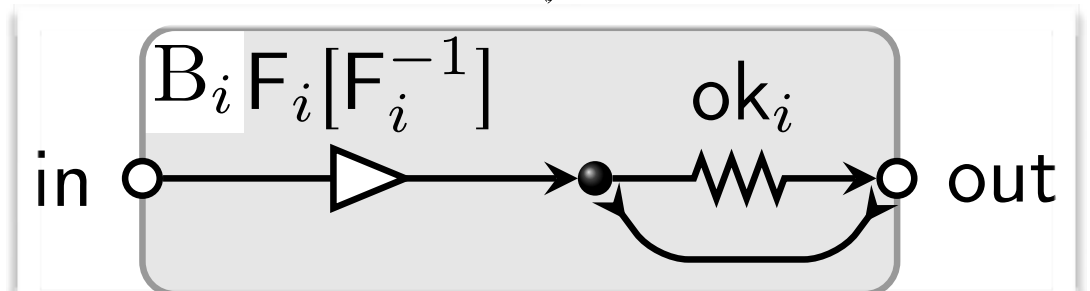
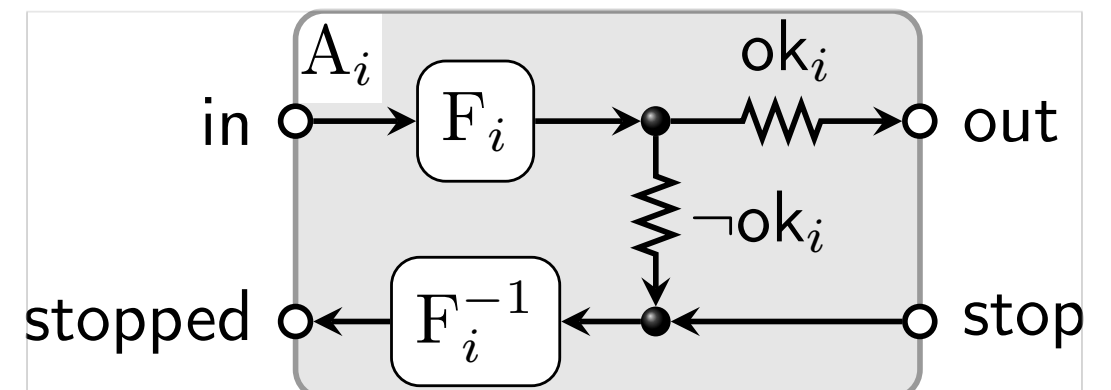
$$\begin{array}{ll}
 x \rightarrow \text{XPred}(\text{ev.dst}, x, 9:35\text{pm}) & x \leftrightarrow y \\
 y \rightarrow \hat{y}_{\text{ev}} := \hat{x}_{\text{ev.dst}} & (y \wedge \hat{y}_{\text{ev}}) \leftrightarrow z
 \end{array}$$

Wrapping up



- Interactive constraint solving
- Expose the atomicity of Reo to components

- Reo and connector colouring



Ongoing experiments

- Avoiding pre-processing (SMT instead of SAT)

$$\exists a, \hat{a} \cdot \psi$$

- Compiling steps:
- Heuristics for identifying potential partial colourings
- Combining local and interactive constraints (probably to a journal)

```

object HotelReservation extends App {

  case class Req(val content:String)

  def srchHotel(i:Int) =
    Function("SearchHotel-"+i){
      case r:Req => i match {
        case 1 => List("F1", "Ibis", "Mercury")
        case 2 => List("B&B", "YHostel")
        case _ => List("HotelA", "HotelB")
      }
    }

  val approve = Predicate("approve"){
    case l:List[String] =>
      println("approve: "+l.mkString(",
        ")+" ". [y,n]")
      readChar() == 'y'
  }

  val book = Function("book"){
    case l : List[String] =>
      println("Options: "+l.mkString(", ")+"
        ". Which one? (1.."+l.length+"")
      val res = readInt()
      l(res-1)
  }

  val cancelB = Function("cancelB"){
    case x => println("canceling "+x+".")
  }

  val invoice = Function("invoice"){
    case x => println("invoice for "+x+".")
  }
}

```

```

val pay = Predicate("paid"){
  case x => if (x == "Ibis") {
    println("paid for Ibis")
    true
  }
  else {
    println("not paid for "+x)
    false
  }
}

// Connector definition
val connector =
  writer("req",List(Req("req1"),
    Req("req2"))) ++
  nexrouter("req",List("h1","h2","h3")) ++
  transf("h1","h1o",srchHotel(1)) ++
  transf("h2","h2o",srchHotel(2)) ++
  transf("h3","h3o",srchHotel(3)) ++
  nmerger(List("h1o","h2o","h3o"),"hs") ++
  filter("hs","ap",approve) ++
  sdrain("hs","ap") ++
  transf("ap","bk",book,cancelB) ++
  monitor("bk","inv",invoice) ++
  filter("inv","paid",paid) ++
  reader("paid",5)

connector.run()
}

```