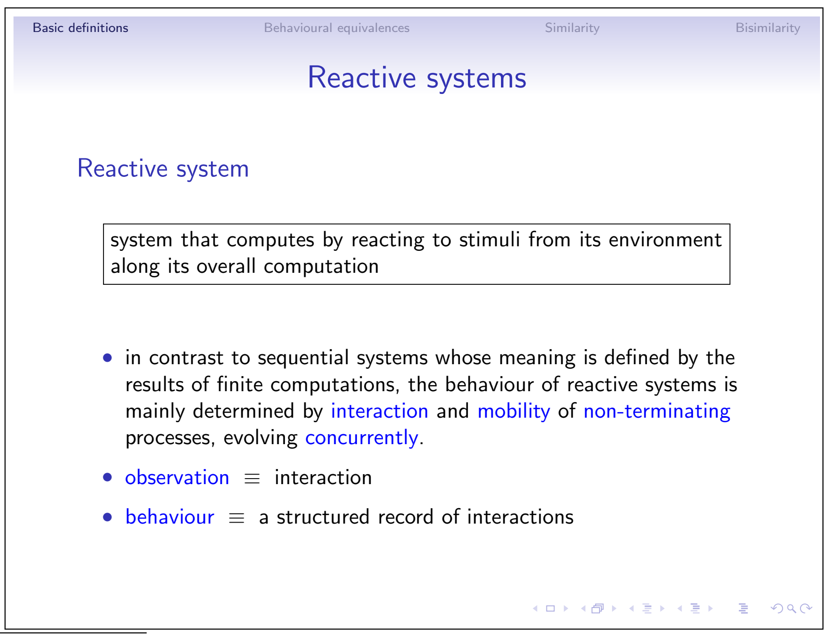
Lecture 2: Introduction to Labelled Transition Systems

Luís Soares Barbosa

Abstract

This lecture offers an introduction to labelled transition systems and their application as a basic semantic structure for non-deterministic reactive systems. Notions of simulation, trace-equivalence and bisimulations are studied in some detail. Students are encouraged to complement the lectures with laboratory, hands-one work with mCRL2.

1 Basic definitions



Lecture notes for *Arquitectura e Cálculo*, MEI profile in *Formal Methods in Software Engineering*, 2014-15.

Labelled Transition System

Definition

A LTS over a set N of names is a tuple $\langle S, N, \longrightarrow \rangle$ where

- $S = \{s_0, s_1, s_2, \ldots\}$ is a set of states
- $\longrightarrow \subseteq S \times N \times S$ is the transition relation, often given as an N-indexed family of binary relations

$$s \xrightarrow{a} s' \;\equiv\; \langle s', a, s \rangle \in \longrightarrow$$

Labelled Transition System

Morphism

A morphism relating two LTS over N, $\langle S, N, \longrightarrow \rangle$ and $\langle S', N, \longrightarrow' \rangle$, is a function $h : S \longrightarrow S'$ st

$$s \xrightarrow{a} s' \quad \Rightarrow \quad h\, s \xrightarrow{a}' h\, s'$$

morphisms preserve transitions

Labelled Transition System

System

Given a LTS $\langle S, N, \longrightarrow \rangle$, each state $s \in S$ determines a system over all states reachable from s and the corresponding restrictions of $\longrightarrow$ and $\downarrow$.

LTS classification

- deterministic
- non deterministic
- finite
- finitely branching
- image finite
- ...

Reachability

Definition

The reachability relation, $\longrightarrow^* \subseteq S \times N^* \times S$, is defined inductively

- $s \xrightarrow{\epsilon}^* s$ for each $s \in S$, where $\epsilon \in N^*$ denotes the empty word;
- if $s \xrightarrow{a} s''$ and $s'' \xrightarrow{\sigma}^* s'$ then $s \xrightarrow{a\sigma}^* s'$, for $a \in N, \sigma \in N^*$

Reachable state

$t \in S$ is reachable from $s \in S$ iff there is a word $\sigma \in N^*$ st $s \xrightarrow{\sigma}^* t$

An alternative characterisation

Coalgebraic characterization (morphism)

A morphism $h : \langle S, \mathsf{next} \rangle \longrightarrow \langle S', \mathsf{next}' \rangle$ is a function $h : S \longrightarrow S'$ st the following diagram commutes

$$\begin{array}{ccc} S \times N & \xrightarrow{\mathsf{next}} & \mathcal{P}S \\ {\scriptstyle h\times \mathsf{id}}\downarrow & & \downarrow{\scriptstyle \mathcal{P}h} \\ S' \times N & \xrightarrow{\mathsf{next}'} & \mathcal{P}S' \end{array}$$

i.e.,

$$\mathcal{P}h \cdot \mathsf{next} \quad = \quad \mathsf{next}' \cdot (h \times \mathsf{id})$$

or, going pointwise,

$$\{h\; x \mid x \in \mathsf{next}\; \langle s, a \rangle\} \quad = \quad \mathsf{next}'\; \langle h\; s, a \rangle$$

An alternative characterisation

Coalgebraic characterization (morphism)

A morphism $h : \langle S, \mathsf{next} \rangle \longrightarrow \langle S', \mathsf{next}' \rangle$

- preseves transitions:

$$s' \in \mathsf{next}\; \langle s, a \rangle \Rightarrow h\; s' \in \mathsf{next}'\; \langle h\; s, a \rangle$$

- reflects transitions:

$$r' \in \mathsf{next}'\; \langle h\; s, a \rangle \Rightarrow \langle \exists\; s' \in S\; :\; s' \in \mathsf{next}\; \langle s, a \rangle :\; r' = h\; s' \rangle$$

(why?)

Comparison

- Both definitions coincide at the object level:

$$\langle s, a, s' \rangle \in T \quad \equiv \quad s' \in \mathsf{next}\ \langle s, a \rangle$$

- Wrt morphisms, the relational definition is more general, corresponding, in coalgebraic terms to

$$\mathcal{P}h \cdot \mathsf{next} \quad \subseteq \quad \mathsf{next}' \cdot (h \times \mathsf{id})$$

How can these notions of morphism be used to compare LTS?

2 Behavioural equivalences

Intuition

Two LTS should be equivalent if they cannot be distinguished by interacting with them.

Equality of functional behaviour

is not preserved by parallel composition: non compositional semantics, cf,

`x:=4; x := x+1` and `x:=5`

Graph isomorphism

is too strong (why?)

Trace

Definition

Let $T = \langle S, N, \longrightarrow \rangle$ be a labelled transition system. The set of traces $\mathsf{Tr}(s)$, for $s \in S$ is the minimal set satisfying

(1) $\epsilon \in \mathsf{Tr}(s)$

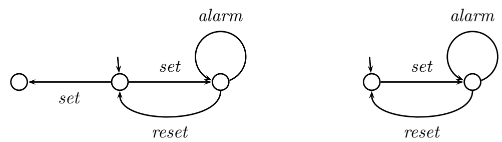
(2) $a\sigma \in \mathsf{Tr}(s) \Rightarrow \langle \exists\, s' \;:\; s' \in S :\; s \xrightarrow{a} s' \wedge \sigma \in \mathsf{Tr}(s') \rangle$

Trace equivalence

Definition

Two states s, r are trace equivalent iff $\mathrm{Tr}(s) = \mathrm{Tr}(r)$
(i.e. if they can perform the same finite sequences of transitions)

Example



Trace equivalence applies when one can neither interact with a system, nor distinguish a slow system from one that has come to a stand still.

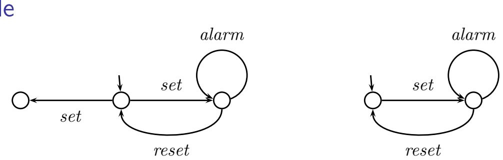
Basic definitions | **Behavioural equivalences** | Similarity | Bisimilarity

Trace equivalence

Definition

Two states s, r are trace equivalent iff $\mathrm{Tr}(s) = \mathrm{Tr}(r)$
(i.e. if they can perform the same finite sequences of transitions)

Example



Trace equivalence applies when one can neither interact with a system, nor distinguish a slow system from one that has come to a stand still.

3 Similarity

Simulation

the quest for a behavioural equality:
able to identify states that cannot be distinguished by any realistic form of observation

Simulation

> A state q simulates another state p if every transition from q is corresponded by a transition from p and this capacity is kept along the whole life of the system to which state space q belongs to.

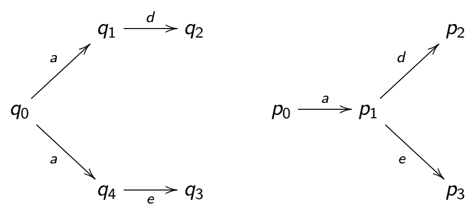
Simulation

Definition

Given $\langle S_1, N, \longrightarrow_1\rangle$ and $\langle S_2, N, \longrightarrow_2\rangle$ over N, relation $R \subseteq S_1 \times S_2$ is a simulation iff, for all $\langle p, q\rangle \in R$ and $a \in N$,

$$(1)\quad p \xrightarrow{a}_1 p' \;\Rightarrow\; \langle \exists\, q' \;:\; q' \in S_2 :\; q \xrightarrow{a}_2 q' \wedge \langle p', q'\rangle \in R\rangle$$

Example



$q_0 \lesssim p_0$ cf. $\{\langle q_0, p_0\rangle, \langle q_1, p_1\rangle, \langle q_4, p_1\rangle, \langle q_2, p_2\rangle, \langle q_3, p_3\rangle\}$

Similarity

Definition

$$p \lesssim q \;\equiv\; \langle \exists\, R \;::\; R \text{ is a simulation and } \langle p, q\rangle \in R\rangle$$

Lemma

The similarity relation is a preorder
(ie, reflexive and transitive)

4 Bisimilarity

Bisimulation

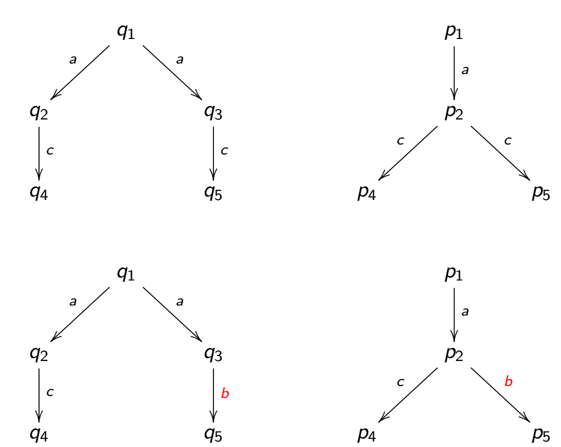
Definition

Given $\langle S_1, N, \longrightarrow_1\rangle$ and $\langle S_2, N, \longrightarrow_2\rangle$ over N, relation $R \subseteq S_1 \times S_2$ is a bisimulation iff both R and its converse R° are simulations.
I.e., whenever $\langle p, q\rangle \in R$ and $a \in N$,

$$(1)\quad p \xrightarrow{a}_1 p' \;\Rightarrow\; \langle \exists\, q' \;:\; q' \in S_2 :\; q \xrightarrow{a}_2 q' \wedge \langle p', q'\rangle \in R\rangle$$

$$(2)\quad q \xrightarrow{a}_2 q' \;\Rightarrow\; \langle \exists\, p' \;:\; p' \in S_1 :\; p \xrightarrow{a}_1 p' \wedge \langle p', q'\rangle \in R\rangle$$

Examples



After thoughts

- Follows a $\forall, \exists$ pattern: p in all its transitions challenge q which is called to find a matchh to each of those (and conversely)
- Tighter correspondence with transitions
- Based on the information that the transitions convey, rather than on the shape of the LTS
- Local checks on states
- Lack of hierarchy on the pairs of the bisimulation (no temporal order on the checks is required)

which means bisimilarity can be used to reason about infinite or circular behaviours.

After thoughts

Compare the definition of bisimilarity with

$p == q$ if, for all $a \in N$

$$(1)\quad p \xrightarrow{a}_1 p' \;\Rightarrow\; \langle \exists\, q' \;:\; q' \in S_2 :\; q \xrightarrow{a}_2 q' \wedge p' == q' \rangle$$
$$(2)\quad q \xrightarrow{a}_2 q' \;\Rightarrow\; \langle \exists\, p' \;:\; p' \in S_1 :\; p \xrightarrow{a}_1 p' \wedge p' == q' \rangle$$

After thoughts

$p == q$ if, for all $a \in N$

$$(1)\quad p \xrightarrow{a}_1 p' \;\Rightarrow\; \langle \exists\, q' \;:\; q' \in S_2 :\; q \xrightarrow{a}_2 q' \wedge p' == q' \rangle$$
$$(2)\quad q \xrightarrow{a}_2 q' \;\Rightarrow\; \langle \exists\, p' \;:\; p' \in S_1 :\; p \xrightarrow{a}_1 p' \wedge p' == q' \rangle$$

- The meaning of $==$ on the pair $\langle p, q \rangle$ requires having already established the meaning of $==$ on the derivatives
- ... therefore the definition is ill-founded if the state space reachable from $\langle p, q \rangle$ is infinite or contain loops
- ... this is a local but inherently inductive definition (to revisit later)

After thoughts

Proof method

To prove that two behaviours are bisimilar, find a bisimulation containing them ...

- ... impredicative character
- coinductive vs inductive definition

Properties

Definition

$$p \sim q \;\equiv\; \langle \exists\, R \;::\; R \text{ is a bisimulation and } \langle p, q\rangle \in R \rangle$$

Lemma

1. The identity relation id is a bisimulation
2. The empty relation $\perp$ is a bisimulation
3. The converse R° of a bisimulation is a bisimulation
4. The composition $S \cdot R$ of two bisimulations S and R is a bisimulation
5. The $\bigcup_{i \in I} R_i$ of a family of bisimulations $\{R_i \mid i \in I\}$ is a bisimulation

Properties

Lemma

The bisimilarity relation is an equivalence relation
(ie, reflexive, symmetric and transitive)

Lemma

The class of all bisimulations between two LTS has the structure of a complete lattice, ordered by set inclusion, whose top is the bisimilarity relation $\sim$.

Properties

Lemma

In a deterministic labelled transition system, two states are bisimilar iff they are trace equivalent, i.e.,

$$s \sim s' \quad \Leftrightarrow \quad \mathrm{Tr}(s) = \mathrm{Tr}(s')$$

Hint: define a relation R as

$$\langle x, y \rangle \in R \ \Leftrightarrow \ \mathrm{Tr}(x) = \mathrm{Tr}(y)$$

and show R is a bisimulation.

Basic definitions Behavioural equivalences Similarity Bisimilarity

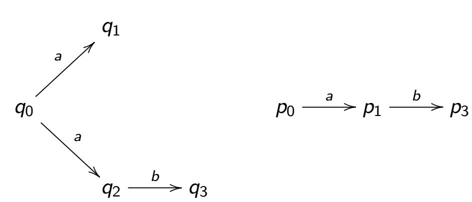
Properties

Warning

The bisimilarity relation $\sim$ is not the symmetric closure of $\lesssim$

Example

$$q_0 \lesssim p_0, \ p_0 \lesssim q_0 \quad \text{but} \quad p_0 \not\sim q_0$$



Notes

Similarity as the greatest simulation

$$\lesssim \triangleq \bigcup\{S \mid S \text{ is a simulation}\}$$

Bisimilarity as the greatest bisimulation

$$\sim \triangleq \bigcup\{S \mid S \text{ is a bisimulation}\}$$