Towards Formal Software Development in VS.NET

MSDN - 14/16th of May, 2002

J.N. Oliveira



SIDEREUS S.A. and DI/UNIV. OF MINHO



MSDN-020514 - p.1/42

Cover Story

Excerpt of article in the CAMBRIDGE EVENING NEWS:

Computer Scientist Gets to the "Bottom" of Financial Scandal

A Cambridge computer professor, Simon Peyton Jones, has made an interesting discovery regarding the Enron collapse. (...) Enron's collapse was due to a nearly impenetrable web of financial contracts that disguised the true financial state of the company (...)

Cover Story (cont.)

(...) Accountants find that even when they are scrupulously honest about the valuation of such contracts there can still be sharp disagreements in regard to the worth of trading reserves, debts, and other components.

Enter Professor Peyton Jones. As part of his research at Microsoft in Cambridge, he developed a computer language for describing and valuing financial contracts. (...)

Cover Story (cont.)

(...) With colleagues Jean-Marc Eber and Julian Seward, they developed a language capable of accurately describing and valuing even the most complex financial instruments. (...)

"While accountants find financial derivatives to be mysterious and difficult, for us they are just ordinary recursive equations,"

says Peyton Jones.

Cover Story (cont.)

(...) "We have been dealing with (...) According to Peyton Jones, his success in the financial world comes from years of research in Haskell (...) Haskell community I would never have been able to do what I've done. It's a jolly wonderful way to program computers"

he stated. (...)

Cover Story (conclusion)

(...)

The Arthur Anderson accounting firm is rumored to have made overtures to Peyton Jones. (...) But Professor Peyton Jones plans to remain where he is.

"I'm flattered that my research has finally been of use to someone but I'm quite happy working on Haskell. Besides, I don't want to have to wear a suit to work every day."

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(CAMBRIDGE EVENING NEWS, 1st of April (!) 2002)

Prof. Peyton Jones' "magic words"

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- ... tools developed by the Haskell community

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In other words:

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and

functional programming

costumer



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costume



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costume



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Why formal / elegant notations?



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From a mobile phone manufacturer:

(...) For each list of calls stored in the mobile phone
(eg. numbers dialed, SMS messages, lost calls), the store operation should work in a way such that (a)
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In Haskell notation:

store :: Call \rightarrow [Call] \rightarrow [Call] store c l = ...

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store :: Call -> [Call] -> [Call]
store c l = [c] ++ l

Notation: x + y means "x catenated with y", eg.

[c] ++ [a,b,c] = [c,a,b,c]

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store :: Call -> [Call] -> [Call]
store c l = [c] ++ filter (/=c) l

Notation: From the **Haskell** Prelude:

filter :: (a -> Bool) -> [a] -> [a] filter p l = [a | a <- l, p a]

(...) such that (a) the more recently a **call** is made the more accessible it is; (b) no number appears twice in a list; (c) only the last 10 entries in each list are stored.

store' :: Call -> \Call] -> [Call]
store' c l = take 10 (store c l)

Notation:

Common practice, in eg. C#

public void store10(string phoneNumber)
{

System.Collections.ArrayList auxList =
 new System.Collections.ArrayList();
auxList.Add(phoneNumber);
auxList.AddRange(
 this.filteratmost9(phoneNumber));
this.callList = auxList;

C[#] version of store (cont.)

```
public System.Collections.ArrayList filteratmost9(string n)
  System.Collections.ArrayList retList =
      new System.Collections.ArrayList();
      int i=0, m=0;
  while((i < this.callList.Count) && (m < 9))</pre>
      if ((string)this.callList[i] != n)
          retList.Add(this.callList[i]);
          m++;
      i++;
  return retList;
```

Comments on C[#] code

Even tolerating code verbosity ...How "good" is this implementation?

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Obs.:

The same requirements in an FM exam paper led to 5 kinds of answer, of which only one (!) was correct!
Comments on C[#] code

Even tolerating code verbosity ...

- How "good" is this implementation?
- Does it meet the 3 properties stated by the mobile phone manufacturer?

Obs.:

- The same requirements in an FM exam paper led to 5 kinds of answer, of which only one (!) was correct!
- Alternatively, FMs provide for <u>correct</u> program construction, eg. by <u>calculation</u>.

Programming by calculation

```
store' c l
```

= . . .

- = take 10 (store c l)
- = take 10 ([c] ++ filter (/=c) l)
- = [c] ++ take 9 filter (/=c) l
- = [c] ++ filteratmost 9 (/=c) l

Notation: calculation stems from formal properties, eg.

take m (x ++ y) = (take m x) ++ (take (m-length x) y)



(implementation)

What

(specification)



(implementation)







Scalling up

Bill of materials (ER):



Bill of materials (SQL)

```
CREATE TABLE COMPONENTS (
Compid CHAR (8) NOT NULL,
CStock NUMBER (10) NOT NULL,
Alarm NUMBER (10) NOT NULL,
Cost NUMBER (3,3) NOT NULL
CONSTRAINT COMPONENTS_pk
PRIMARY KEY(Compid)
);
```

```
CREATE TABLE EQUIPMENTS (
EqId CHAR (8) NOT NULL,
Description CHAR (73) NOT NULL,
EStock NUMBER (10) NOT NULL,
CONSTRAINT EQUIPMENTS_pk PRIMARY KEY (EqId)
);
```

(...)

Bill of materials (Haskell)

The entities:

Bill of materials (Haskell)

The relationships:

```
data Part_of = Part_of_Record {
       comp :: String,
       equip :: String,
       howManyC :: Int
data Sub_Block_of = Sub_Block_of_Record {
       equipL :: String,
       equipS :: String,
       howManyE :: Int
The relational tables themselves:
type Components = Set Component
type Equipments = Set Equipment
type Parts of = Set Part of
```

type Sub_Blocks_of = Set Sub_Block_of

Data processing is functional

Every function

$$f: B \to A$$

is a kind of "data miner": it extracts the A-view of every piece of B-data.

In fact:

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In fact:

data-mining is functional
 document processing (eg. XML) is functional

Haskell at work

Not only ...
HaXML — library for XML processing
HaskelIDB — communication with the data access layer

Haskell at work

Not only ...

- HaXML library for XML processing
- HaskelIDB communication with the data access layer

but also

- HaskellScript library for the integration of Haskell with COM/ActiveX.
- HSP a compositional alternative to ASP/PHP
- HDirect IDL compiler helping to interact with external code.

Need for interoperability

- Software architectures are getting more and more complex
- Need to reuse / costumize pre-existing software components
- Programming = "gluing" software together

Rôle of interoperability

Interoperability in .NET

The (old) idea of UCSD P-code, AS-400, JMV ...

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- Microsoft Research and the .NET product group have attracted a number of computer scientists and language implementors to improve support for a wide range of programming paradigms
- Thus novel features in CLI such as "tail-calls" (the tail.call instruction) saving stack space in recursion-based languages such as Haskell.

Haskell for .NET

Haskell.NET = Haskell (GHC) + Mondrian

Glossary:

- GHC— Glasgow Haskell compiler
- Mondrian— OO/functional hybrid used to communicate with .NET
- hs/ms Haskell/Mondrian source code
- mc Mondrian internal code

About Mondrian

 Mondrian — is a non-strict functional language designed for an OO environment

- its syntax is a meld of that of Haskell and Java/C[#]
- disjoint union types are modelled using subtyping, eg

class List {}
class Nil extends List {}
class Cons extends List
{ head : Object;
 tail : List Object
};

Mondrian (cont.)

Object (the topmost class) plays the rôle of a type variable, cf. parameterization.

- functions are compiled to classes embodying a standard evaluation method
- supports concurrent programming using threads and exceptions
- provides access to "foreign" language objects



Haskell.NET in a diagram



Haskell.NET in a diagram



Mobile phone revisited

```
using mondrian.prelude;
using mondrian.runtime;
using store;
• • •
public class SMain
       public static void store10(string c)
       {
             // Store new call in current list
             callList = store.Apply(c,callList) ;
             // Keep only 10 in current list
             callList = take.Apply(10,callList);
```

Potential of .NET for FMs

Aim of the Programming Principles and Tools (PPT) group at MR:

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 Preliminary ideas for Haskell/Mondrian scripting were presented by Erik Meijer (Program Manager in the CLR Group) at 6408.70aAFP'98 (Braga) organized by DI/UMinho.

Are FMs cost-effective?

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In favour:

- FM = discipline, rigour and good documentation ("safety net" for HR mobility)
- FM+FP = rapid prototyping, early feedback on what one is doing
- **FM+FP+CLI** = "time to market" integration
- FM = the only way to complex problems

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Against:

Human factorsLack of FM-trained people

Are FMs a "may" or a "must"?

Control and automation:

- Safety-critical software development requires FMs
- Dependable computing

Data processing services are becoming critical:

- Many IS servers required to be available 365 \times 24 hour / year
- Poor-quality data lead to wrong management decisions
Our background

By 2004: 20 years of FM teaching at the Univ. of Minho \sim 10 years ago: Industrial application of FMs based on FP tested at INESC-BRAGA Spin-off of INESC-BRAGA (1996): **SIDEREUS S.A.** - Rigorous Solutions for Software Systems (Porto)

Interest in .NET

 Sidereus and DI/UM are experimenting with the .NET as a platform for formal/informal tool integration

- Some experiments/products follow:
 - ADO.NET
 - KMig
 - K-reverse
 - FRMS

Modelling ADO.NET

(...) The DataSet is an in-memory cache of data retrieved from a database (...)

```
data DataSet = DataSet {
```

tables :: DataTableCollection, relations :: DataRelationCollection, caseSensitive :: Bool }

(...) While DataTable objects contain the data, the DataRelationCollection allows you to navigate through the table hierarchy (...)

type DataRelationCollection =

Map RelationId DataRelation

From ADO.NET to KMig

- KMig is a data-migration package developed by Sidereus which interoperates with DTS.
- Data-migrations are formally specified in KMig internal language (M2L) and so can be checked for correctness before they are executed.
- KMig and M2L have been fully formally specified (most work carried out at specification level)

K-Reverse and FRMS

Semi-automatic tool for relational database reverse specification.

Can be used for data-quality checking, re-documentation, re-engineering (in connection with KMig).

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Semi-automatic tool for relational database reverse specification.

- Can be used for data-quality checking, re-documentation, re-engineering (in connection with KMig).
- FRMS is an advanced search engine incorporating fuzzy reciprocal matching.
- FRMS requirements are complex: it would have been a risk to approach them informally...

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Closing

Ist of April joke by John Peterson (maintainer of haskell.org) inspired in

> Simon L. Peyton Jones, Jean-Marc Eber, and Julian Seward. Composing contracts: an adventure in financial engineering, functional pearl. In International Conference on Functional Programming, pages 280–292, 2000.

(thanks to Andrei Serjantov for the broadcasting...)

More about Microsoft Research PPT group: http:/research.microsoft.com/ppt/ More about FMs http://www.fmeurope.org More about Haskell http://www.haskell.org More about Mondrian http://www.mondrian-script.org More about us http://www.sidereus.pt http://www.di.uminho.pt